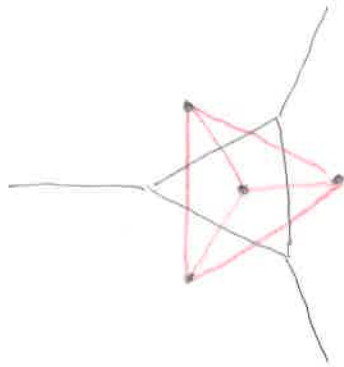


Delaunay Triangulation

(L11)



Connect two sites iff they share a Voronoi Edge
 i.e. $D(P)$ is the planar dual of $V(P)$

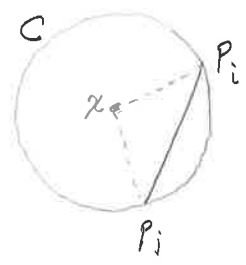
Properties (1) $P_i P_j$ is Delaunay edge iff (by def)

$V(P_i)$ and $V(P_j)$ share Vor. Edge iff

\exists point x (not a site) that is closer (and equal distance) to P_i and P_j than any other site iff

\exists Circle C (with center x) through P_i and P_j that is empty of other sites

Note: $\overline{xP_i}$ is in $V(P_i)$
 $\overline{xP_j}$ is in $V(P_j)$



(2) No Delaunay Edges Cross.

Suppose $P_i P_j$ and $P_k P_l$ cross

P_k and P_l must be outside circle C

$\Rightarrow P_k$ and P_l outside $\Delta P_i x P_j$

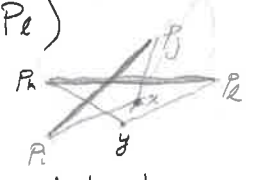
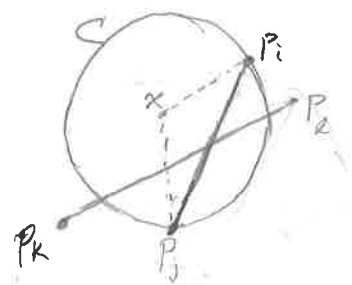
$\Rightarrow \overline{P_k P_l}$ intersects $\overline{xP_i}$ or $\overline{xP_j}$

Similarly $\overline{P_i P_j}$ intersects $\overline{yP_k}$ or $\overline{yP_l}$ (for empty circle with center y through P_k and P_l)

\Rightarrow one of $\overline{xP_i}$ or $\overline{xP_j}$ intersects $\overline{yP_k}$ or $\overline{yP_l}$

Say...

But $\overline{xP_i} \in V(P_i)$ and $\overline{yP_k} \in V(P_k)$
 $\Rightarrow \Leftarrow$

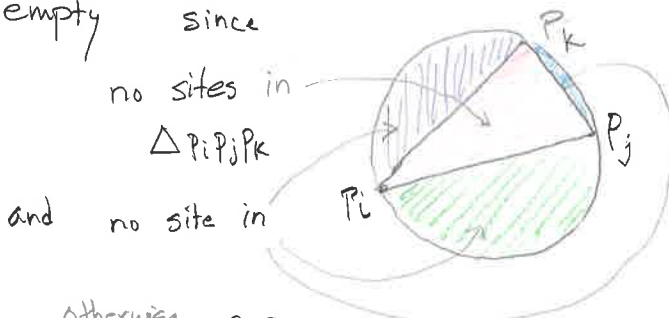


Hard to draw b/c impossible.

(3) Sites P_i, P_j, P_k form a face in Del. Δ tion iff there is an empty circle through them

\Leftarrow empty circle implies $\overline{P_i P_j}, \overline{P_j P_k}, \overline{P_k P_i}$ are Del. Edges and nothing can cross them so its a face

\Rightarrow if $P_i P_j P_k$ is a face then $\bigcirc_{P_i P_j P_k}$ is empty since



no sites in $\Delta_{P_i P_j P_k}$

and no site in

otherwise

$P_i P_j$ ~~edge~~
 $P_j P_k$ ~~edge~~
 $P_k P_i$ ~~edge~~

would not be Del Edge

Alternative definition of Delaunay triangulation

All triangles* whose bounding circle contains no other sites

* assuming no 4 colinear points

Fortune's Sweep parabolic front

Incremental Delaunay