# Two Technique Papers on High Dimensionality

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## Papers

- A. Morrison, G. Ross, and M. Chalmers. Fast multidimensional scaling through sampling, springs and interpolation. In *Information Visualization*, pages 68-77, 2003.
- F. Jourdan and G. Melançon. Multiscale hybrid MDS. In *Intl. Conf. on Information Visualization (London)*, pages 338-393, 2004.
- well written, clear, appropriately detailed
- High-dim and MDS can be complicated

### Dimensionality reduction

- Mapping high-dimensional data to 2D space
- Could be done many different ways
- Different techniques satisfy different goals
- Familiar example projection of 3D to 2D preserves geometric relationships
- Abstract data may not need that

#### Multidimensional scaling (MDS)

- Display multivariate abstract point data in 2D
  - Data from bioinformatics, financial sector, etc.
  - No inherent mapping in 2D space
  - □ *p*-dim embedding of *q*-dim space (p < q) where inter-object relationships are approximated in low-dimensional space
- Proximity in high-D -> proximity in 2D
  - High-dim distance between points (similarity) determines relative (x,y) position
  - Absolute (x,y) positions are not meaningful
- Clusters show closely associated points

#### Multidimensional scaling (MDS)

- Eigenvector analysis of N x N matrix O(N<sup>3</sup>)
  Need to recompute if data changes slightly
- Iterative O(N<sup>2</sup>) algorithm Chalmers, 1996
- This paper  $O(N\sqrt{N})$
- Next paper O(N log N)

#### Multidimensional scaling (MDS)

#### Proximity data

- □ In social sciences, geology, archaeology, etc.
- □ E.g. library catalogue query find similar points
  - Multi-dimensional scatterplot not possible
- Want to see clusters, curves, etc.
  - Features that stand out from the noise
- Distance function
  - Typically use Euclidean distance intuitive

## Spring models

- Used instead of statistical techniques (PCA)
  - Better convergence to optimal solution
  - Iterative steerable Munzner et al, 2004
- Good aesthetic results symmetry, edge lengths
- Basic algorithm  $O(N^3)$ 
  - Start: place points randomly in 2D space
  - Springs reflect diff btwn high-D and 2D distance
  - #iterations required is generally O(N)

#### Chalmers' 1996 algorithm

- Approximate solution works well
- Caching, stochastic sampling  $O(N^2)$ 
  - Perform each iteration in O(N) instead of  $O(N^2)$
  - Keep constant-size set of neighbours
  - Constants as low as 5 worked well
- Still only worked on datasets up to few 1000s

#### Hybrid methods of clustering and layout

- Diff clustering algorithms have diff strengths
  Kohonen's self-organising feature maps (SOM)
  K-means iterative centroid-based divisive alg.
- Hybrid methods have produced benefits
- Neural networks, machine learning literature

#### New hybrid MDS approach

- Start: run spring model on subset of size  $\sqrt{N}$ Completes in O(N) ( $O(\sqrt{N} \cdot \sqrt{N})$ )
- For each remaining point:
  - Place close to closest 'anchor'
  - Adjust by adding spring forces to other anchors
- Overall complexity  $O(N\sqrt{N})$

#### Experimental results

- 3-D data sets: 5000 50,000 points
- 13-D data sets: 2000 24,000 points
- Took less than 1/3 the time of the  $O(N^2)$
- Achieved lower stress when done
- Also compared against original O(N<sup>3</sup>) model
  9 seconds vs. 577; and 24 vs. 3642
  - Achieved much lower stress (0.06 vs. 0.2)

#### Experimental results



*Figure 3.* Run time to completion for different sizes of 3D 'S' data.



*Figure 4.* Stress of completed layout over different sizes of 3D 'S' data.





#### Future work

#### Hashing

- Pivots Morrison, Chalmers, 2003 Achieved  $O(N\sqrt[4]{N})$
- Dynamically resizing anchor set
- Proximity grid
  - Do MDS, then transform continuous layout into discrete topology

## Jourdan and Melançon

- Multiscale hybrid MDS
- Extension of previous paper
- Achieves O(N log N) time complexity
- Good introduction of Chalmers et al papers
- Like Chalmers, begins by embedding subset S of size  $\sqrt{N}$

Improving parent-finding strategy

- Select constant-size subset  $P \subset S$
- For each p in P create sorted list  $L_p$
- For each remaining point *u*, binary search L<sub>p</sub> for point u<sub>p</sub> as distant from *p* as *u* is
  Implies that *u* and u<sub>p</sub> are very close
- Place u according to location of  $u_p$

### Comparison



Figure 2. Comparison of the  $N^{5/4}$  and  $N \log N$  curves on the scale  $10^3 - 10^6$ .



Figure 3. Comparison of the  $N^{1/4}$  and  $\log N$  curves on the scale  $10^3 - 10^6$ .

#### Chalmers et al is better for N < 5500</p>

Main diff is in parent-finding, represented by Fig. 3

#### Comparison



Figure 4. Comparing the actual time spent on the parent-finding phase.



Figure 5. Comparison of Stress for the  $\mathcal{O}(N^{3/2})$ ,  $\mathcal{O}(N^{5/4})$  (grayed) and the  $\mathcal{O}(N \log N)$  (darker) MDS strategies.

Experimental study confirms theoretical results
 This technique becomes better for N > 70,000

## Quality of output



Figure 6. Small world network induced from randomly selected points in 2D (500 node elements).



Figure 7. MDS output obtained from the network in Figure 6.

- MDS theory uses stress to objectively determine quality of placement of points
- Subjective determinations can be made too
  - □ 2D small world network example (500 80,000 nodes)

#### Multiscale MDS



Figure 8. Comparison of Stress for all four MDS algorithms considered here. The darker curve at the bottom reports Stress values reached by the multiscale MDS.



Figure 9. Comparison of the actual running time for all four MDS algorithms considered here.

 Recursively defining the initial kernel set of points can yield much better real-time performance

#### Conclusions and future work

- Series of results yielding progressively better time complexities for MDS
- 2D mappings provide good representations
- Further examination of multiscale approach
- User-steerable MDS could be fruitful