

Introduction to Security

(better late than never!)

Slides borrowed from CS 240 by **Marco** Canini
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Selected content adapted from D. Boneh.

How Systems Fail

Systems may fail for many reasons, including:

- **Reliability** deals with accidental failures
- **Usability** deals with problems arising from operating mistakes made by users
- **Security** deals with **intentional** failures created by **intelligent** parties
 - Security is about computing in the presence of an adversary
 - But **security, reliability, and usability** are all related

The computer security problem

Two factors:

- **Lots of buggy software** (and gullible users)
- **Money can be made from finding and exploiting vulnerabilities**

1. Marketplace for vulnerabilities
2. Marketplace for owned machines (PPI) ← **Pay per install**
3. Many methods to profit from owned client machines

current state of computer security

Why own machines:

1. IP address and bandwidth stealing

Attacker's goal: look like a random Internet user

Use the IP address of infected machine or phone for:

- **Spam** (e.g. the storm botnet)

Spamalytics: 1:12M pharma spams leads to purchase
 1:260K greeting card spams leads to infection

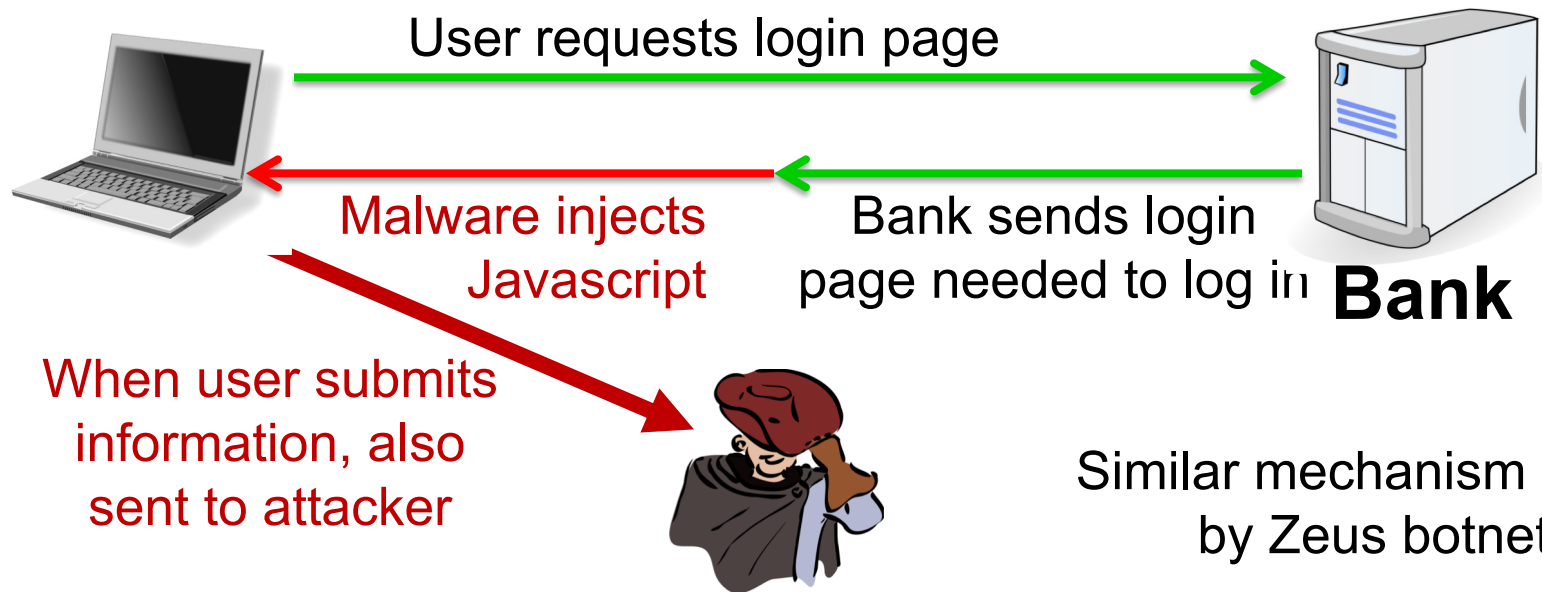
- **Denial of Service:** Services: 1h (20\$), 24h (100\$)
- **Click fraud** (e.g. Clickbot.a)

Why own machines:

2. Steal user credentials

keylog for banking passwords, web pwds., gaming pwds.

Example: SilentBanker (and many like it)



Why own machines:

3. Spread to isolated systems

Example: **Stuxnet**

Windows infection ⇒

Siemens PCS 7 SCADA control software on Windows ⇒

Siemens device controller on isolated network

Challenges: What is “Security”?

- What does **security** mean?
 - Often the hardest part of building a secure system is figuring out what security means
 - What are the **assets** to protect?
 - What are the **threats** to those assets?
 - Who are the **adversaries**, and what are their **resources**?
 - What is the **security policy or goals**?
- Perfect security does not exist!
 - Security is not a binary property
 - Security is about risk management

Current events, security reviews, and other discussions are designed to exercise our thinking about these issues.

Theme 1: Security Mindset

- Thinking critically about designs, challenging assumptions
- Being curious, thinking like an attacker
- “That new product X sounds awesome, I can’t wait to use it!” versus “That new product X sounds cool, but I wonder what would happen if someone did Y with it...”
- Why it’s important
 - Technology changes, so learning to think like a security person is more important than learning specifics of today
 - Will help you design better systems/solutions
 - Interactions with broader context: law, policy, ethics, etc.

Example



Example – What Do You See?



Example – What Do You See?



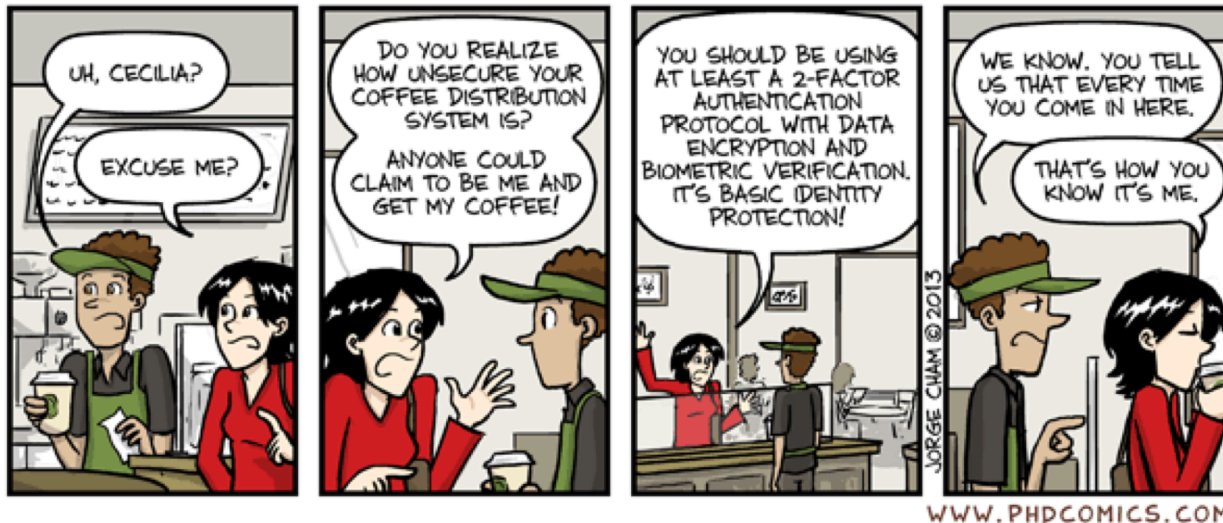
“Security is mostly a superstition” –

Helen Keller (1880-1968), American writer and activist

- Security is all about trade-offs
 - Performance
 - Cost
 - Usability
 - Functionality
- The right question is: how do you know when something is secure enough?
 - Manage security risks vs benefits
 - Requires understanding of the trade-offs involved

How to think about trade-offs?

- What are you trying to protect? How valuable is it?
 - Nuclear missile launch station vs. ... coffee machine



- In what way is it valuable?
 - May be important only to one person (e.g. private e-mail or passwords)
 - May be important because accurate and reliable (e.g. bank's accounting logs)
 - May be important because of a service it provides (e.g. Google's web servers)

High level plan

- Policy: the goal you want to achieve
 - e.g. only Alice should read file F
- Threat model: assumptions about what the attacker could do
 - e.g. can guess passwords, cannot physically grab file server
 - Better to err on the side of assuming attacker can do something
- Mechanism: knobs that your system provides to help uphold policy
 - e.g. user accounts, passwords, file permissions, encryption
- Resulting goal: no way for adversary within threat model to violate policy
 - Note that goal has nothing to say about mechanism

Security goals

- Prevent common vulnerabilities from occurring (e.g. buffer overflows)
 - Recover from attacks
- Traceability, accountability and auditing of security-relevant actions
 - Monitoring
- Detect attacks
 - Privacy, confidentiality, anonymity
 - Protect secrets
- Authenticity
 - Needed for access control, authorization, etc.
- Integrity
 - Prevent unwanted modification or tampering
- Availability and reliability
 - Reduce risk of DoS

Classic CIA triad

- **Confidentiality**

- NO unauthorized disclosure of information

- E.g. a credit card transaction system attempts to enforce confidentiality by encrypting credit card details over the Internet and in the transaction processing network

- **Integrity**

- NO unauthorized information modification

- E.g. traditional Unix file permissions can be an important factor in single system measures for protecting data integrity

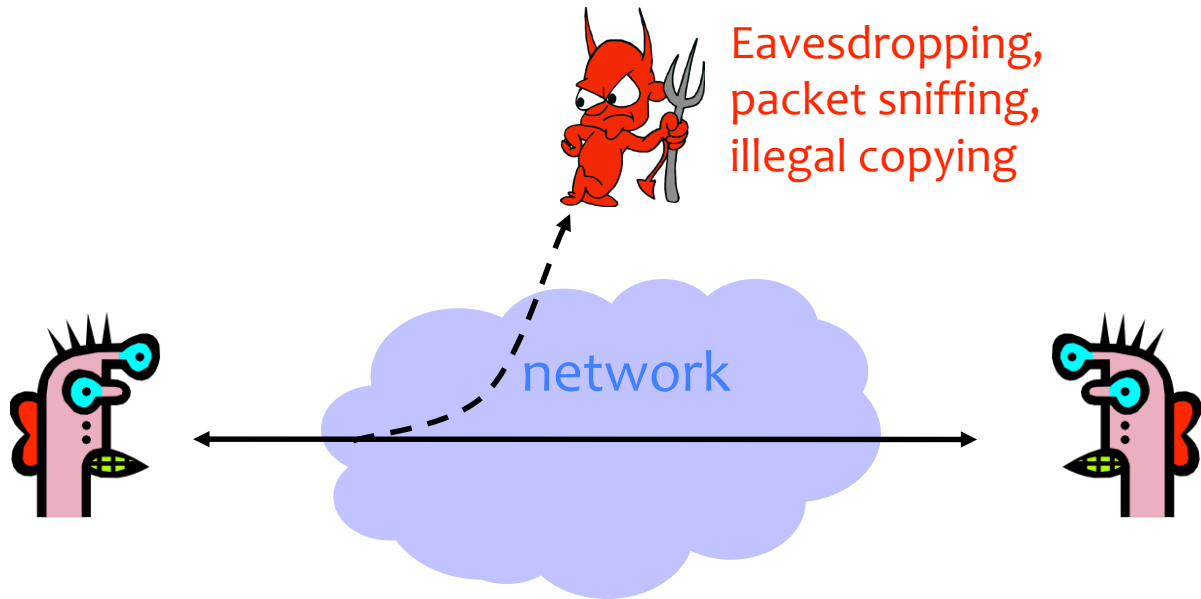
- **Availability + Authenticity (non-standard)**

- Information or system remains available despite attacks

- High availability systems aim to remain available at all times, preventing disruptions due to power outages, upgrades, hardware failures, Denial of Service (DoS) attacks, ...

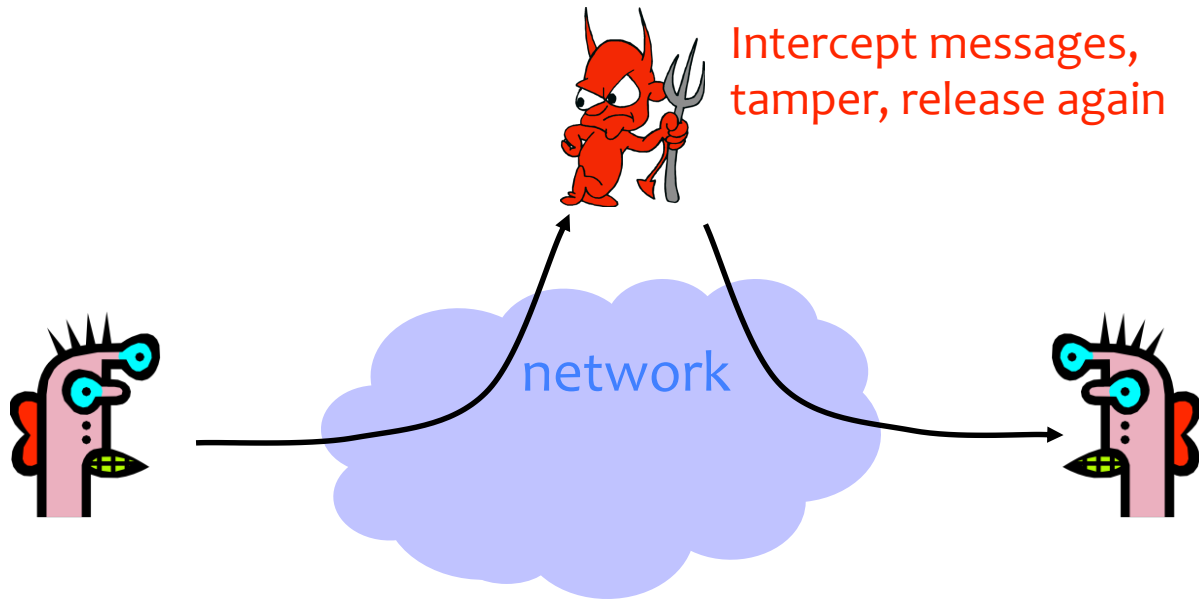
Confidentiality (Privacy)

- Confidentiality is concealment of information.



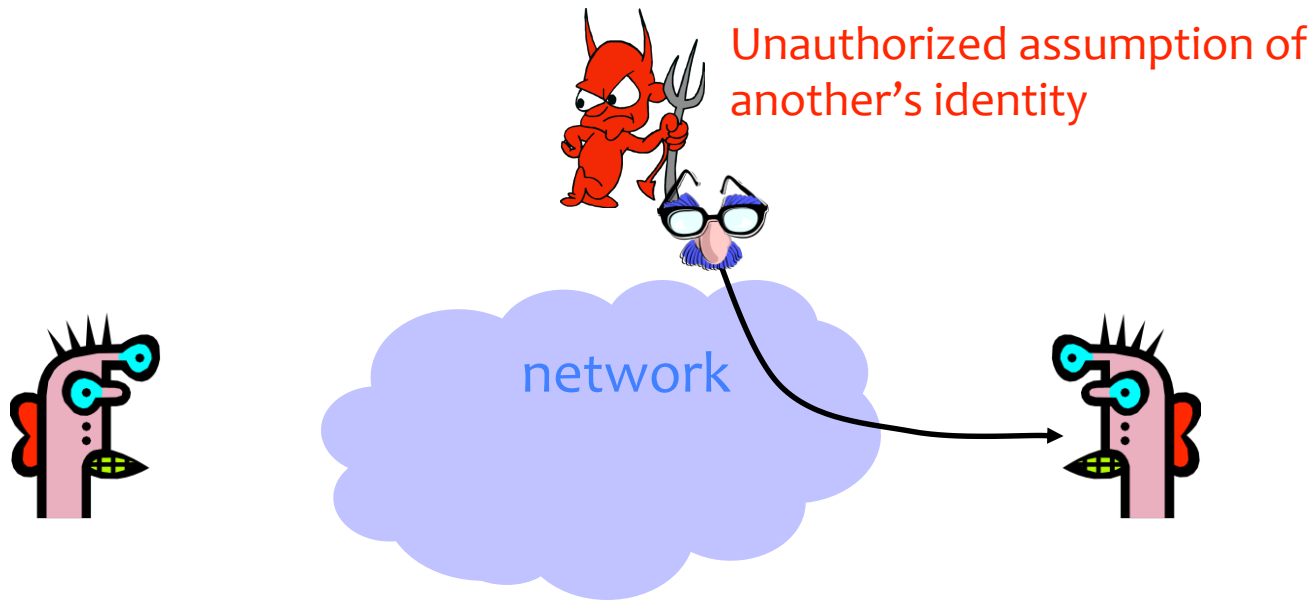
Integrity

- Integrity is prevention of unauthorized changes.



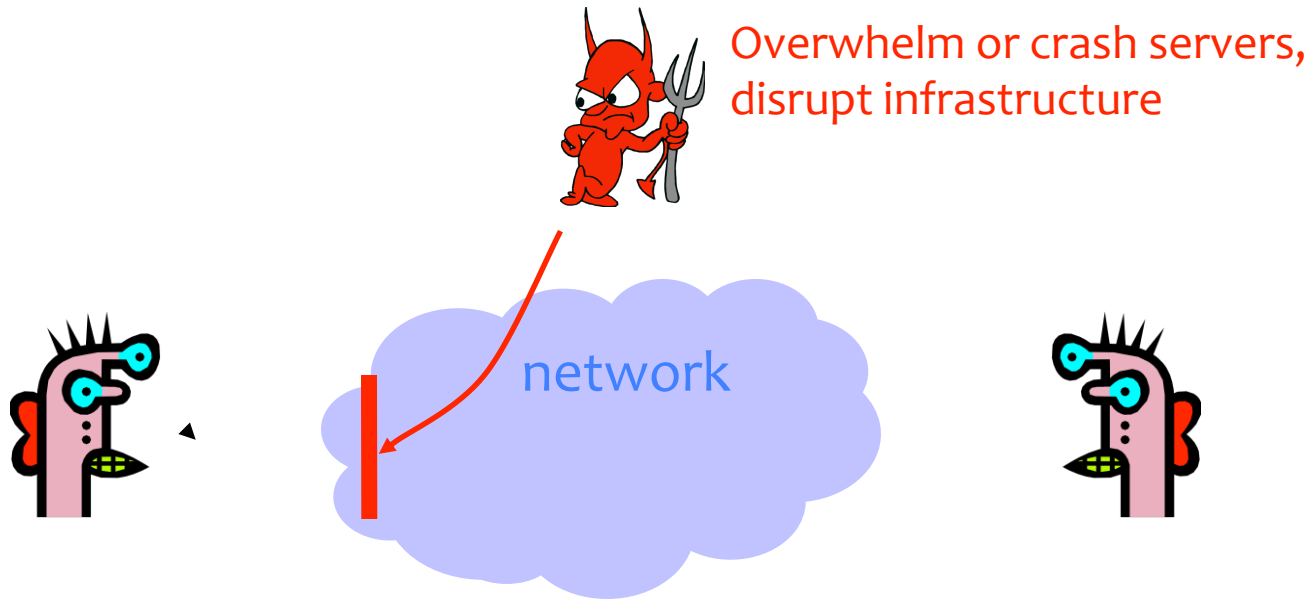
Authenticity

- Authenticity is **knowing who you're talking to**.



Availability

- Availability is ability to use information or resources.



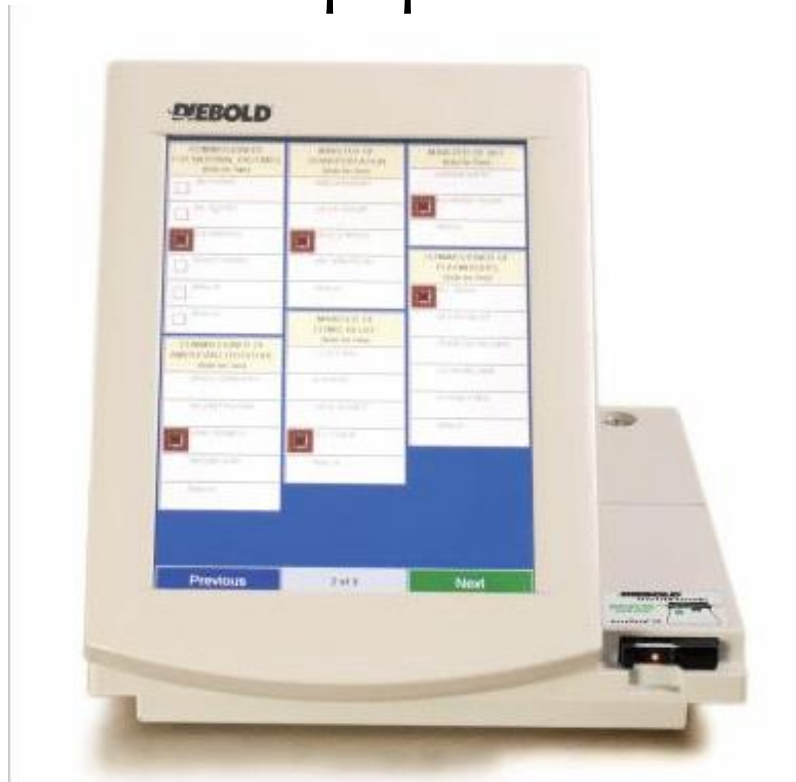
THREAT MODELING

Threat Modeling (Security Reviews)

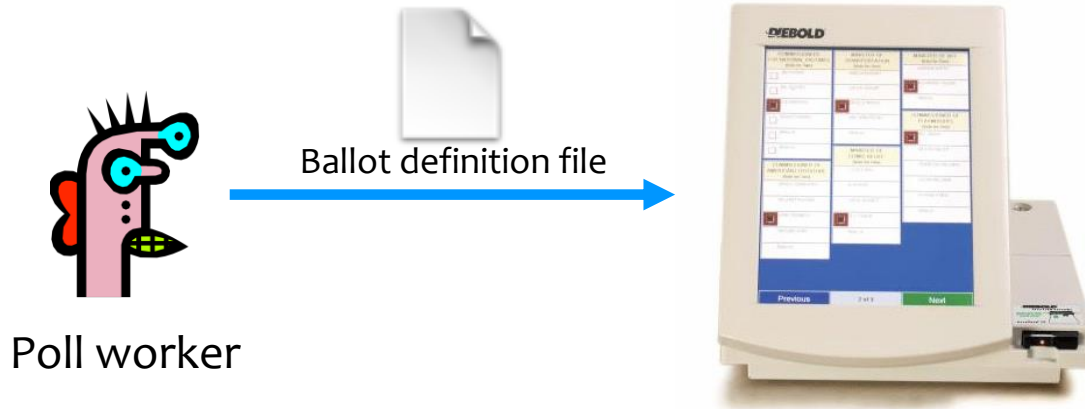
- **Assets:** What are we trying to protect? How valuable are those assets?
- **Adversaries:** Who might try to attack, and why?
- **Vulnerabilities:** How might the system be weak?
- **Threats:** What actions might an adversary take to exploit vulnerabilities?
- **Risk:** How important are assets? How likely is exploit?
- **Possible Defenses**

Example: Electronic Voting

- Popular replacement to traditional paper ballots

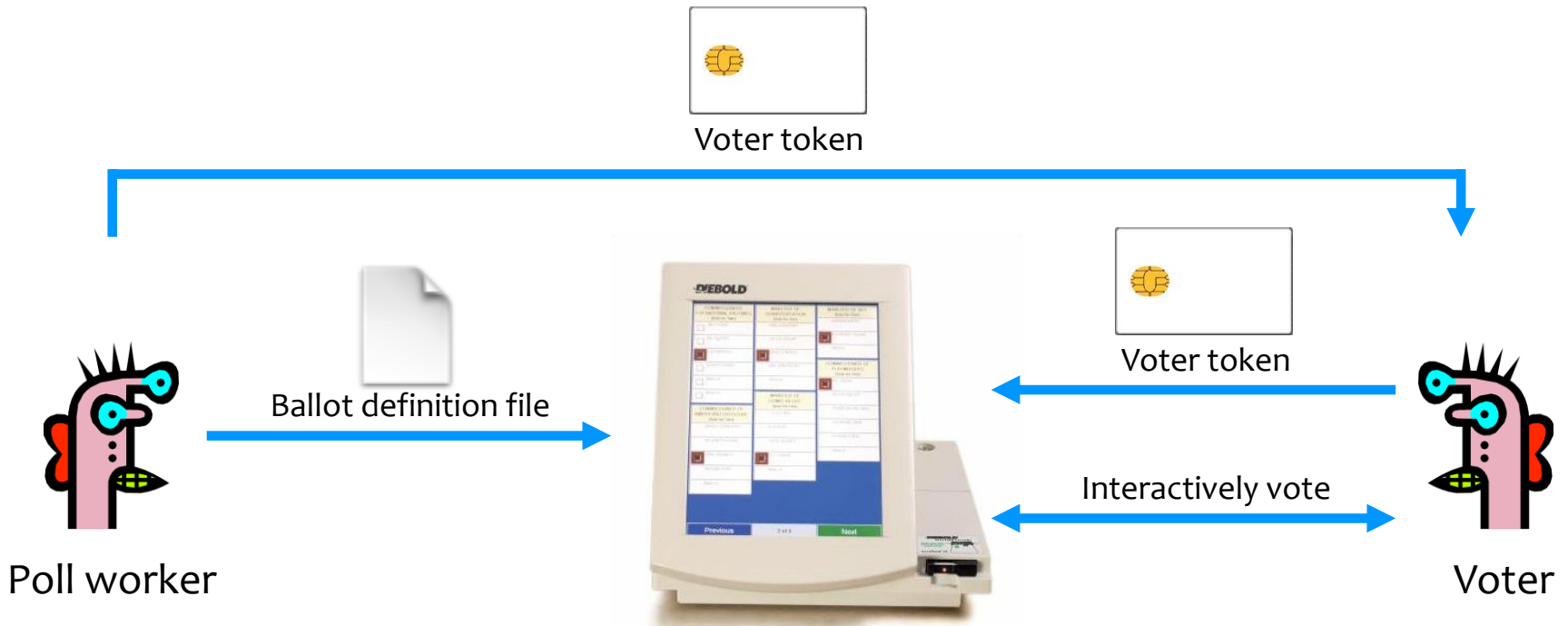


Pre-Election



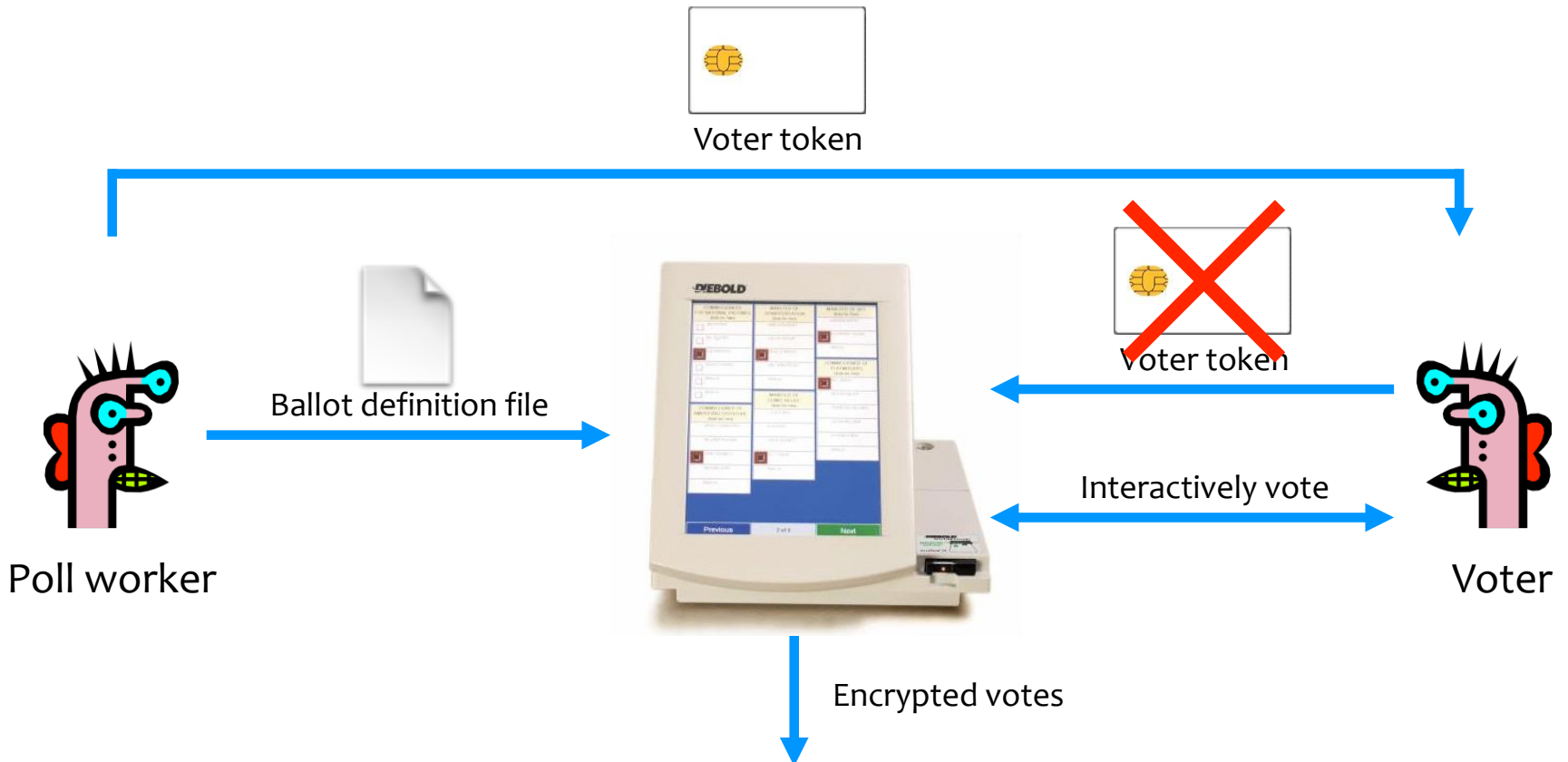
Pre-election: Poll workers load “ballot definition files” on voting machine.

Active Voting



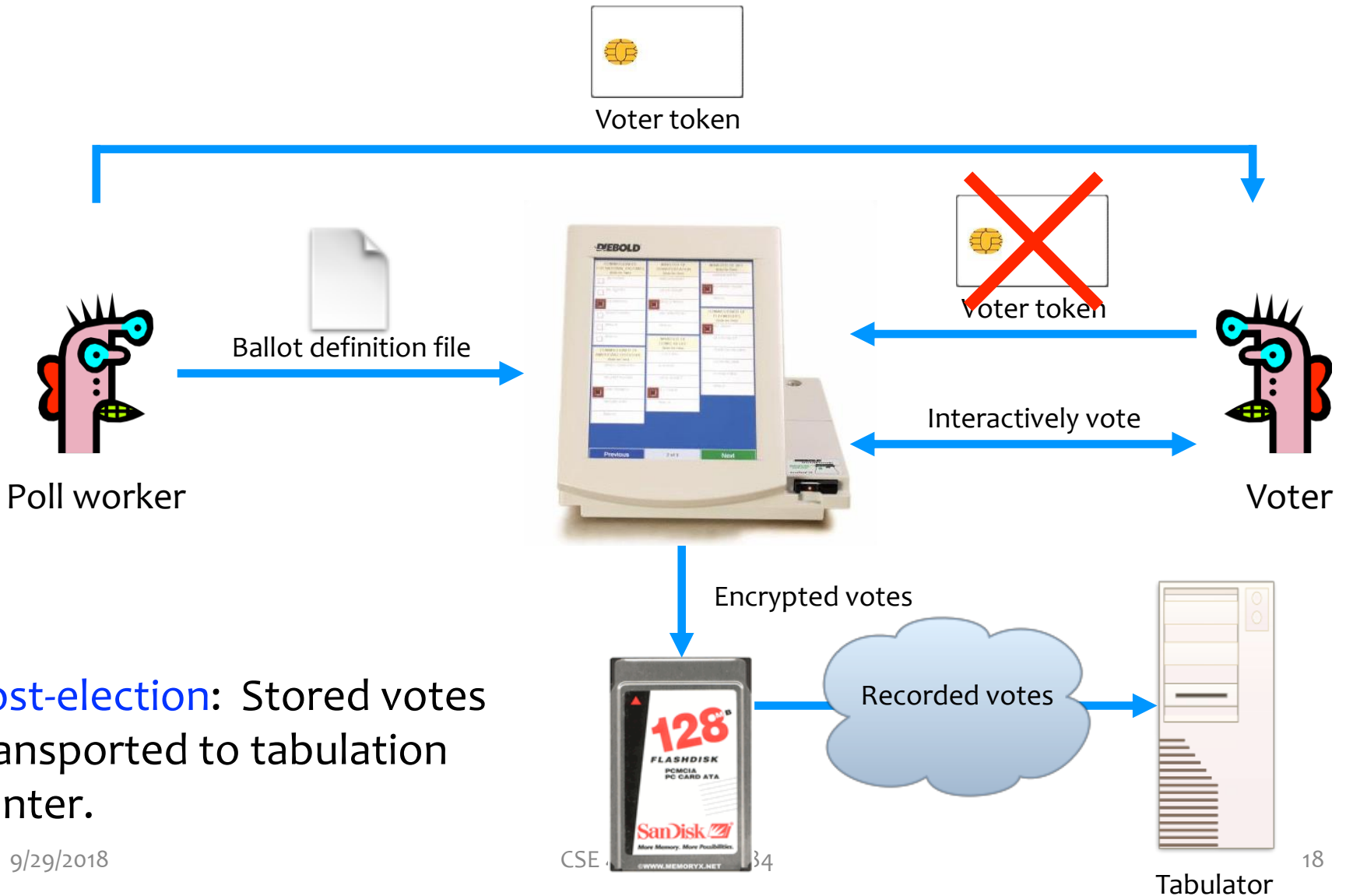
Active voting: Voters obtain **single-use** tokens from poll workers. Voters use tokens to **activate machines** and vote.

Active Voting



Active voting: Votes encrypted and stored. Voter token canceled.

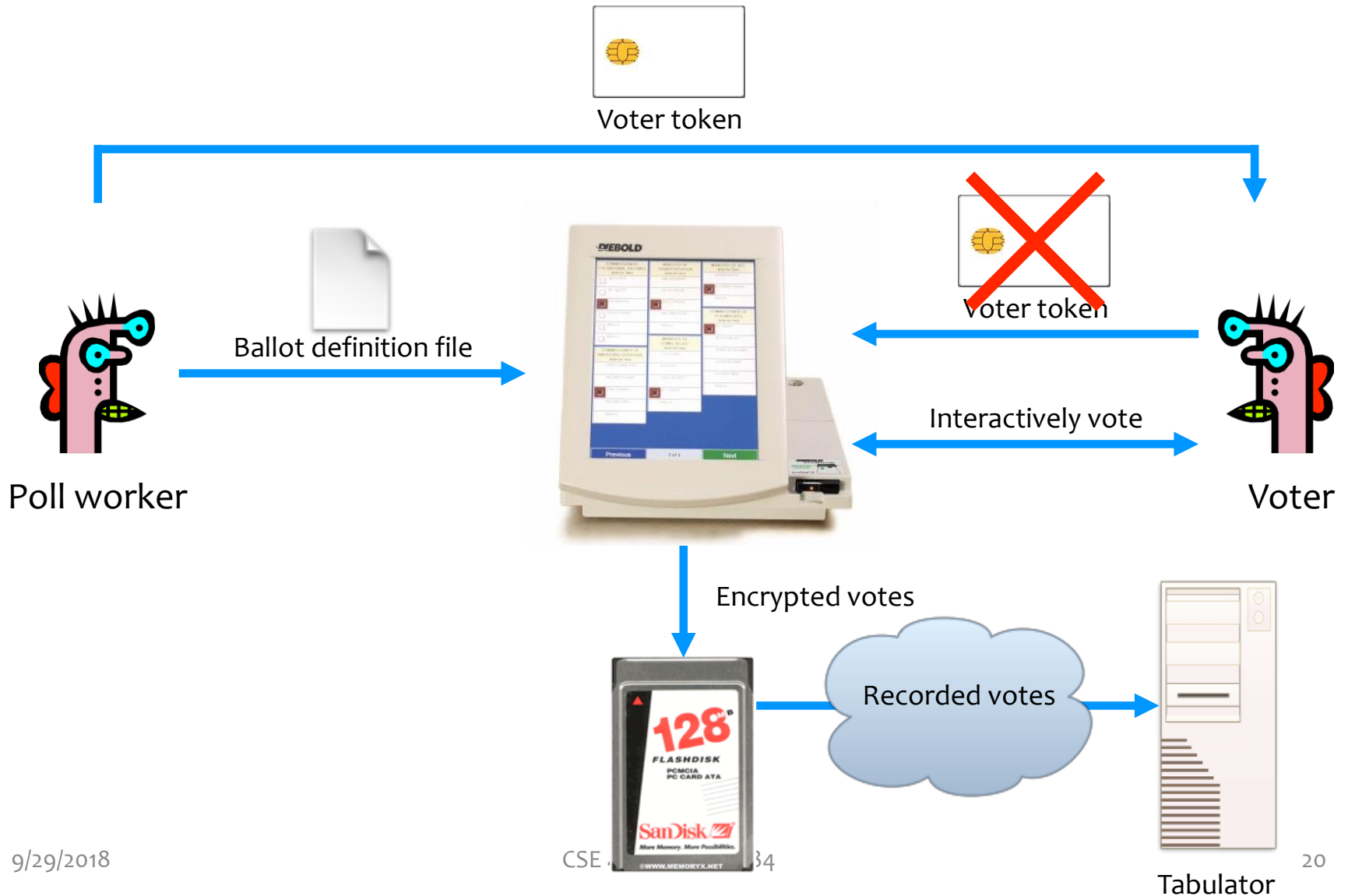
Post-Election



Security and E-Voting (Simplified)

- Functionality goals:
 - Easy to use, reduce mistakes/confusion
- Security goals:
 - Adversary should not be able to tamper with the election outcome
 - By changing votes (**integrity**)
 - By voting on behalf of someone (**authenticity**)
 - By denying voters the right to vote (**availability**)
 - Adversary should not be able to figure out how voters vote (**confidentiality**)

Can You Spot Any Potential Issues?



Potential Adversaries

- Voters
- Election officials
- Employees of voting machine manufacturer
 - Software/hardware engineers
 - Maintenance people
- Other engineers
 - Makers of hardware
 - Makers of underlying software or add-on components
 - Makers of compiler
- ...
- Or any combination of the above

What Software is Running?



Problem: An adversary (e.g., a poll worker, software developer, or company representative) able to control the software or the underlying hardware could do whatever he or she wanted.

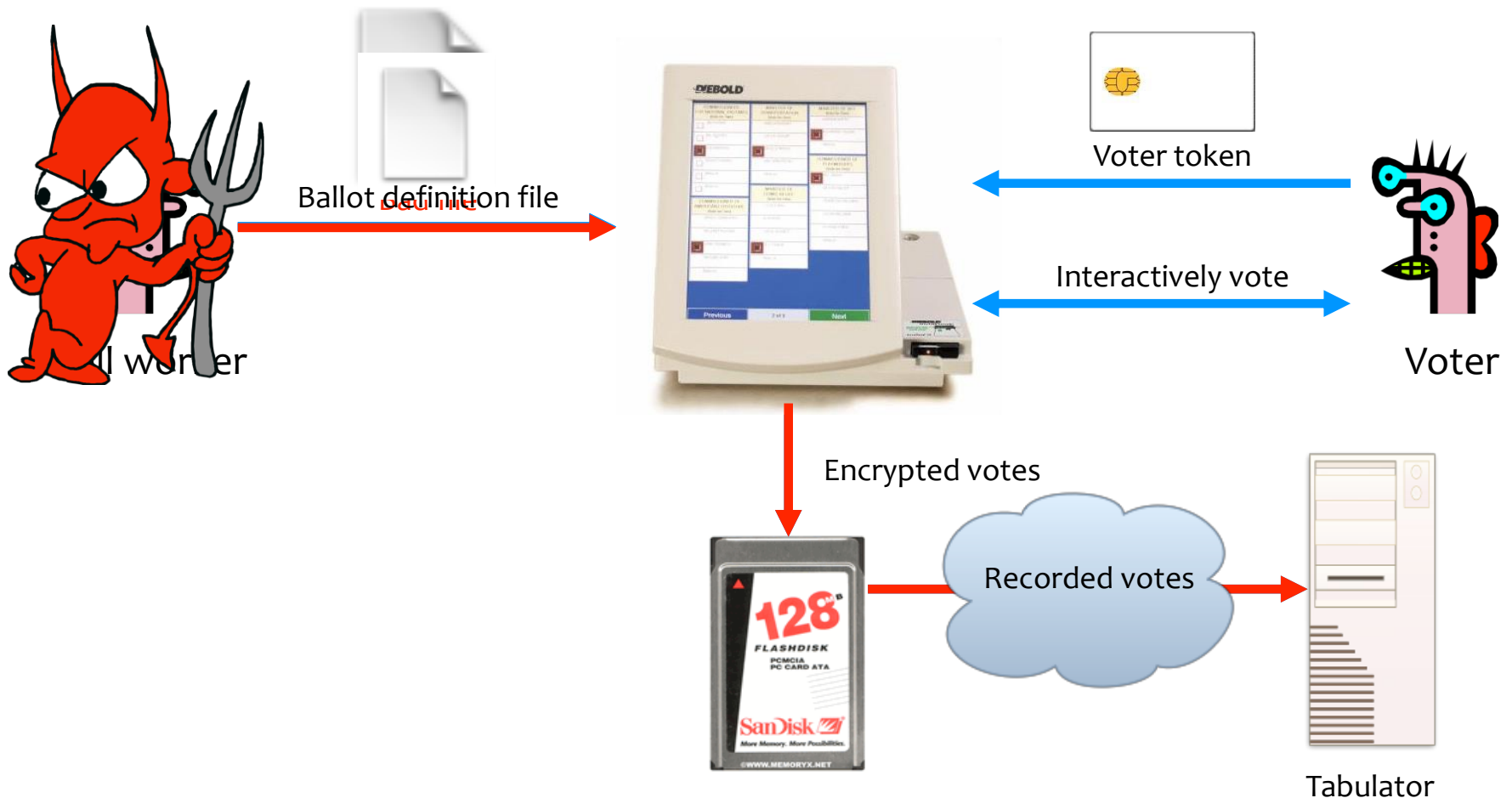


KEYS TO THE KINGDOM

Photo taken from Diebold's online store. The keys that open every Diebold touch-screen voting machine. Working copies have been made from the photo.

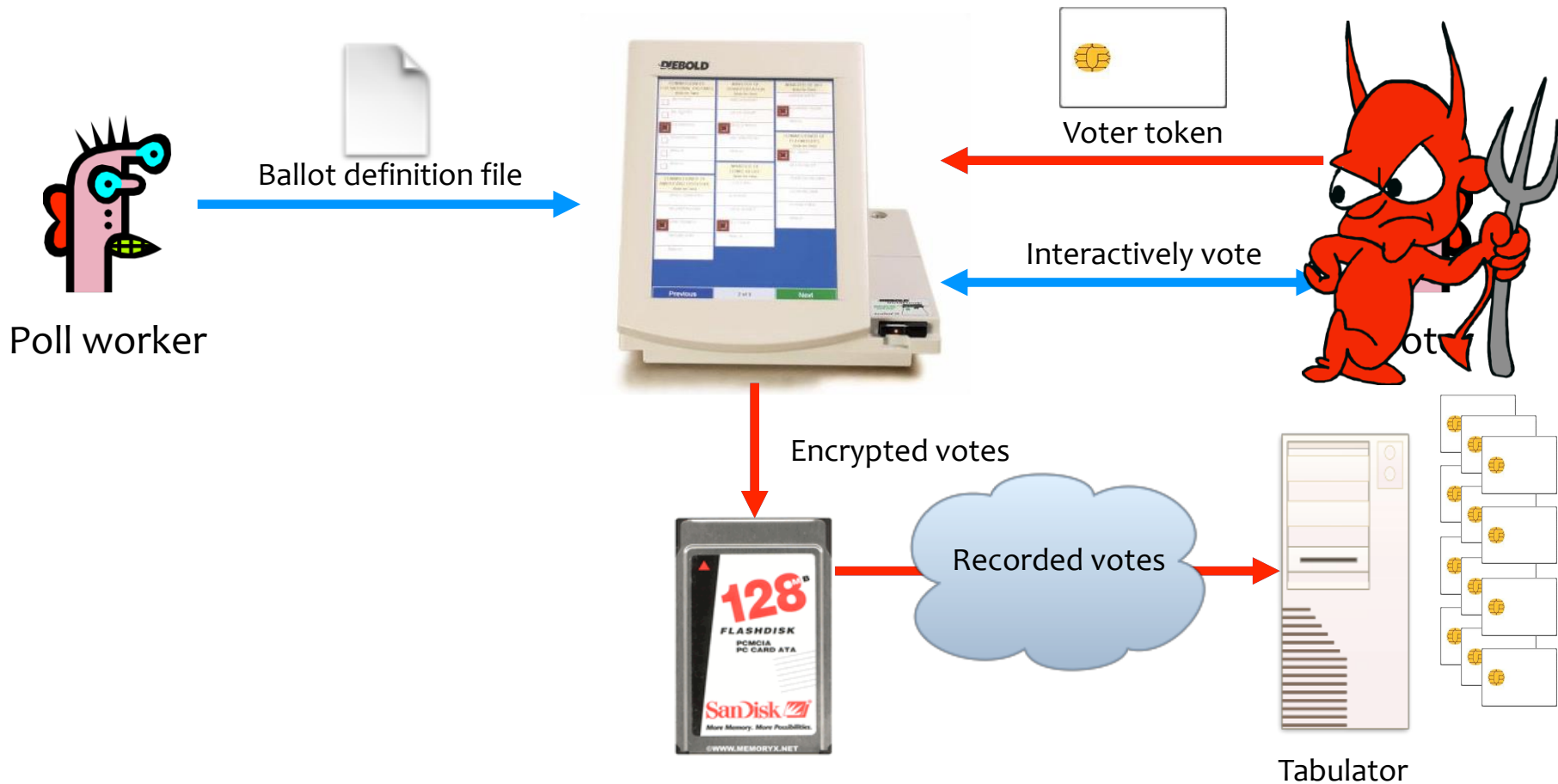
Problem: Ballot definition files are not authenticated.

Example attack: A malicious poll worker could modify ballot definition files so that votes cast for “Mickey Mouse” are recorded for “Donald Duck.”



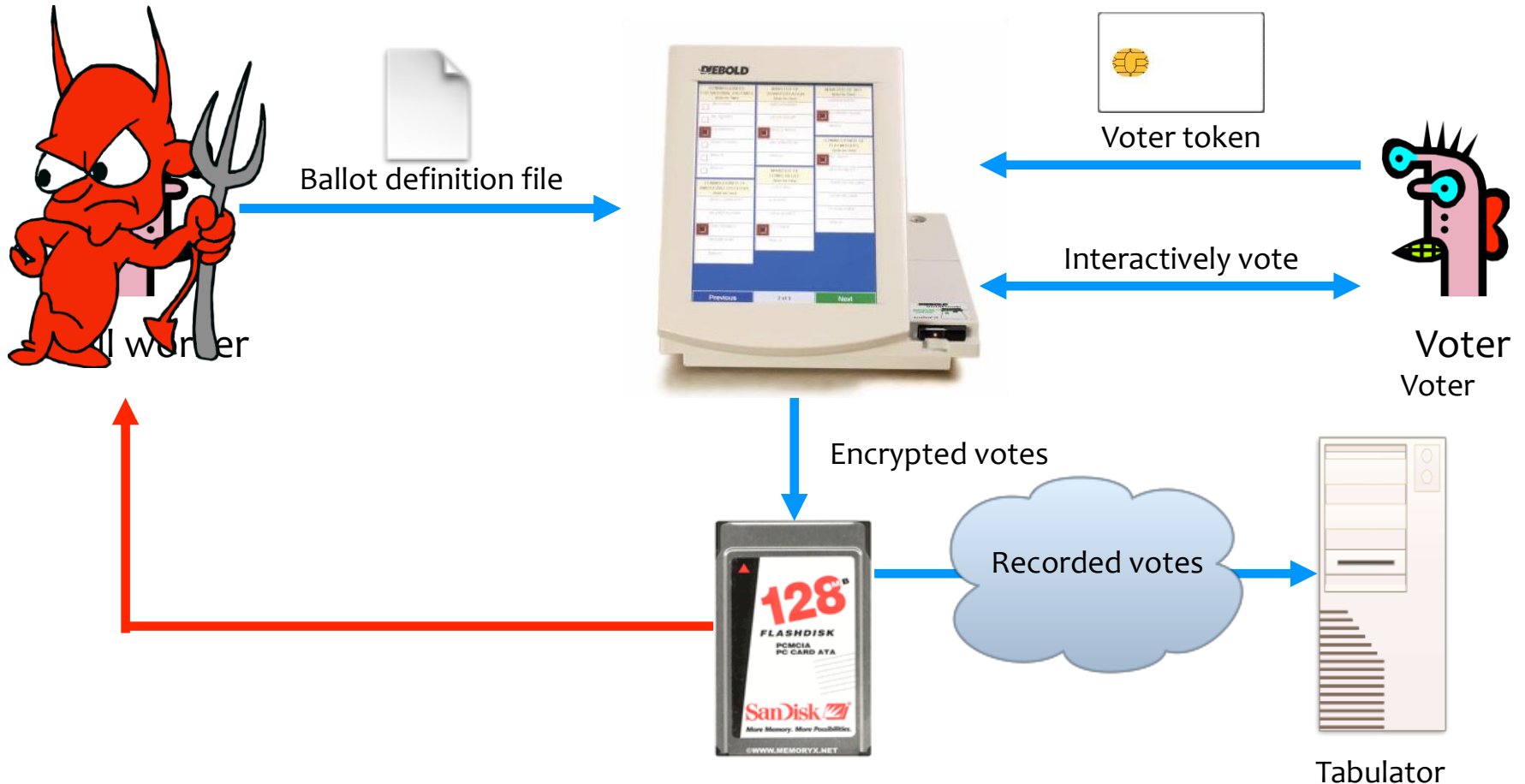
Problem: Smartcards can perform cryptographic operations. But there is **no authentication from voter token to terminal**.

Example attack: A regular voter could make his or her own voter token and **vote multiple times**.



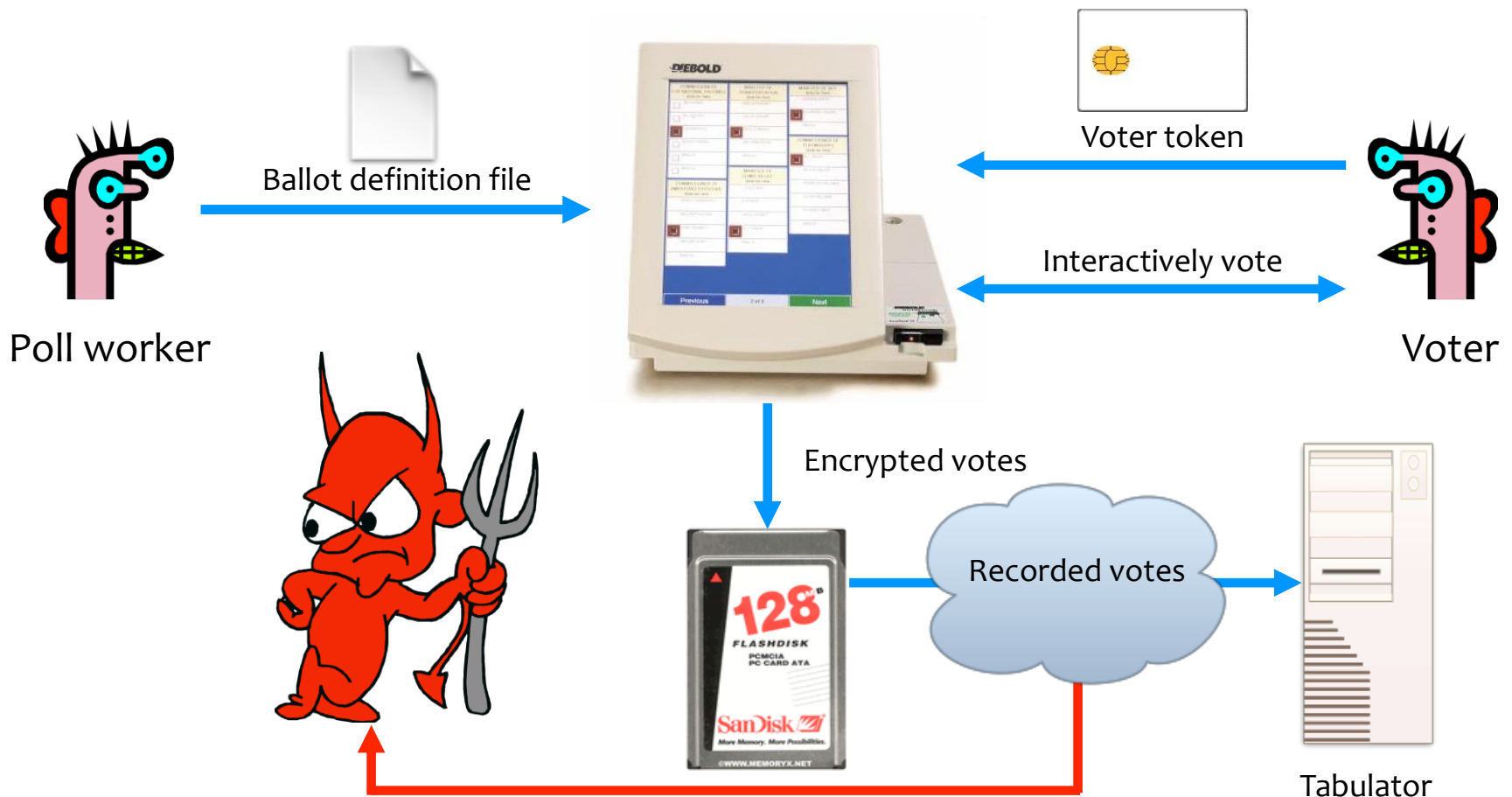
Problem: Encryption key (“F2654hD4”) **hard-coded** into the software since (at least) 1998. Votes stored in the order cast.

Example attack: A poll worker could **determine how voters vote.**



Problem: When votes transmitted to tabulator over the Internet or a dialup connection, they are **decrypted first**; the cleartext results are sent the the tabulator.

Example attack: A sophisticated outsider could determine how voters vote.



Tables Often Help!



Example Table 1

| Attacker “Positions” | Machine Manufacturer | Poll Worker | Voter | Power Company Employee |
|-----------------------------------|-------------------------|-------------|-------|------------------------------|
| Voter Privacy | | | | |
| Vote Integrity | | | | |
| Voting Machine Availability | | | | |
| ... | | | | |

- What can different parties do? Each cell would have an action or actions that these parties might try to do
- Note that some parties could collaborate

Example Table 2

| Attack Methods | Modify Software | Produce Fake Voter Tokens | Steal Flash Drive | Intercept Network Connections |
|-----------------------------|-----------------|---------------------------|-------------------|-------------------------------|
| Voter Privacy | | | | |
| Vote Integrity | | | | |
| Voting Machine Availability | | | | |
| ... | | | | |

- What different attack methods are there? (Columns)
- Who could mount these different attacks? What are the attack details (the cells)
- How easy is it to implement each of these attack methods?

Table from Paper

<https://homes.cs.washington.edu/~yoshi/papers/eVoting/vote.pdf>

| | Voter (with forged smartcard) | Poll Worker (with access to storage media) | Poll Worker (with access to network traffic) | Internet Provider (with access to network traffic) | OS Developer | Voting Device Developer | Section |
|---|-------------------------------------|--|--|--|-----------------|-------------------------------|----------|
| Vote multiple times using forged smartcard | • | • | • | | | | 3.2 |
| Access administrative functions or close polling station | • | • | | | • | • | 3.3 |
| Modify system configuration | | • | | | • | • | 4.1 |
| Modify ballot definition (e.g., party affiliation) | | • | • | • | • | • | 4.2 |
| Cause votes to be miscounted by tampering with configuration | | • | • | • | • | • | 4.2 |
| Impersonate legitimate voting machine to tallying authority | | • | • | • | • | • | 4.3 |
| Create, delete, and modify votes | | • | • | • | • | • | 4.3, 4.5 |
| Link voters with their votes | | • | • | • | • | • | 4.5 |
| Tamper with audit logs | | • | | | • | • | 4.6 |
| Delay the start of an election | | • | • | • | • | • | 4.7 |
| Insert backdoors into code | | | | | • | • | 5.3 |

Table 1: This table summarizes some of the more important attacks on the system.

TOWARDS DEFENSES

Approaches to Security

- Prevention
 - Stop an attack
- Detection
 - Detect an ongoing or past attack
- Response
 - Respond to attacks
- The threat of a response may be enough to deter some attackers

Example security mechanisms

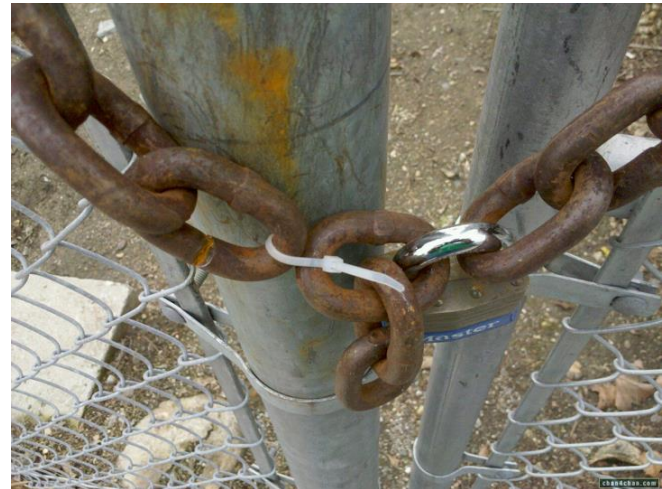
- Verifying the identity of a prospective user by demanding a password
 - Authentication
- Shielding the computer to prevent interception and subsequent interpretation of electromagnetic radiation
 - Covert channels
- Enciphering information sent via communication channels
 - Cryptography
- Locking the room containing the computer
 - Physical aspects of security
- Controlling who is allowed to make changes to a computer system
 - Social aspects of security

Whole System is Critical

- Securing a system involves a **whole-system view**
 - Cryptography
 - Implementation
 - People
 - Physical security
 - Everything in between
- This is because “security is only as strong as the weakest link,” and security can fail in many places
 - No reason to attack the strongest part of a system if you can walk right around it.

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Whole System is Critical



In reality

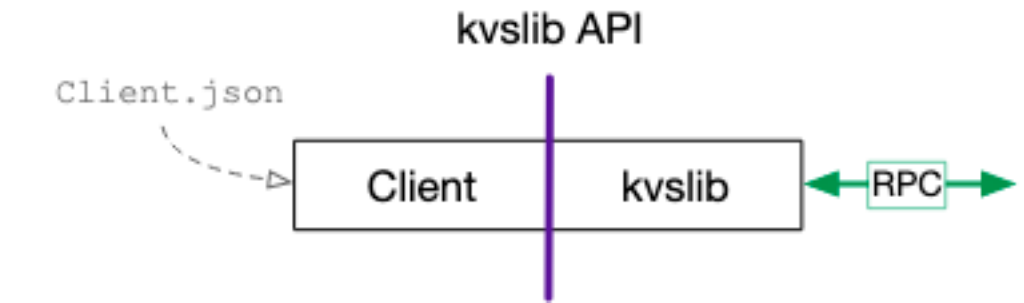


Thank you for taking CPSC 416!

And attending 8am classes

- **Upcoming events:**

- ▶ A6 due on Friday 11:59p
- ▶ Final exam on Sunday @ 3:30p



```

func clientCore chan int {
    tracer := tracing.NewTracerFromFile("client_core.go.json")
    defer tracer.Close()

    ttrace := tracer.CreateTrace()

    client, err := rpc.Dial("tcp", serverPort)
    if err != nil {
        log.Fatalf("dialing: %v", err)
    }
    ttrace.RecordAction(ClientStart(serverPort: serverPort))

    args := Args{Token: trace.GenerateToken()}
    var reply *Reply
    err = client.Call("Person.GetName", args, &reply)
    if err != nil {
        log.Fatalf("person error: %v", err)
    }
    fmt.Printf("GetName: %v\n", reply.Name)
    ttrace.RecordAction(ReplyToken)

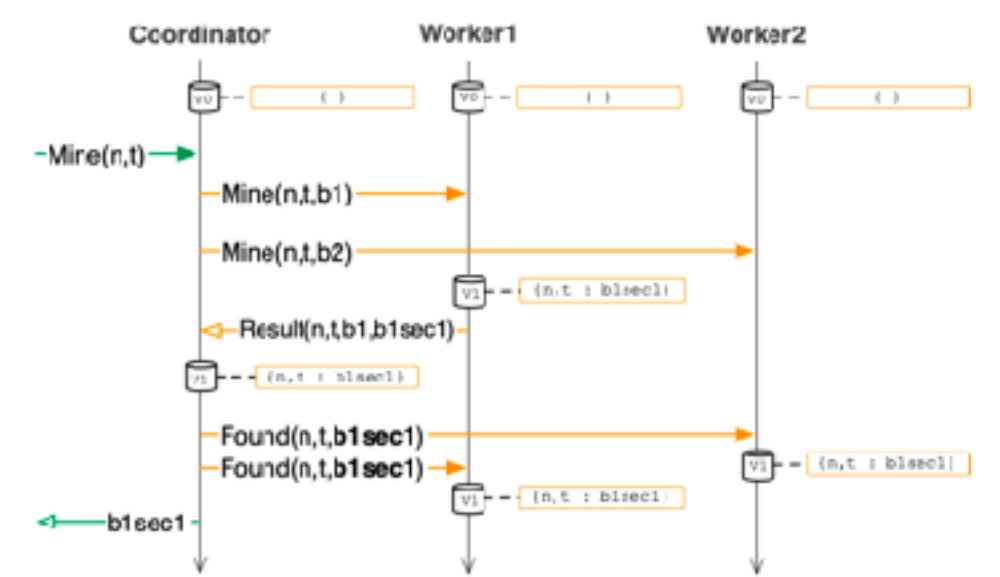
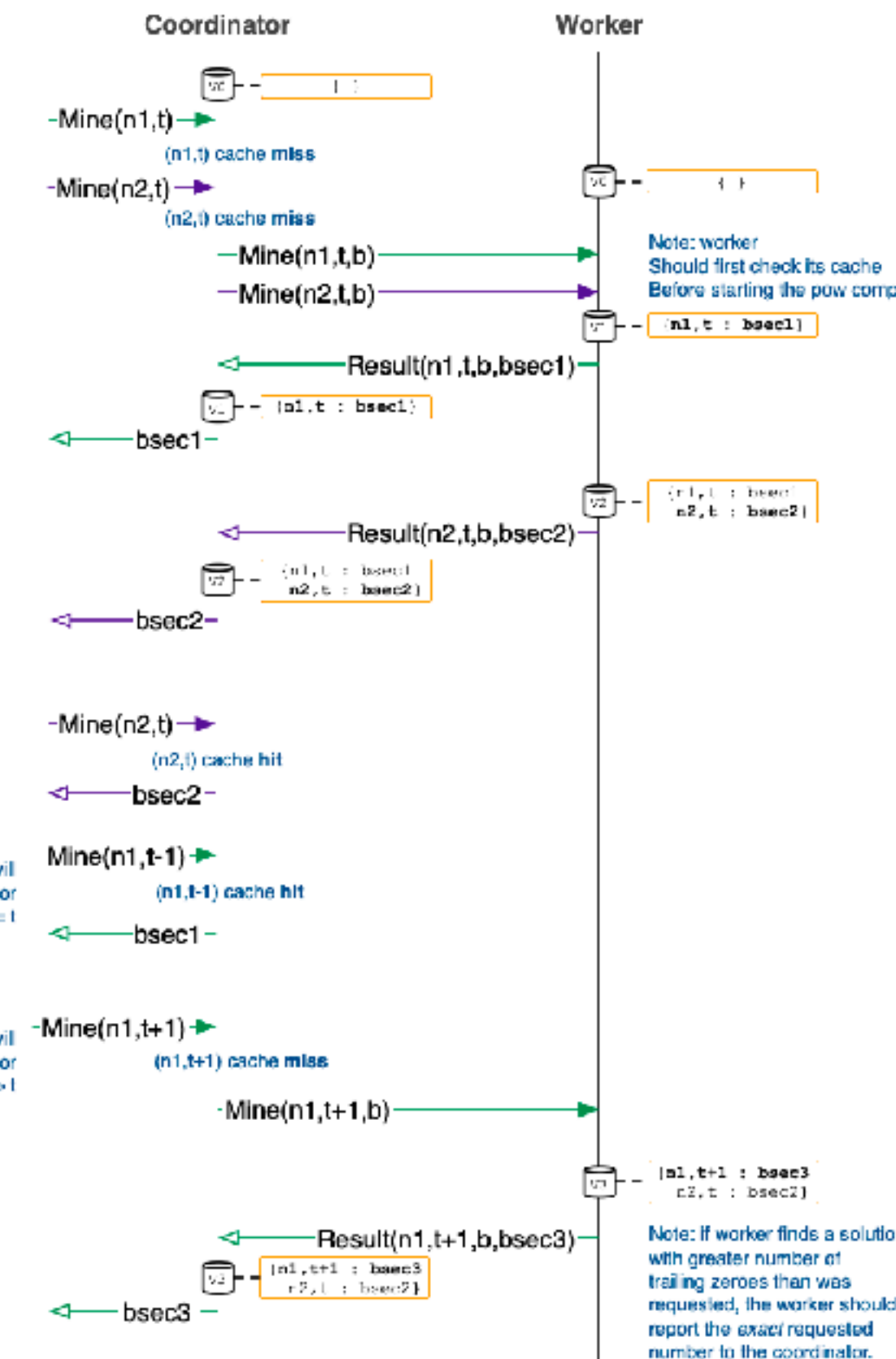
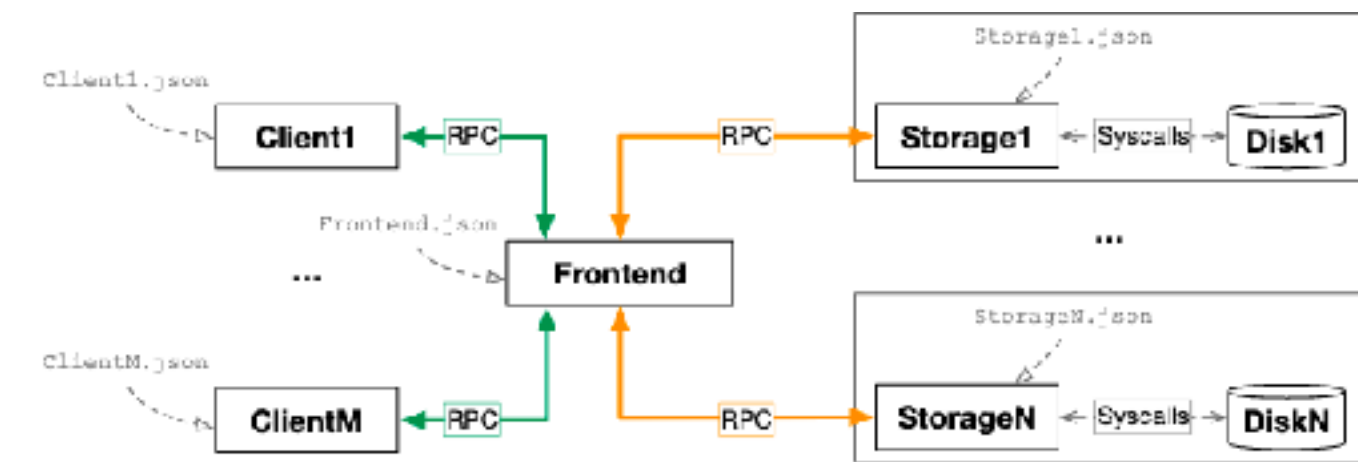
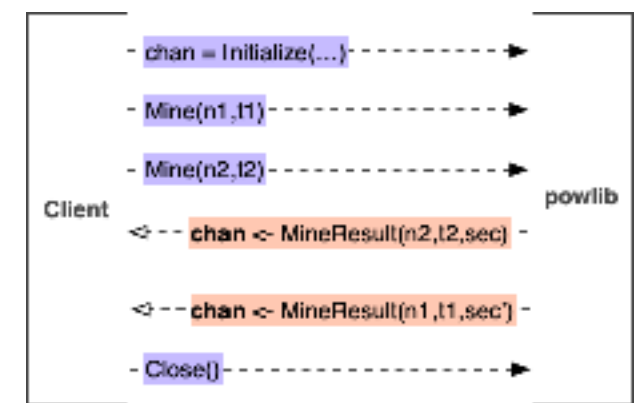
    ttrace.RecordAction(ClientFinish(serverPort: serverPort))
    core <- 1
}
    
```

```

type Args struct {
    Token tracing.TracingToken
}

type Reply struct {
    Name string
    Token tracing.TracingToken
}

func (p *Person) GetName(args Args, reply *Reply) error {
    trace := p.tracer.ReceiveToken(args.Token)
    reply.Name = p.name
    reply.Token = trace.GenerateToken()
    return nil
}
    
```



Note: an (n,t) cached entry will generate a CacheHit(n,k) for any request (n,k) with k <= t

Note: an (n,t) cache entry will generate a CacheMiss(n,k) for any (n,k) with k > t

You took a challenging course
(during a pandemic)
I hope it was a rewarding
experience. You should be proud!