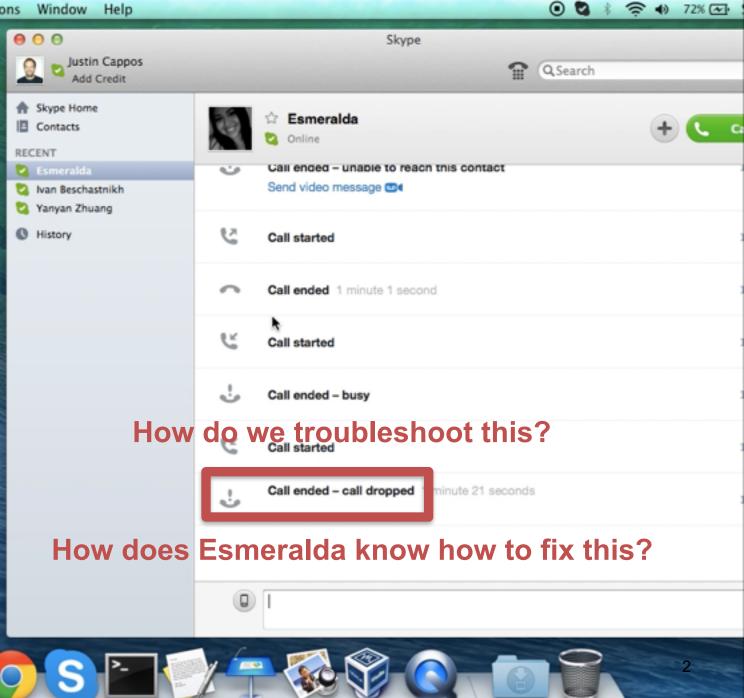
NetCheck: Network Diagnoses from Blackbox Traces

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Goal Find bugs in networked applications Large complex unknown applications

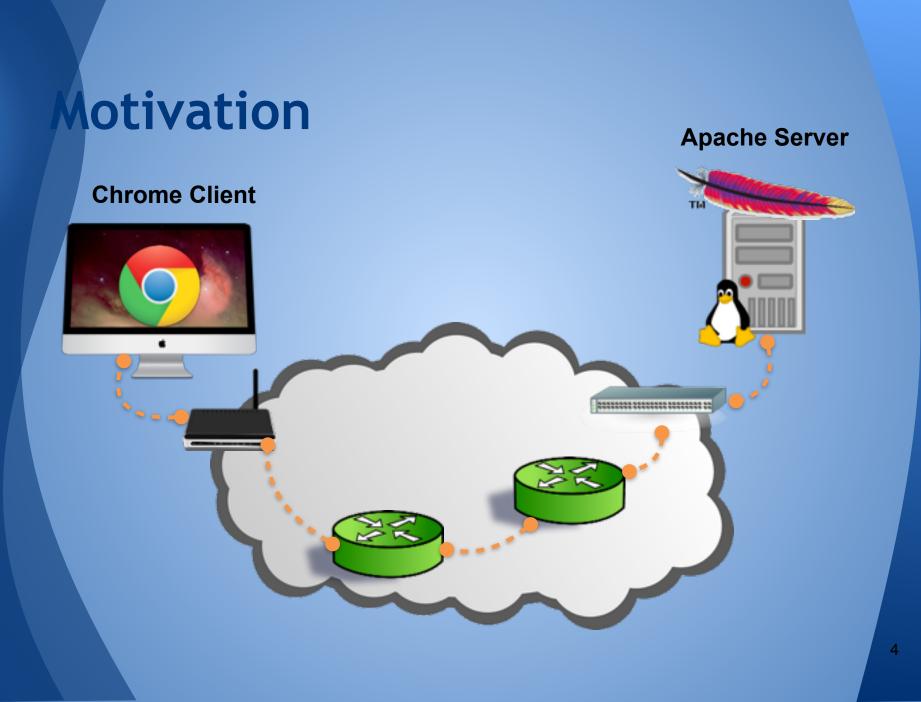






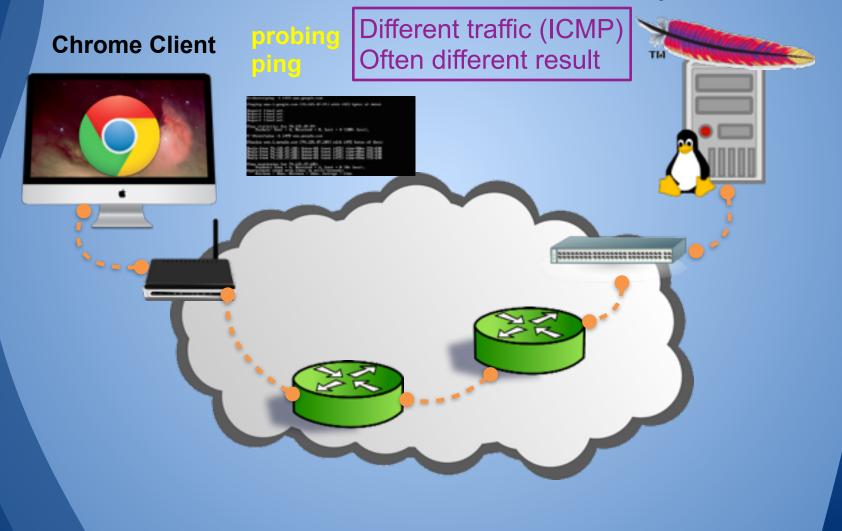


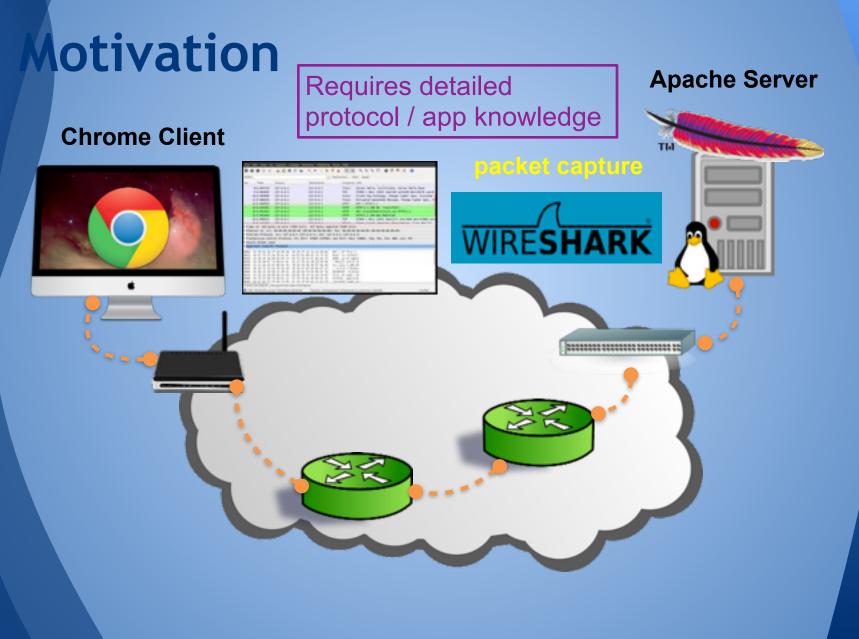
Understandable output / fix

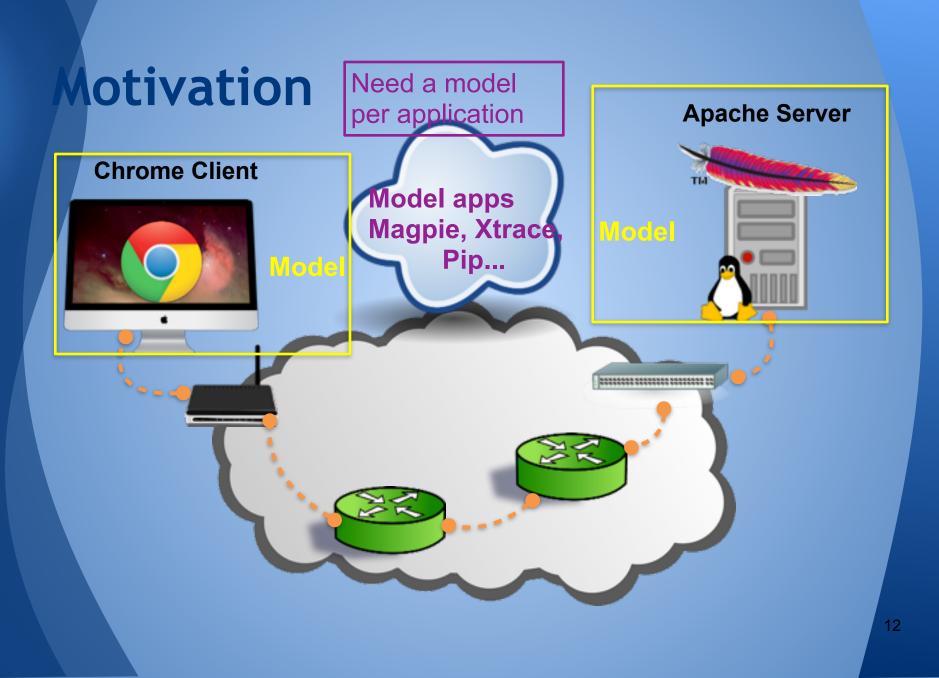


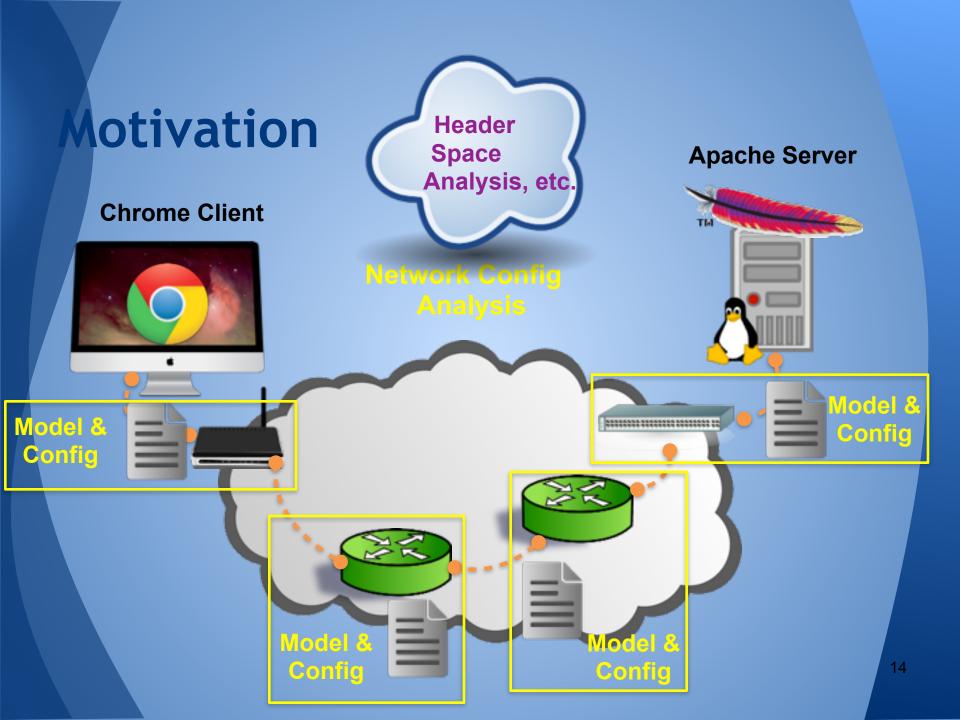
otivation

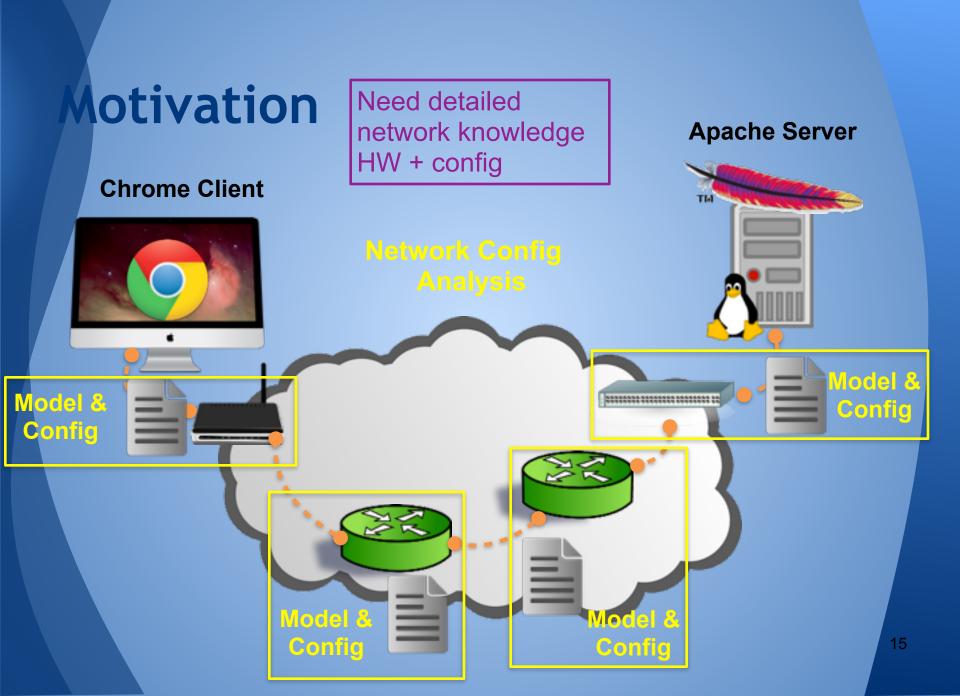
Apache Server

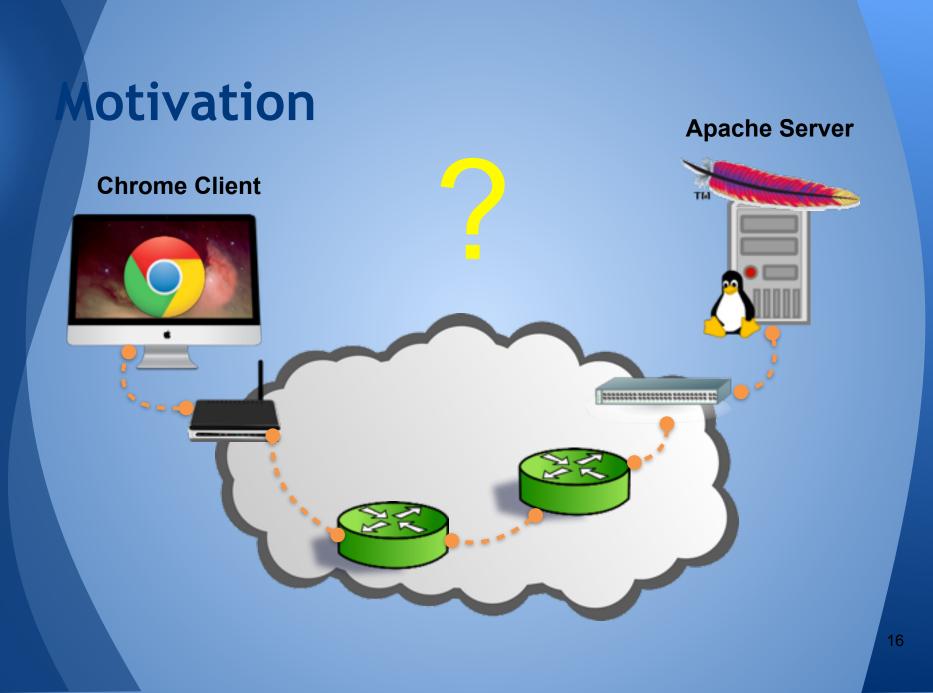


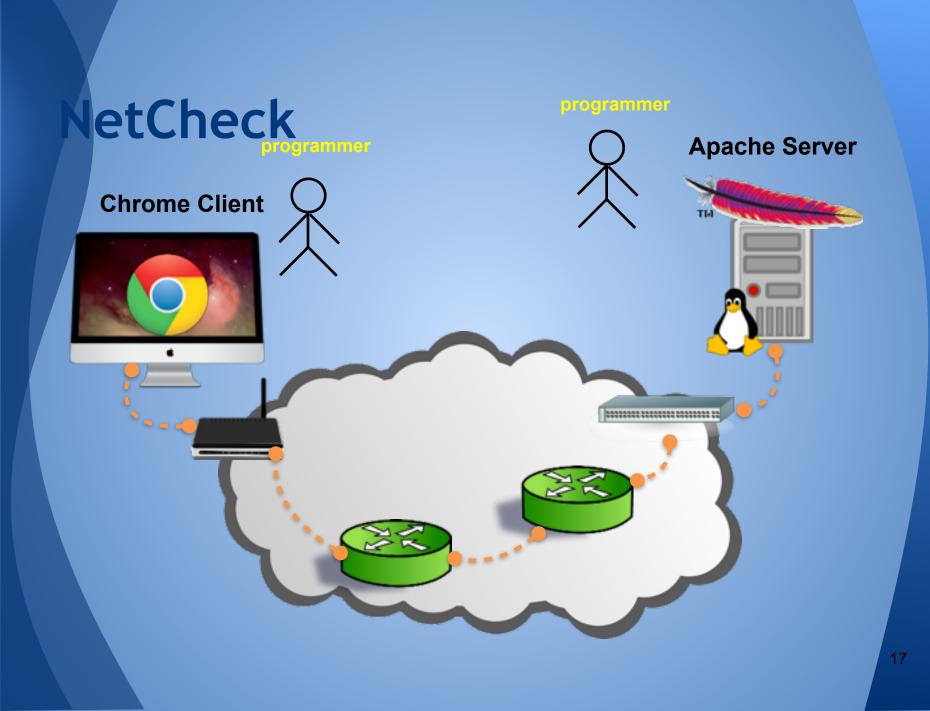


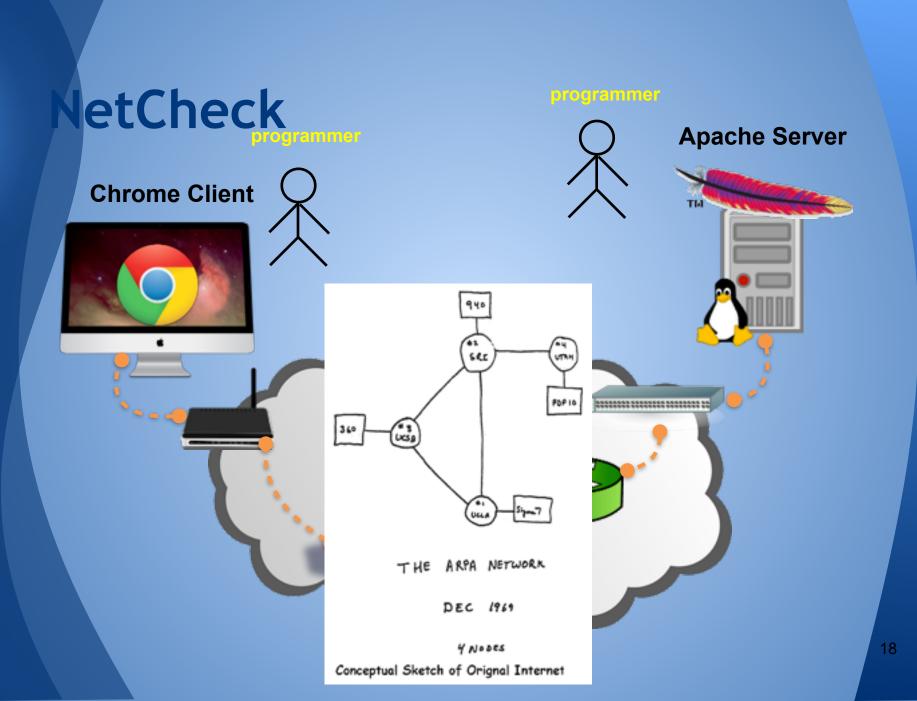


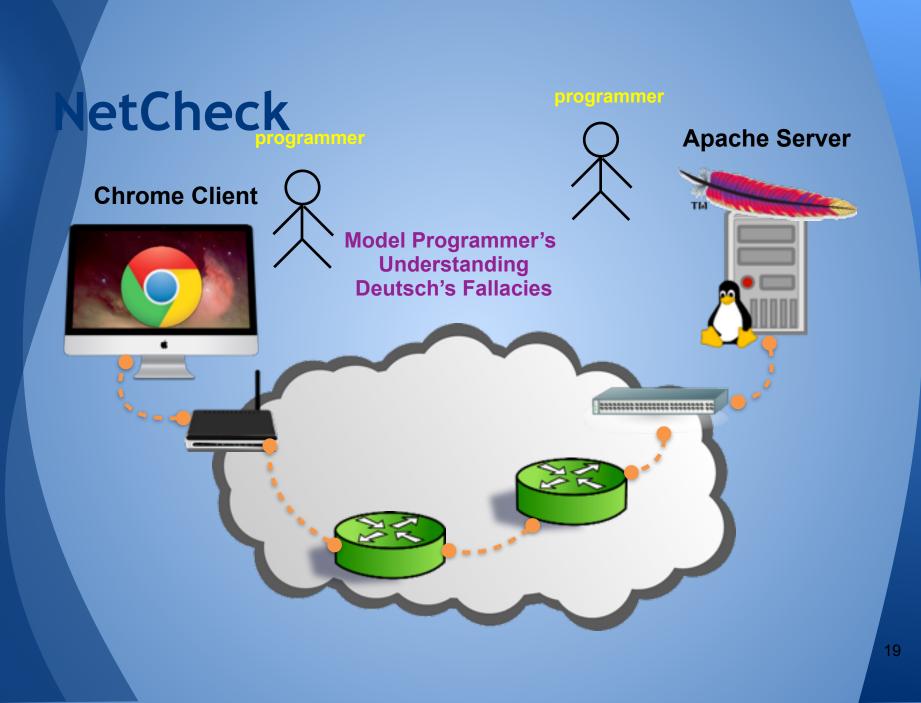






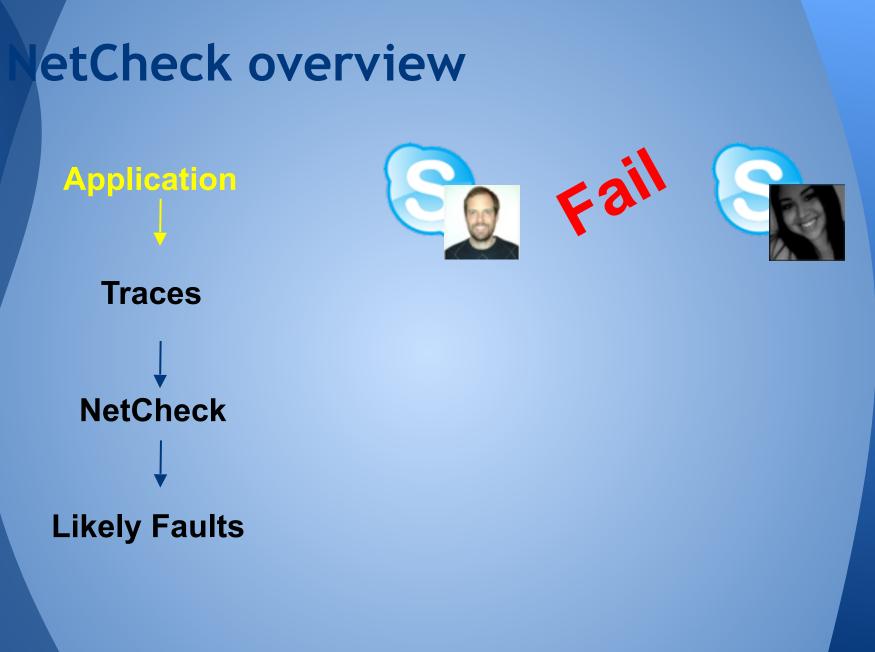


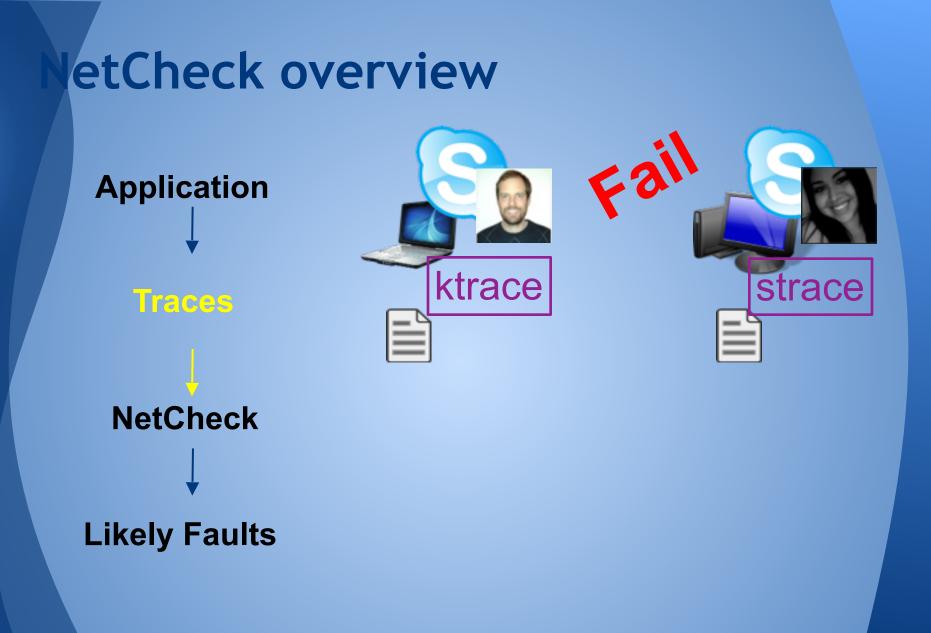


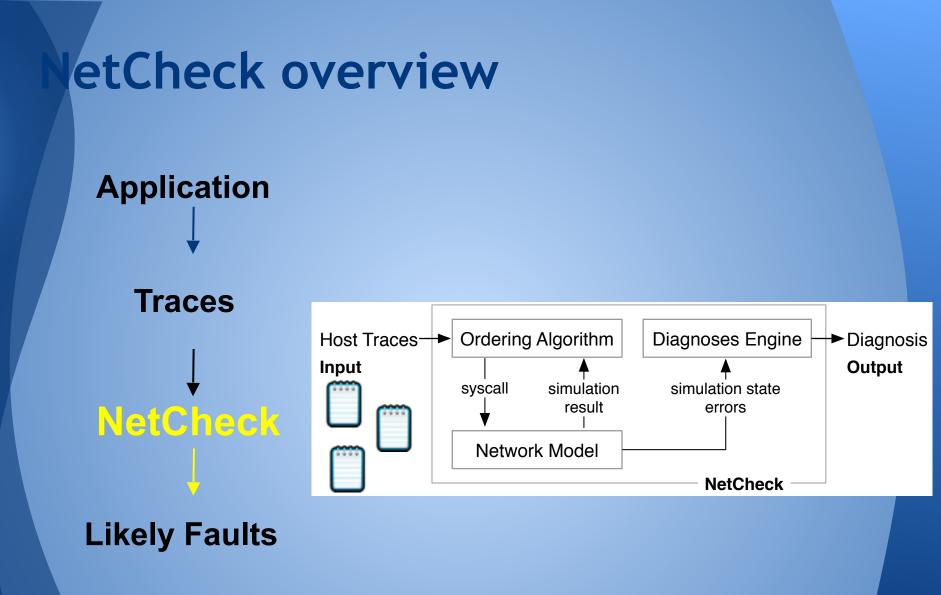


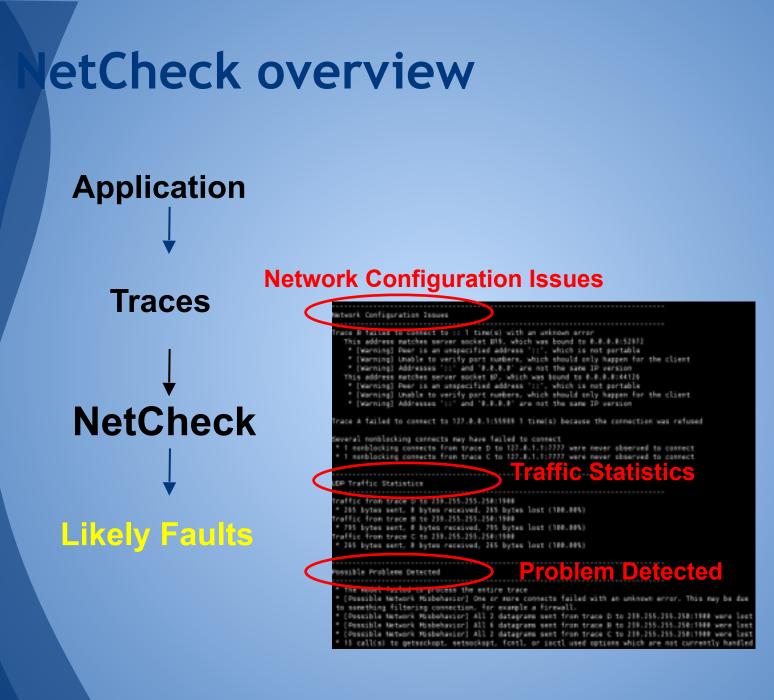
Outline

Motivation
NetCheck Overview
Trace Ordering
Network Model
Fault Classification
Results / Conclusion









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Trace

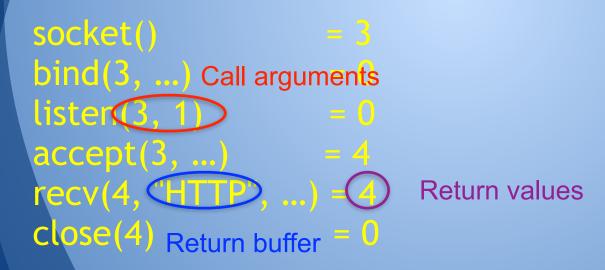
Ordering

(a)

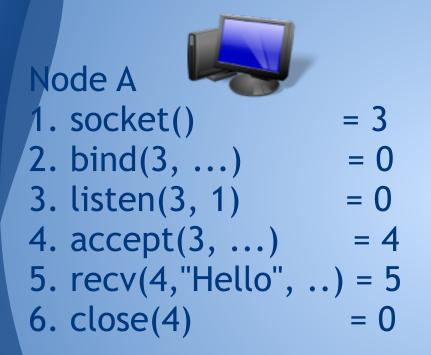
Traces

Traces

Series of locally ordered system calls Don't want to modify apps or use a global clock Gathered by strace, ktrace, systrace, truss, etc. Call arguments and "return values"



What we see is this:



Node B 1. socket() = 3 2. connect(3,...) = 0 3. send(3, "Hello",.) = 5 4. close(3) = 0

- one trace per host
- local order but no global order
- Q: how do we reconstruct what really happened?

What we want is this

A1. socket() = 3 B1. socket() A2. bind(3,) A3. listen(3, 1) B2. connect(3,...) = 0 A4. accept(3, ...)B3. send(3, "Hello", ...) = 5 A5. recv(4, "Hello", ...) = 5 B4. close(3)= 0 A6. close(4)

B

-The ground truth

What we want is this

A1. socket() = 3 B1. socket() A2. bind(3,) A3. listen(3, 1) B2. connect(3,...) = 0 A4. accept(3, ...) B3. send(3, "Hello", ...) = 5 A5. recv(4, "Hello", ...) = 5 B4. close(3)= 0 A6. close(4)

The ground truth



Goal: find <u>an</u> equivalent interleaving

Observation 1: Order Equivalence

Node A 1. socket() = 3 2. bind(3, ...) = 0 3. listen(3, 1) = 0 4. accept(3, ...) = 4 5. recv(4, "Hello", ...) = 5 6. close(4) = 0 Node B 1. socket() = 3 2. connect(3,...) = 0 3. send(3, "Hello",.) = 5 4. close(3) = 0

one trace per host

 - local order but no global order
 Q: how do we reconstruct what really happened?
 The socket() calls are not visible to the other side Some orders are equivalent! 30

Observation 2: Return Values Guide Ordering

 Node A

 1. socket()
 = 3

 2. bind(3, ...)
 = 0

 3. listen(3, 1)
 = 0

 4. accept(3, ...)
 = 4

 5. recv(4, "Hello", ...)
 = 5

 6. close(4)
 = 0

Node B 1. socket() = 3 2. connect(3,...) = 0 3. send(3, "Hello",.) = 5 4. close(3) = 0

- one trace per host

- local order but no global order

Q: how do we reconstruct what really happened?

Return values guide ordering

A2. bind(3, ...) = 0A3. listen(3, 1) = 0B2. connect(3, ...) = 0

One valid ordering: all syscalls returned successfully.

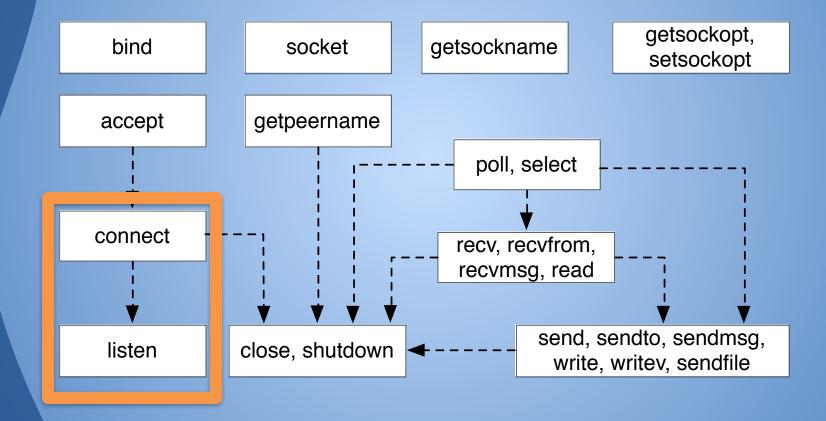
A2. bind(3, ...) = 0B2. connect(3, ...) = -1, ECONNREFUSED A3. listen(3, 1) = 0A3. listen(3, 1) = 0

A second valid ordering connect failed with ECONNREFUSED.

A call's return value may-depend-on a remote call's action

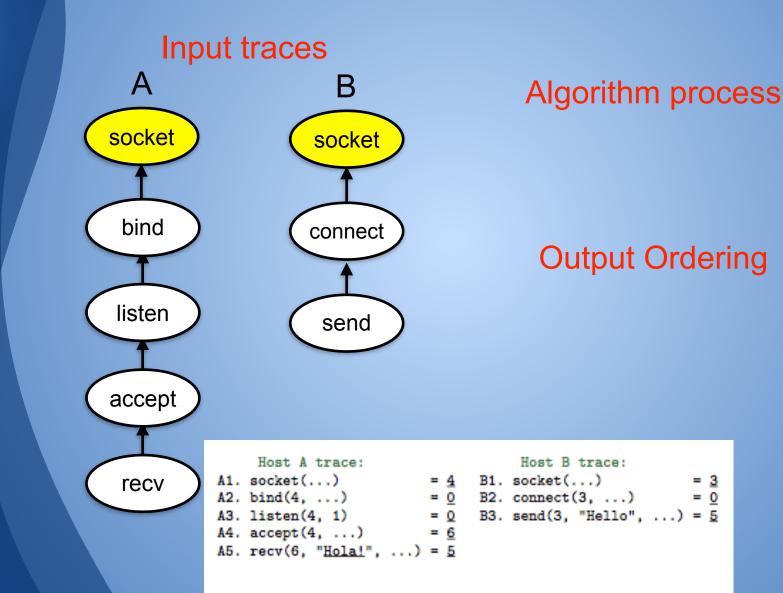
Result indicates order of calls

Deciding call order

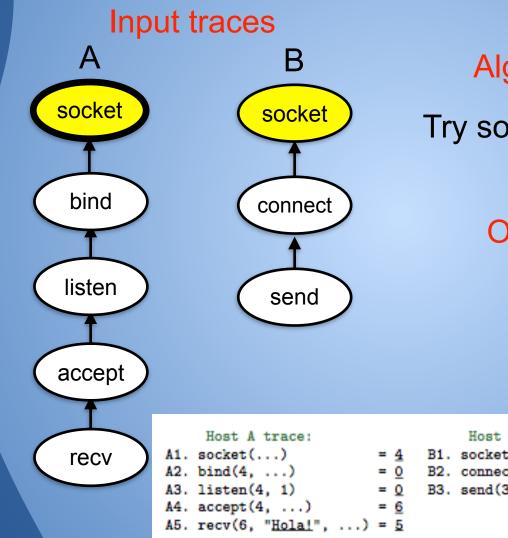


full set of may-depend-on relations

Prdering Algorithm



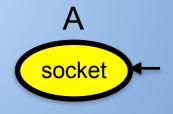
Prdering Algorithm



Algorithm process

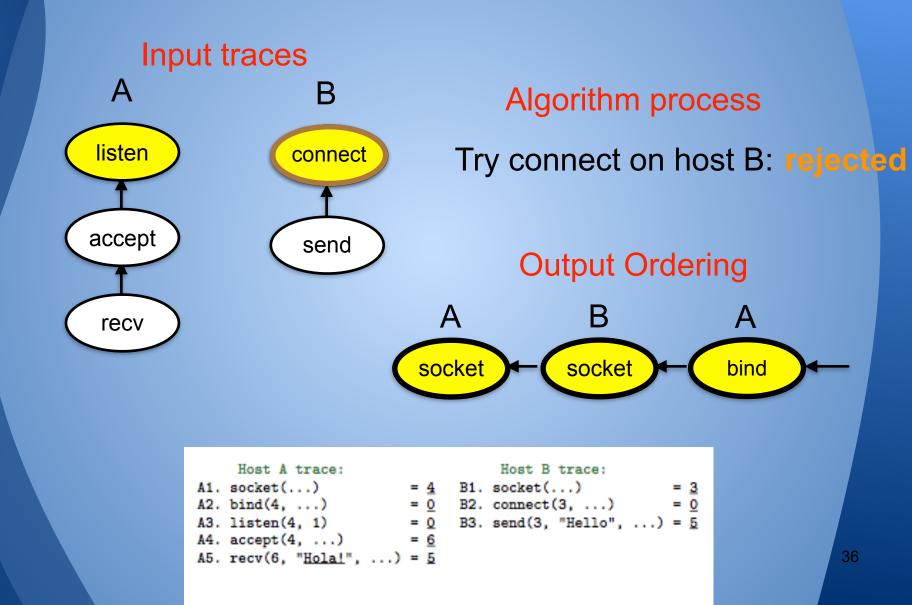
Try socket on host A: accepted

Output Ordering

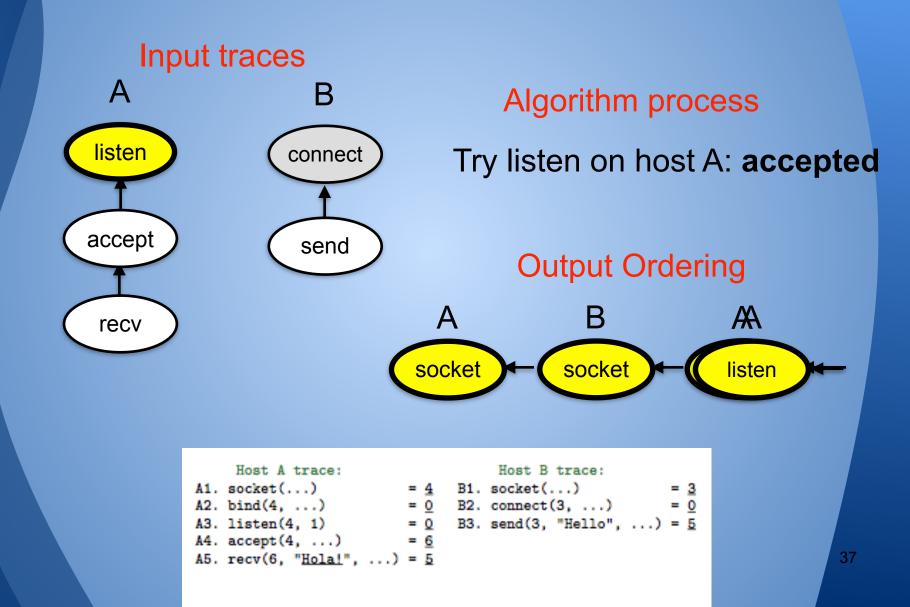


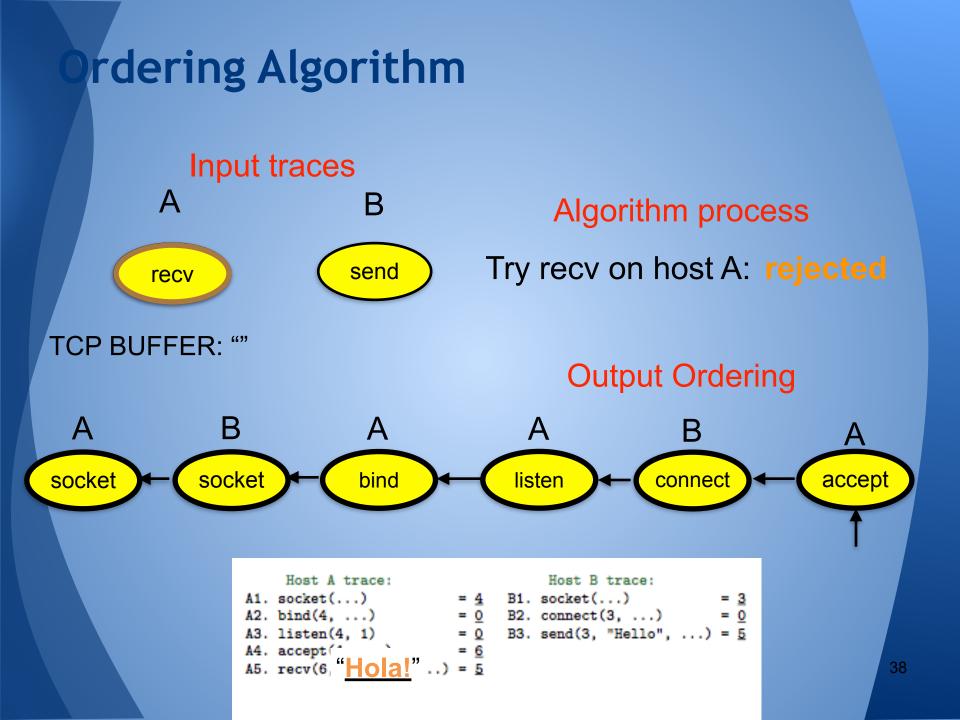
- Host B trace:
- = <u>4</u> B1. socket(...) = <u>3</u> = <u>0</u> B2. connect(3, ...) = <u>0</u> = <u>0</u> B3. send(3, "Hello", ...) = <u>5</u>

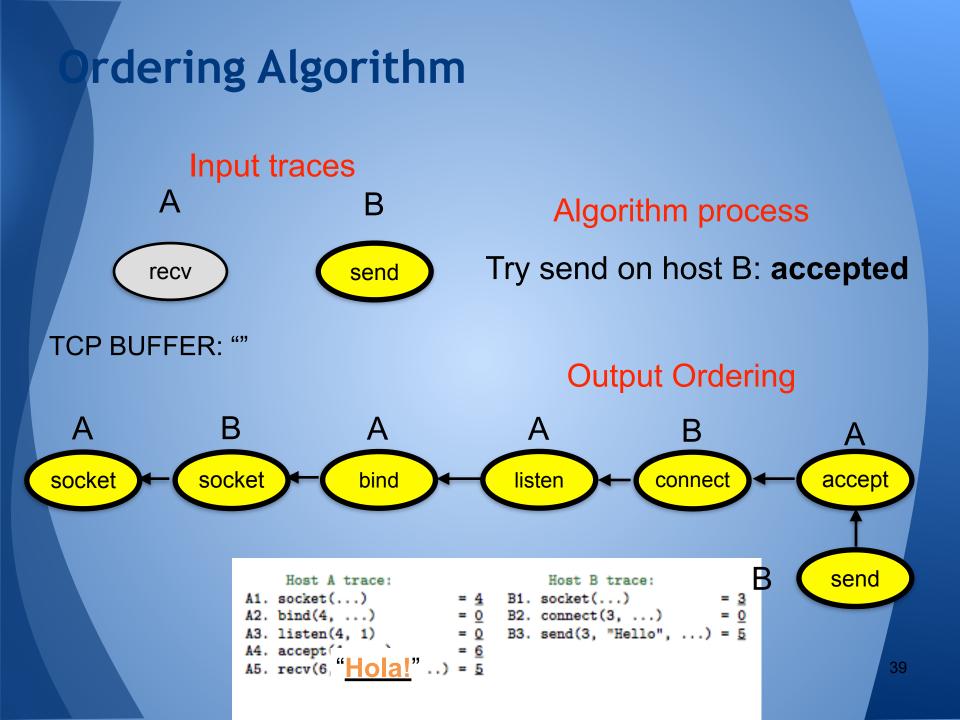
rdering Algorithm

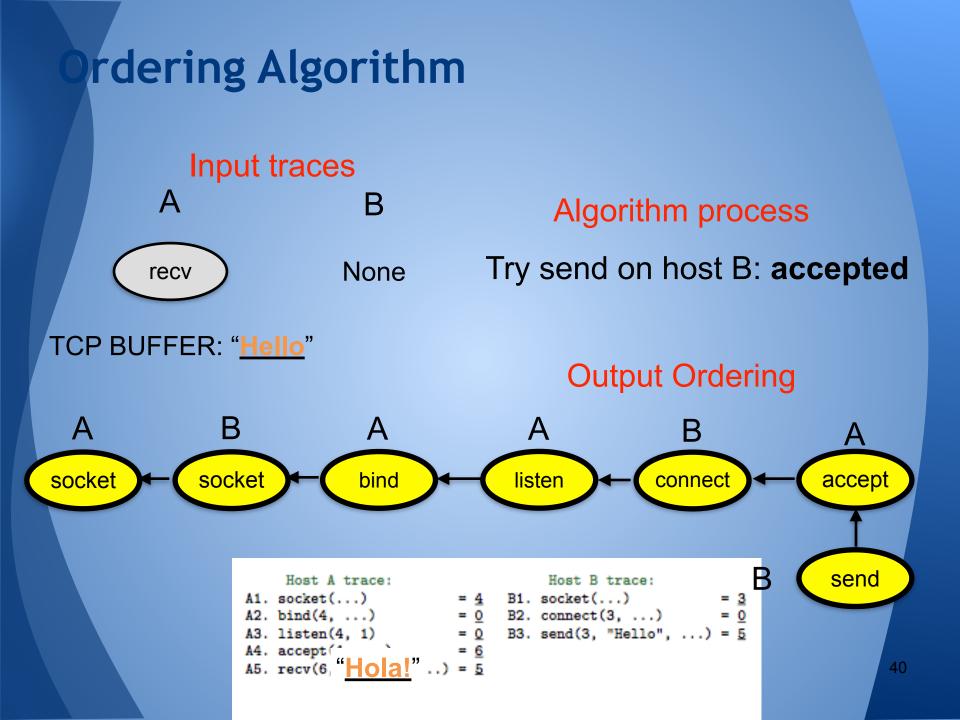


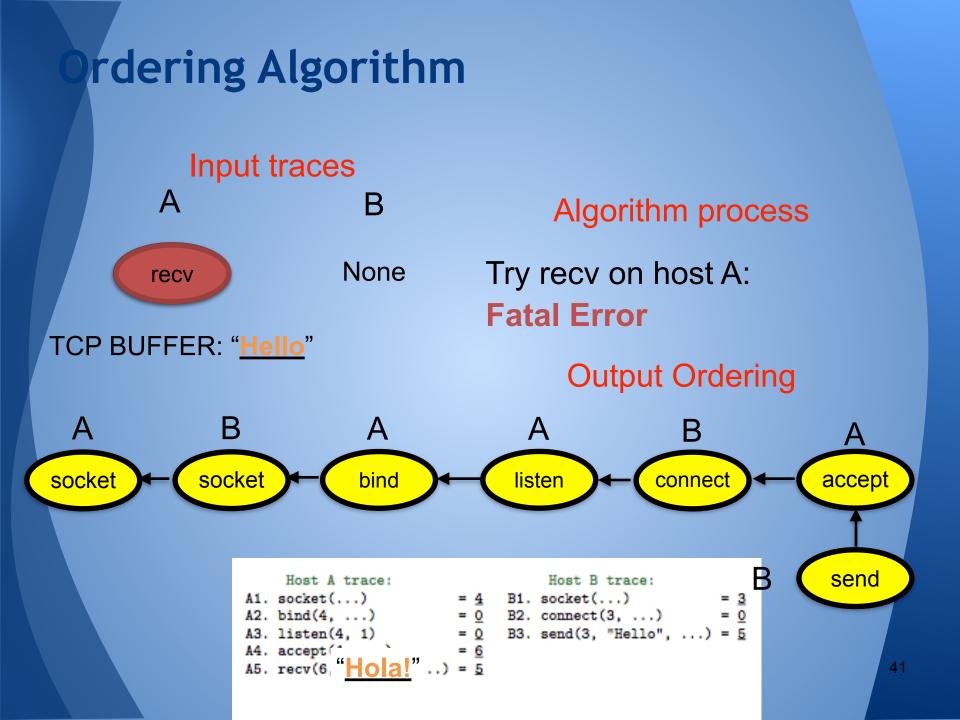
Prdering Algorithm





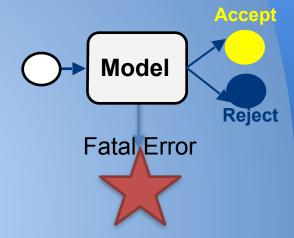






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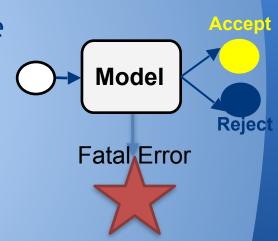
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etwork Model

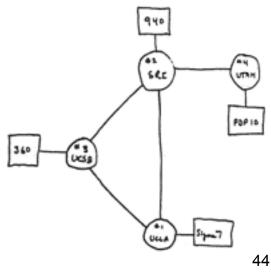
Simulates invocation of a syscall

- datagrams sent/lost
 - reordering / duplication is notable
- track pending connections
- buffer lengths and contents
 - send -> put data into buffer
 - recv -> pop data from buffer
- Simulation outcome
 - Accept → can process (correct buffer)
 - *Reject* → wrong order (incomplete buffer)
 - Permanent reject → abnormal behavior (incorrect buffer)



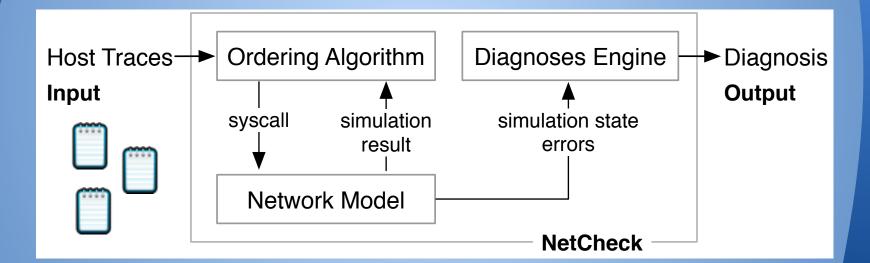
etwork Model

- Simulates invocation of a syscall
- Capture programmer assumptions
 - Assumes a simplified network view
 - Assume transitive connectivity
 - Little, random loss
 - No middle boxes
 - Assume uniform platform
 - Flag OS differences



ow Model Return Values Impact race Ordering

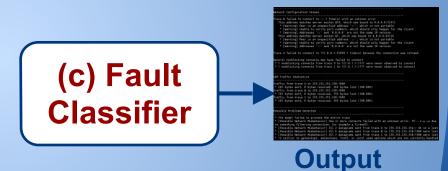
Blackbox Tracing mechanism



Trace Ordering: linear running time (total trace length) * number of traces

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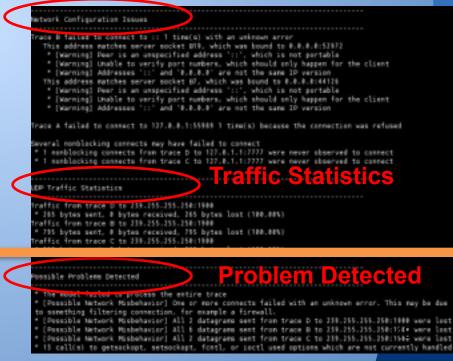
Fault Classifier

- Goal: Decide what to output
- Problem: Show relevant information
- Fault classifier: global (rather than local) view
 - uncovers high-level patterns by extracting low-level features
 - Examples: middleboxes, non-transitive connectivity, MTU, mobility, network disconnection
 - All look like loss, but have different patterns in the context of other flows

Fault Classifier

- Options to show different levels of detail
 Network admins / developers
 - detailed info
- End users
 - Classification
 - Recommendations

Network Configuration Issues



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Valuation: Production Application

- Reproduce reported bugs from bug trackers (Python, Apache, Ruby, Firefox, etc.)
 - A total of 71 bugs
 - Grouped into 23 categories
 - Virtualization incurred/portability bugs
 - SO_REUSEADDR behaves differently across OSes
 - accept inherit O_NONBLOCK
 -
 - Correct analysis of >95% bugs

valuation: Observed Network aults

- Twenty faults observed in practice on a live network
 - MTU bug
 - Intermediary device
 - Port forward
 - Traffic sent to non-relevant addresses
 - Provide supplemental info
 - packet loss
 - buffers being closed with data in
 - 90% of cases correctly detected

General Findings in Practice

Middle boxes

• Multiple unaccepted connections

client behind NAT in FTP
 TCP/UDP

non-transitive connectivity in VLC

Complex failures

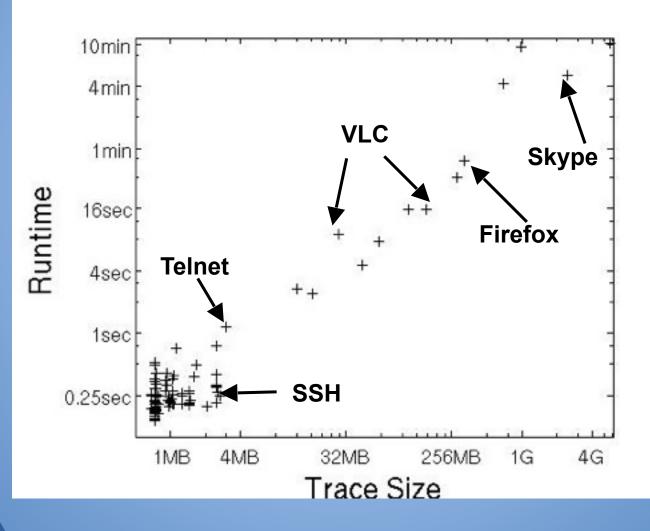
• VirtualBox send data larger than buffer size

• Pidgin returned IP different from bind

Skype NAT + close socket from a different thread
 Used on Seattle Testbed seattle.poly.edu



etCheck Performance Overhead



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Conclusion

Built and evaluated NetCheck, a tool to diagnose network failures in complex apps

Key insights:

model the programmer's misconceptions
 relation between calls → reconstruct order

NetCheck is effective

- Everyday applications & networks
- Real network / application bugs
- No per-network knowledge
- No per-application knowledge

Try it here: https://netcheck.poly.edu/