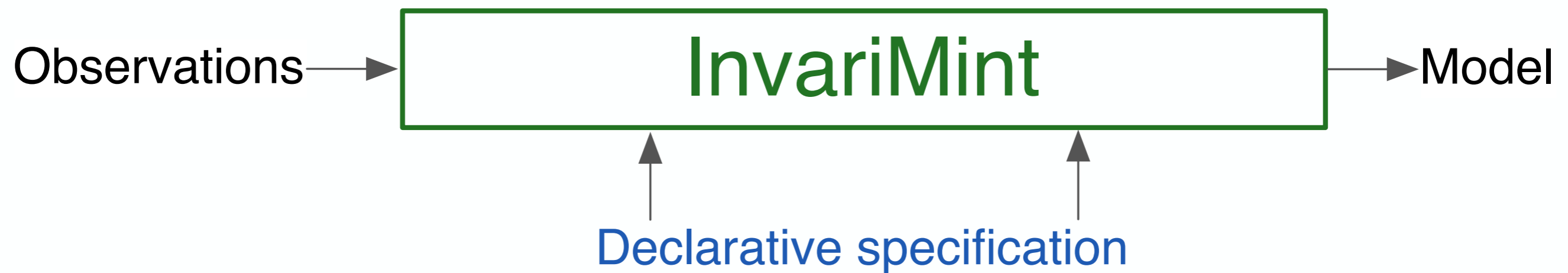


# Declarative Specification of FSM-Inference Algorithms



University of Washington

Ivan Beschastnikh  
Yuriy Brun  
Jenny Abrahamson  
Michael D. Ernst  
Arvind Krishnamurthy



UMass. Amherst













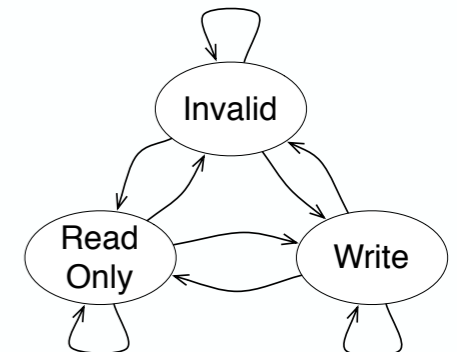
# FSM-inference representations

```
src : 2, dst : 0, timestamp : 0, type : prepare
src : 2, dst : 1, timestamp : 1, type : prepare
src : 0, dst : 2, timestamp : 2, type : commit
src : 1, dst : 2, timestamp : 3, type : commit
src : 2, dst : 0, timestamp : 4, type : tx_commit
src : 2, dst : 1, timestamp : 5, type : tx_commit
src : 0, dst : 2, timestamp : 6, type : ack
src : 1, dst : 2, timestamp : 7, type : ack
src : 2, dst : 0, timestamp : 8, type : prepare
src : 2, dst : 1, timestamp : 9, type : prepare
src : 0, dst : 2, timestamp : 10, type : commit
src : 1, dst : 2, timestamp : 11, type : commit
src : 2, dst : 0, timestamp : 12, type : tx_commit
src : 2, dst : 1, timestamp : 13, type : tx_commit
src : 0, dst : 2, timestamp : 14, type : ack
src : 1, dst : 2, timestamp : 15, type : ack
src : 2, dst : 0, timestamp : 16, type : prepare
src : 2, dst : 1, timestamp : 17, type : prepare
src : 0, dst : 2, timestamp : 18, type : commit
src : 1, dst : 2, timestamp : 19, type : commit
src : 2, dst : 0, timestamp : 20, type : tx_commit
src : 2, dst : 1, timestamp : 21, type : tx_commit
src : 0, dst : 2, timestamp : 22, type : ack
src : 1, dst : 2, timestamp : 23, type : ack
src : 2, dst : 0, timestamp : 0, type : prepare
src : 2, dst : 1, timestamp : 1, type : prepare
src : 0, dst : 2, timestamp : 2, type : commit
src : 1, dst : 2, timestamp : 3, type : commit
src : 2, dst : 0, timestamp : 4, type : tx_commit
```

Observations



Inference  
algorithm



Model



# FSM-inference representations

```
src : 2, dst : 0, timestamp : 0, type : prepare
src : 2, dst : 1, timestamp : 1, type : prepare
src : 0, dst : 2, timestamp : 2, type : commit
src : 1, dst : 2, timestamp : 3, type : commit
src : 2, dst : 0, timestamp : 4, type : tx_commit
src : 2, dst : 1, timestamp : 5, type : tx_commit
src : 0, dst : 2, timestamp : 6, type : ack
src : 1, dst : 2, timestamp : 7, type : ack
src : 2, dst : 0, timestamp : 8, type : prepare
src : 2, dst : 1, timestamp : 9, type : prepare
src : 0, dst : 2, timestamp : 10, type : commit
src : 1, dst : 2, timestamp : 11, type : commit
src : 2, dst : 0, timestamp : 12, type : tx_commit
src : 2, dst : 1, timestamp : 13, type : tx_commit
src : 0, dst : 2, timestamp : 14, type : ack
src : 1, dst : 2, timestamp : 15, type : ack
src : 2, dst : 0, timestamp : 16, type : prepare
src : 2, dst : 1, timestamp : 17, type : prepare
src : 0, dst : 2, timestamp : 18, type : commit
src : 1, dst : 2, timestamp : 19, type : commit
src : 2, dst : 0, timestamp : 20, type : tx_commit
src : 2, dst : 1, timestamp : 21, type : tx_commit
src : 0, dst : 2, timestamp : 22, type : ack
src : 1, dst : 2, timestamp : 23, type : ack
src : 2, dst : 0, timestamp : 0, type : prepare
src : 2, dst : 1, timestamp : 1, type : prepare
src : 0, dst : 2, timestamp : 2, type : commit
src : 1, dst : 2, timestamp : 3, type : commit
src : 2, dst : 0, timestamp : 4, type : tx_commit
```

Observations

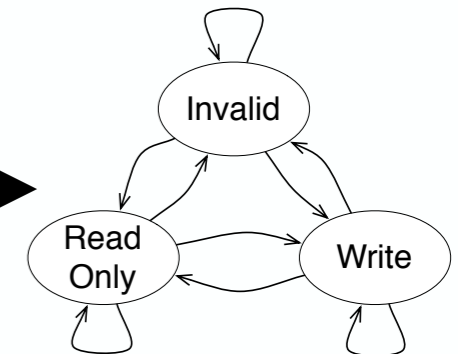


```
/**
 * Implements the KTails algorithm as defined in Biermann & Feldman '72.
 */
public class KTails {
    public static Logger logger;
    static {
        logger = Logger.getLogger("KTails");
    }

    /**
     * Constructs and returns a PartitionGraph generated by applying kTails with
     * the given k value to the
     */
    public static PartitionGraph generate(PartitionGraph g, int k) {
        PartitionGraph pGraph = g.clone();
        attemptMerge(pGraph, k);
        return pGraph;
    }

    /**
     * Finds all possible merges in pGraph. Requires making a new call to
     * attemptMerge after every merge in case previously un-merge-able pairs
     * become merge-able.
     */
    private static void attemptMerge(PartitionGraph pGraph, int k) {
        // Keeps track of the merges that we want to perform.
        Set<PartitionMultiMerge> merges = new LinkedHashSet<PartitionMultiMerge>();
    }
}
```

Code



Model

# FSM-inference representations

```

/**
 * Implements the KTails algorithm as defined in Biermann & Feldman '72.
 */
public class KTails {
    public static Logger logger;
    static {
        logger = Logger.getLogger("KTails");
    }

    /**
     * Constructs and returns a PartitionGraph by applying kTails with
     * the given k value to the given pGraph.
     */
    public static PartitionGraph kTails(PartitionGraph pGraph, int k) {
        PartitionGraph pGraph = kTails(pGraph, k, false, null);
        attemptMerge(pGraph, k);
        return pGraph;
    }

    /**
     * Finds all possible merges in pGraph. Requires making a new call to
     * attemptMerge after every merge in case previously un-merge-able pairs
     * become merge-able.
     */
    private static void attemptMerge(PartitionGraph pGraph, int k) {
        // Keeps track of the merges that we want to perform.
        Set<PartitionMultiMerge> merges = new LinkedHashSet<PartitionMultiMerge>();
    }
}

```

Code

```

src : 2, dst : 0, timestamp : 0, type : prepare
src : 2, dst : 1, timestamp : 1, type : prepare
src : 0, dst : 2, timestamp : 2, type : commit
src : 1, dst : 2, timestamp : 3, type : commit
src : 2, dst : 0, timestamp : 4, type : tx_commit
src : 2, dst : 1, timestamp : 5, type : tx_commit
src : 0, dst : 2, timestamp : 6, type : ack
src : 1, dst : 2, timestamp : 7, type : ack
src : 2, dst : 0, timestamp : 8, type : prepare
src : 2, dst : 1, timestamp : 9, type : prepare
src : 0, dst : 2, timestamp : 10, type : commit
src : 1, dst : 2, timestamp : 11, type : commit
src : 2, dst : 0, timestamp : 12, type : tx_commit
src : 2, dst : 1, timestamp : 13, type : tx_commit
src : 0, dst : 2, timestamp : 14, type : ack
src : 1, dst : 2, timestamp : 15, type : ack
src : 2, dst : 0, timestamp : 16, type : prepare
src : 2, dst : 1, timestamp : 17, type : prepare
src : 0, dst : 2, timestamp : 18, type : commit
src : 1, dst : 2, timestamp : 19, type : commit
src : 2, dst : 0, timestamp : 20, type : tx_commit
src : 2, dst : 1, timestamp : 21, type : tx_commit
src : 0, dst : 2, timestamp : 22, type : ack
src : 1, dst : 2, timestamp : 23, type : ack
src : 2, dst : 0, timestamp : 0, type : prepare
src : 2, dst : 1, timestamp : 1, type : prepare
src : 0, dst : 2, timestamp : 2, type : commit
src : 1, dst : 2, timestamp : 3, type : commit
src : 2, dst : 0, timestamp : 4, type : tx_commit

```

Observations



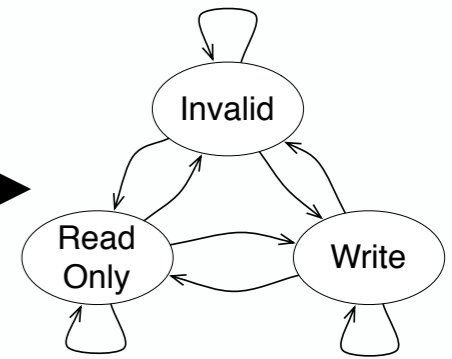
Formally, given two sequences of length  $k$   $seq1 = (x_1^1, P_1^1) \dots (x_k^1, P_k^1)$  and  $seq2 = (x_1^2, P_1^2) \dots (x_k^2, P_k^2)$ , we say that

- $seq1 = seq2$  iff  $\forall i = 1, \dots, k \ x_i^1 = x_i^2$  and  $P_i^1 \Leftrightarrow P_i^2$
- $seq1 \subseteq seq2$  iff  $\forall i = 1, \dots, k \ x_i^1 = x_i^2$  and  $P_i^1 \Rightarrow P_i^2$

Given two  $k$ -futures  $f1 = \{seq_1^1, \dots, seq_{n_1}^1\}$ , with  $seq_i^1 = (x_1^1, P_1^1) \dots (x_k^1, P_k^1)$  and  $f2 = \{seq_1^2, \dots, seq_{n_2}^2\}$ , with  $seq_j^2 = (x_1^2, P_1^2) \dots (x_k^2, P_k^2)$

**Set theory**

- $f1$  is equivalent to  $f2$  iff  $n_1 = n_2$  and  $\forall i = 1, \dots, n_1 \ \exists j = 1, \dots, n_2$  s.t.  $seq_i^1 = seq_j^2$
- $f1$  weakly subsumes  $f2$  iff  $n_1 = n_2$ ,  $\forall j = 1, \dots, n_1 \ \exists i = 1, \dots, n_2$  s.t.  $seq_j^1 \subseteq seq_i^2$ , and vice versa  $\forall i = 1, \dots, n_1 \ \exists j = 1, \dots, n_2$  s.t.  $seq_i^2 \subseteq seq_j^1$
- $f1$  strongly subsumes  $f2$  iff  $\forall j = 1, \dots, n_2 \ \exists i = 1, \dots, n_1$  s.t.  $seq_j^2 \subseteq seq_i^1$



Model

# FSM-inference representations

```

/**
 * Implements the KTails algorithm as defined in Biermann & Feldman '72.
 */
public class KTails {
    public static Logger logger;
    static {
        logger = Logger.getLogger("KTails");
    }

    /**
     * Constructs and returns a PartitionGraph represented by applying kTails with
     * the given k value to the given PartitionGraph.
     */
    public static PartitionGraph kTails(PartitionGraph pGraph, int k) {
        PartitionGraph pGraph = kTails(pGraph, k, false, null);
        attemptMerge(pGraph, k);
        return pGraph;
    }

    /**
     * Finds all possible merges in pGraph. Requires making a new call to
     * attemptMerge after every merge in case previously un-merge-able pairs
     * become merge-able.
     */
    private static void attemptMerge(PartitionGraph pGraph, int k) {
        // Keeps track of the merges that we want to perform.
        Set<PartitionMultiMerge> merges = new LinkedHashSet<PartitionMultiMerge>();
    }
}

```

Code

Formally, given two sequences of length  $k$   $seq1 = (x_1^1, P_1^1) \dots (x_k^1, P_k^1)$  and  $seq2 = (x_1^2, P_1^2) \dots (x_k^2, P_k^2)$ , we say that

- $seq1 = seq2$  iff  $\forall i = 1, \dots, k$   $x_i^1 = x_i^2$  and  $P_i^1 \Leftrightarrow P_i^2$
- $seq1 \subseteq seq2$  iff  $\forall i = 1, \dots, k$   $x_i^1 = x_i^2$  and  $P_i^1 \Rightarrow P_i^2$

Given two  $k$ -futures  $f1 = \{seq_1^1, \dots, seq_{n_1}^1\}$ , with  $seq_i^1 = (x_1^1, P_1^1) \dots (x_k^1, P_k^1)$  and  $f2 = \{seq_1^2, \dots, seq_{n_2}^2\}$  with  $seq_i^2 = (x_1^2, P_1^2) \dots (x_k^2, P_k^2)$

- $f1$  is equivalent to  $f2$  iff  $n_1 = n_2$  and  $\forall i = 1, \dots, n_1$   $\exists j = 1, \dots, n_2$  s.t.  $seq_i^1 = seq_j^2$
- $f1$  weakly subsumes  $f2$  iff  $n_1 = n_2$ ,  $\forall j = 1, \dots, n_1$   $\exists i = 1, \dots, n_2$  s.t.  $seq_j^1 \subseteq seq_i^2$ , and vice versa  $\forall i = 1, \dots, n_1$   $\exists j = 1, \dots, n_2$  s.t.  $seq_j^2 \subseteq seq_i^1$
- $f1$  strongly subsumes  $f2$  iff  $\forall j = 1, \dots, n_2$   $\exists i = 1, \dots, n_1$  s.t.  $seq_j^2 \subseteq seq_i^1$

Set theory

```

src : 2, dst : 0, timestamp : 0, type : prepare
src : 2, dst : 1, timestamp : 1, type : prepare
src : 0, dst : 2, timestamp : 2, type : commit
src : 1, dst : 2, timestamp : 3, type : commit
src : 2, dst : 0, timestamp : 4, type : tx_commit
src : 2, dst : 1, timestamp : 5, type : tx_commit
src : 0, dst : 2, timestamp : 6, type : ack
src : 1, dst : 2, timestamp : 7, type : ack
src : 2, dst : 0, timestamp : 8, type : prepare
src : 2, dst : 1, timestamp : 9, type : prepare
src : 0, dst : 2, timestamp : 10, type : commit
src : 1, dst : 2, timestamp : 11, type : commit
src : 2, dst : 0, timestamp : 12, type : tx_commit
src : 2, dst : 1, timestamp : 13, type : tx_commit
src : 0, dst : 2, timestamp : 14, type : ack
src : 1, dst : 2, timestamp : 15, type : ack
src : 2, dst : 0, timestamp : 16, type : prepare
src : 2, dst : 1, timestamp : 17, type : prepare
src : 0, dst : 2, timestamp : 18, type : commit
src : 1, dst : 2, timestamp : 19, type : commit
src : 2, dst : 0, timestamp : 20, type : tx_commit
src : 2, dst : 1, timestamp : 21, type : tx_commit
src : 0, dst : 2, timestamp : 22, type : ack
src : 1, dst : 2, timestamp : 23, type : ack
src : 2, dst : 0, timestamp : 0, type : prepare
src : 2, dst : 1, timestamp : 1, type : prepare
src : 0, dst : 2, timestamp : 2, type : commit
src : 1, dst : 2, timestamp : 3, type : commit
src : 2, dst : 0, timestamp : 4, type : tx_commit

```

Observations

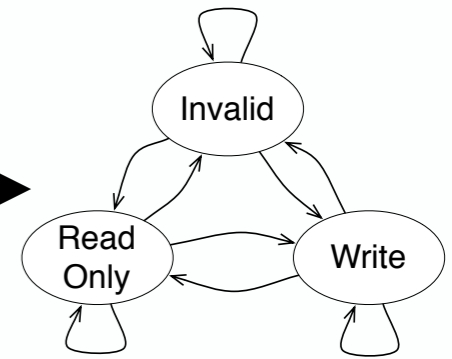


```

1  Input: event log L
2  let initialGraph = extract(L)
3  let I = mineInvariants(initialGraph)
4  let (V, E) = partition(initialGraph)
5  while (V, E) does not satisfy invariants I
6      // p: event → boolean, π: partition that will be split
7      let (p, π) = selectSplit((V, E), I)
8      let π1 =
9      let π2 =
10     V := (V \ π1) ∪ π2
11     E := { (π3, π4, r) ∈ V × V × R | ∃ event1 ∈ π3, ∃ event2 ∈ π4
12           : event1 r event2 ∈ initialGraph }
13 end while
14 if (hybrid)
15     (V, E) := kTail((V, E), 0, I)
16 Output: (V, E)

```

Pseudo-code



Model

# FSM-inference representations

```

/**
 * Implements the KTails algorithm as defined in Biermann & Feldman '72.
 */
public class KTails {
    public static Logger logger;
    static {
        logger = Logger.getLogger("KTails");
    }

    /**
     * Constructs and returns a PartitionGraph represented by applying kTails with
     * the given k value to the given graph.
     */
    public static PartitionGraph kTails(PartitionGraph pGraph, int k) {
        PartitionGraph pGraph = kTails(pGraph, k, false, null);
        attemptMerge(pGraph, k);
        return pGraph;
    }

    /**
     * Finds all possible merges in pGraph. Requires making a new call to
     * attemptMerge after every merge in case previously un-merge-able pairs
     * become merge-able.
     */
    private static void attemptMerge(PartitionGraph pGraph, int k) {
        // Keeps track of the merges that we want to perform.
        Set<PartitionMultiMerge> merges = new LinkedHashSet<PartitionMultiMerge>();
    }
}

```

Code

Formally, given two sequences of length  $k$   $seq1 = (x_1^1, P_1^1) \dots (x_k^1, P_k^1)$  and  $seq2 = (x_1^2, P_1^2) \dots (x_k^2, P_k^2)$ , we say that

- $seq1 = seq2$  iff  $\forall i = 1, \dots, k$   $x_i^1 = x_i^2$  and  $P_i^1 \Leftrightarrow P_i^2$
- $seq1 \subseteq seq2$  iff  $\forall i = 1, \dots, k$   $x_i^1 = x_i^2$  and  $P_i^1 \Rightarrow P_i^2$

Given two  $k$ -futures  $f1 = \{seq_1^1, \dots, seq_{n_1}^1\}$ , with  $seq_i^1 = (x_1^1, P_1^1) \dots (x_i^1, P_i^1)$  and  $f2 = \{seq_1^2, \dots, seq_{n_2}^2\}$  with  $seq_j^2 = (x_1^2, P_1^2) \dots (x_j^2, P_j^2)$

- $f1$  is eq  $f2$  iff  $n_1 = n_2$  and  $\exists j = 1, \dots, n_1$  s.t.  $seq_j^2 \subseteq seq_i^1$
- $f1$  weakly subsumes  $f2$  iff  $n_1 = n_2$ ,  $\forall j = 1, \dots, n_1$   $\exists i = 1, \dots, n_1$  s.t.  $seq_j^2 \subseteq seq_i^1$ , and vice versa  $\forall i = 1, \dots, n_1$   $\exists j = 1, \dots, n_1$  s.t.  $seq_j^2 \subseteq seq_i^1$
- $f1$  strongly subsumes  $f2$  iff  $\forall j = 1, \dots, n_2$   $\exists i = 1, \dots, n_1$  s.t.  $seq_j^2 \subseteq seq_i^1$

Set theory

```

1 Input: event log L
2 let initialGraph = extract(L)
3 let I = mineInvariants(initialGraph)
4 let (V, E) = partition(initialGraph)
5 while (V, E) does not satisfy invariants I
6   // p: event → boolean, π: partition that will be split
7   let (p, π) = findSplit(V, E, I)
8   let π' = merge(π, p)
9   let π'' = merge(π', p)
10  V := V - {event | event ∈ π''}
11  E := {(π3, π4, r) ∈ V × V × R | ∃ event1 ∈ π3, ∃ event2 ∈ π4
12    : event1 r event2 ∈ initialGraph}
13 end while
14 if (hybrid)
15   (V, E) := kTail((V, E), 0, I)
16 Output: (V, E)

```

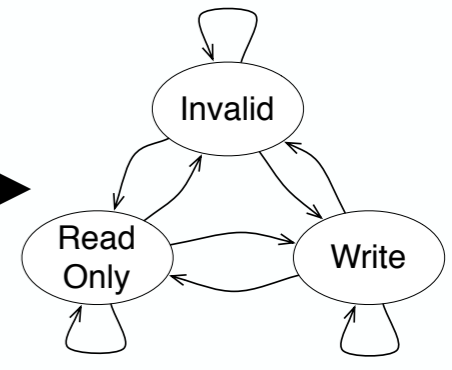
Pseudo-code

```

src : 2, dst : 0, timestamp : 0, type : prepare
src : 2, dst : 1, timestamp : 1, type : prepare
src : 0, dst : 2, timestamp : 2, type : commit
src : 1, dst : 2, timestamp : 3, type : commit
src : 2, dst : 0, timestamp : 4, type : tx_commit
src : 2, dst : 1, timestamp : 5, type : tx_commit
src : 0, dst : 2, timestamp : 6, type : ack
src : 1, dst : 2, timestamp : 7, type : ack
src : 2, dst : 0, timestamp : 8, type : prepare
src : 2, dst : 1, timestamp : 9, type : prepare
src : 0, dst : 2, timestamp : 10, type : commit
src : 1, dst : 2, timestamp : 11, type : commit
src : 2, dst : 0, timestamp : 12, type : tx_commit
src : 2, dst : 1, timestamp : 13, type : tx_commit
src : 0, dst : 2, timestamp : 14, type : ack
src : 1, dst : 2, timestamp : 15, type : ack
src : 2, dst : 0, timestamp : 16, type : prepare
src : 2, dst : 1, timestamp : 17, type : prepare
src : 0, dst : 2, timestamp : 18, type : commit
src : 1, dst : 2, timestamp : 19, type : commit
src : 2, dst : 0, timestamp : 20, type : tx_commit
src : 2, dst : 1, timestamp : 21, type : tx_commit
src : 0, dst : 2, timestamp : 22, type : ack
src : 1, dst : 2, timestamp : 23, type : ack
src : 2, dst : 0, timestamp : 0, type : prepare
src : 2, dst : 1, timestamp : 1, type : prepare
src : 0, dst : 2, timestamp : 2, type : commit
src : 1, dst : 2, timestamp : 3, type : commit
src : 2, dst : 0, timestamp : 4, type : tx_commit

```

Observations



Model

# FSM-inference representations

```

/**
 * Implements the KTails algorithm as defined in Biermann & Feldman '72.
 */
public class KTails {
    public static Logger logger;
    static {
        logger = Logger.getLogger("KTails");
    }

    /**
     * Constructs and returns a PartitionGraph represented by applying kTails with
     * the given k value to the given PartitionGraph.
     */
    public static PartitionGraph kTails(PartitionGraph pGraph, int k) {
        PartitionGraph pGraph = kTails(pGraph, k, false, null);
        attemptMerge(pGraph, k);
        return pGraph;
    }

    /**
     * Finds all possible merges in pGraph. Requires making a new call to
     * attemptMerge after every merge in case previously un-merge-able pairs
     * become merge-able.
     */
    private static void attemptMerge(PartitionGraph pGraph, int k) {
        // Keeps track of the merges that we want to perform.
        Set<PartitionMultiMerge> merges = new LinkedHashSet<PartitionMultiMerge>();
    }
}

```

Code

Formally, given two sequences of length  $k$   $seq1 = (x_1^1, P_1^1) \dots (x_k^1, P_k^1)$  and  $seq2 = (x_1^2, P_1^2) \dots (x_k^2, P_k^2)$ , we say that

- $seq1 = seq2$  iff  $\forall i = 1, \dots, k$   $x_i^1 = x_i^2$  and  $P_i^1 \Leftrightarrow P_i^2$
- $seq1 \subseteq seq2$  iff  $\forall i = 1, \dots, k$   $x_i^1 = x_i^2$  and  $P_i^1 \Rightarrow P_i^2$

Given two  $k$ -futures  $f1 = \{seq_1^1, \dots, seq_{n_1}^1\}$ , with  $seq_i^1 = (x_1^1, P_1^1) \dots (x_i^1, P_i^1)$  and  $f2 = \{seq_1^2, \dots, seq_{n_2}^2\}$  with  $seq_j^2 = (x_1^2, P_1^2) \dots (x_j^2, P_j^2)$

- $f1$  is equal to  $f2$  iff  $n_1 = n_2$  and  $\forall j = 1, \dots, n_1$   $\exists i = 1, \dots, n_2$  s.t.  $seq_j^1 = seq_i^2$
- $f1$  weakly subsumes  $f2$  iff  $n_1 = n_2$ ,  $\forall j = 1, \dots, n_1$   $\exists i = 1, \dots, n_2$  s.t.  $seq_j^1 \subseteq seq_i^2$ , and vice versa  $\forall i = 1, \dots, n_2$   $\exists j = 1, \dots, n_1$  s.t.  $seq_j^1 \subseteq seq_i^2$
- $f1$  strongly subsumes  $f2$  iff  $\forall j = 1, \dots, n_2$   $\exists i = 1, \dots, n_1$  s.t.  $seq_j^2 \subseteq seq_i^1$

Set theory

```

1 Input: event log L
2 let initialGraph = extract(L)
3 let I = mineInvariants(initialGraph)
4 let (V, E) = partition(initialGraph)
5 while (V, E) does not satisfy invariants I
6   // p: event → boolean, π: partition that will be split
7   let (p, π) = findSplit(V, E, I)
8   let π' = refine(π, p)
9   let π'' = merge(π', p)
10  V := V - {event | event ∈ π''}
11  E := {(π3, π4, r) ∈ V × V × R | ∃ event1 ∈ π3, ∃ event2 ∈ π4
12      : event1 r event2 ∈ initialGraph}
13 end while
14 if (hybrid)
15   (V, E) := kTail((V, E), 0, I)
16 Output: (V, E)

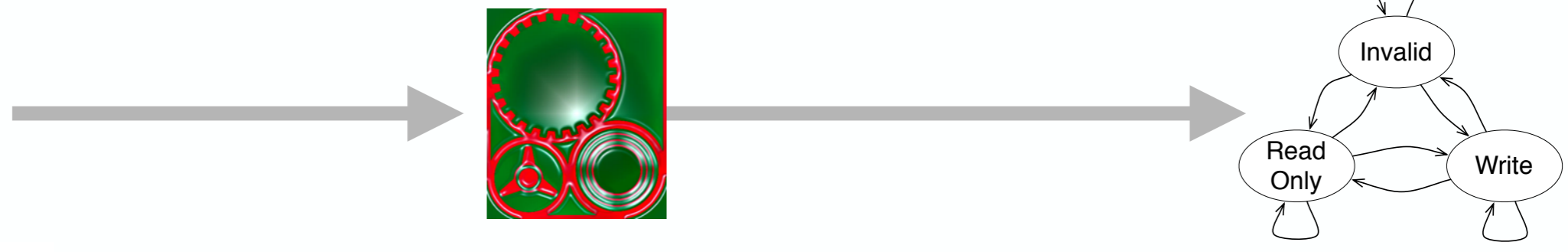
```

Pseudo-code

```

src : 2, dst : 0, timestamp : 0, type : prepare
src : 2, dst : 1, timestamp : 1, type : prepare
src : 0, dst : 2, timestamp : 2, type : commit
src : 1, dst : 2, timestamp : 3, type : commit
src : 2, dst : 0, timestamp : 4, type : tx_commit
src : 2, dst : 1, timestamp : 5, type : tx_commit
src : 0, dst : 2, timestamp : 6, type : ack
src : 1, dst : 2, timestamp : 7, type : ack
src : 2, dst : 0, timestamp : 8, type : prepare
src : 2, dst : 1, timestamp : 9, type : prepare
src : 0, dst : 2, timestamp : 10, type : commit
src : 1, dst : 2, timestamp : 11, type : commit
src : 2, dst : 0, timestamp : 12, type : tx_commit
src : 2, dst : 1, timestamp : 13, type : tx_commit
src : 0, dst : 2, timestamp : 14, type : ack
src : 1, dst : 2, timestamp : 15, type : ack
src : 2, dst : 0, timestamp : 16, type : prepare
src : 2, dst : 1, timestamp : 17, type : prepare
src : 0, dst : 2, timestamp : 18, type : commit
src : 1, dst : 2, timestamp : 19, type : commit
src : 2, dst : 0, timestamp : 20, type : tx_commit
src : 2, dst : 1, timestamp : 21, type : tx_commit
src : 0, dst : 2, timestamp : 22, type : ack
src : 1, dst : 2, timestamp : 23, type : ack
src : 2, dst : 0, timestamp : 0, type : prepare
src : 2, dst : 1, timestamp : 1, type : prepare
src : 0, dst : 2, timestamp : 2, type : commit
src : 1, dst : 2, timestamp : 3, type : commit
src : 2, dst : 0, timestamp : 4, type : tx_commit

```



Observations

Model

Limitations:

- ✗ Transparency
- ✗ Extensibility

# FSM-inference representations

```

/**
 * Implements the KTails algorithm as defined in Biermann & Feldman '72.
 */
public class KTails {
    public static Logger logger;
    static {
        logger = Logger.getLogger("KTails");
    }

    /**
     * Constructs and returns a PartitionGraph represented by applying kTails with
     * the given k value to the given PartitionGraph.
     */
    public static PartitionGraph kTails(PartitionGraph pGraph, int k) {
        PartitionGraph pGraph = kTails(pGraph, k, false, null);
        attemptMerge(pGraph, k);
        return pGraph;
    }

    /**
     * Finds all possible merges in pGraph. Requires making a new call to
     * attemptMerge after every merge in case previously un-merge-able pairs
     * become merge-able.
     */
    private static void attemptMerge(PartitionGraph pGraph, int k) {
        // Keeps track of the merges that we want to perform.
        Set<PartitionMultiMerge> merges = new LinkedHashSet<PartitionMultiMerge>();
    }
}

```

Code

Formally, given two sequences of length  $k$   $seq1 = (x_1^1, P_1^1) \dots (x_k^1, P_k^1)$  and  $seq2 = (x_1^2, P_1^2) \dots (x_k^2, P_k^2)$ , we say that

- $seq1 = seq2$  iff  $\forall i = 1, \dots, k$   $x_i^1 = x_i^2$  and  $P_i^1 \Leftrightarrow P_i^2$
- $seq1 \subseteq seq2$  iff  $\forall i = 1, \dots, k$   $x_i^1 = x_i^2$  and  $P_i^1 \Rightarrow P_i^2$

Given two  $k$ -futures  $f1 = \{seq_1^1, \dots, seq_{n_1}^1\}$ , with  $seq_i^1 = (x_1^1, P_1^1) \dots (x_i^1, P_i^1)$  and  $f2 = \{seq_1^2, \dots, seq_{n_2}^2\}$  with  $seq_j^2 = (x_1^2, P_1^2) \dots (x_j^2, P_j^2)$

- $f1$  is equivalent to  $f2$  iff  $n_1 = n_2$  and  $\forall j = 1, \dots, n_1$   $\exists i = 1, \dots, n_2$  s.t.  $seq_j^1 \subseteq seq_i^2$  and vice versa  $\forall i = 1, \dots, n_2$   $\exists j = 1, \dots, n_1$  s.t.  $seq_i^2 \subseteq seq_j^1$
- $f1$  weakly subsumes  $f2$  iff  $n_1 = n_2$ ,  $\forall j = 1, \dots, n_1$   $\exists i = 1, \dots, n_2$  s.t.  $seq_j^1 \subseteq seq_i^2$ , and vice versa  $\forall i = 1, \dots, n_2$   $\exists j = 1, \dots, n_1$  s.t.  $seq_i^2 \subseteq seq_j^1$
- $f1$  strongly subsumes  $f2$  iff  $\forall j = 1, \dots, n_2$   $\exists i = 1, \dots, n_1$  s.t.  $seq_j^2 \subseteq seq_i^1$

Set theory

```

1 Input: event log L
2 let initialGraph = extract(L)
3 let I = mineInvariants(initialGraph)
4 let (V, E) = partition(initialGraph)
5 while (V, E) does not satisfy invariants I
6   // p: event → boolean, π: partition that will be split
7   let (p, π) = findSplit(V, E, I)
8   let π' = merge(π, p)
9   let π'' = merge(π', p)
10  V := V - {event | event ∈ π''}
11  E := {(π3, π4, r) ∈ V × V × R | ∃ event1 ∈ π3, ∃ event2 ∈ π4
12      : event1 r event2 ∈ initialGraph}
13 end while
14 if (hybrid)
15   (V, E) := kTail((V, E), 0, I)
16 Output: (V, E)

```

Pseudo-code

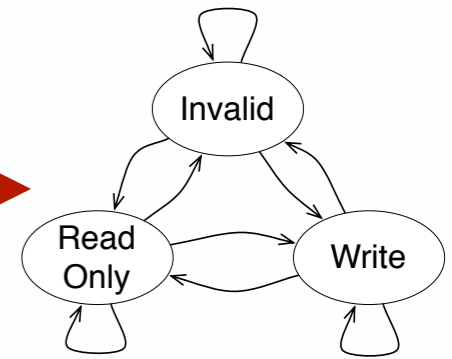
```

src : 2, dst : 0, timestamp : 0, type : prepare
src : 2, dst : 1, timestamp : 1, type : prepare
src : 0, dst : 2, timestamp : 2, type : commit
src : 1, dst : 2, timestamp : 3, type : commit
src : 2, dst : 0, timestamp : 4, type : tx_commit
src : 2, dst : 1, timestamp : 5, type : tx_commit
src : 0, dst : 2, timestamp : 6, type : ack
src : 1, dst : 2, timestamp : 7, type : ack
src : 2, dst : 0, timestamp : 8, type : prepare
src : 0, dst : 2, timestamp : 10, type : commit
src : 1, dst : 2, timestamp : 11, type : commit

```

Observed trace

Trace admitted by the model?



Model

Observations

Limitations:

- ✗ Transparency
- ✗ Extensibility

# FSM-inference representations

```

/**
 * Implements the KTails algorithm as defined in Biermann & Feldman '72.
 */
public class KTails {
    public static Logger logger;
    static {
        logger = Logger.getLogger("KTails");
    }

    /**
     * Constructs and returns a PartitionGraph represented by applying kTails with
     * the given k value to the given graph.
     */
    public static PartitionGraph kTails(PartitionGraph pGraph, int k) {
        PartitionGraph pGraph = kTails(pGraph, k, false, null);
        attemptMerge(pGraph, k);
        return pGraph;
    }

    /**
     * Finds all possible merges in pGraph. Requires making a new call to
     * attemptMerge after every merge in case previously un-merge-able pairs
     * become merge-able.
     */
    private static void attemptMerge(PartitionGraph pGraph, int k) {
        // Keeps track of the merges that we want to perform.
        Set<PartitionMultiMerge> merges = new LinkedHashSet<PartitionMultiMerge>();
    }
}

```

Code

Formally, given two sequences of length  $k$   $seq1 = (x_1^1, P_1^1) \dots (x_k^1, P_k^1)$  and  $seq2 = (x_1^2, P_1^2) \dots (x_k^2, P_k^2)$ , we say that

- $seq1 = seq2$  iff  $\forall i = 1, \dots, k$   $x_i^1 = x_i^2$  and  $P_i^1 \Leftrightarrow P_i^2$
- $seq1 \subseteq seq2$  iff  $\forall i = 1, \dots, k$   $x_i^1 = x_i^2$  and  $P_i^1 \Rightarrow P_i^2$

Given two  $k$ -futures  $f1 = \{seq_1^1, \dots, seq_{n_1}^1\}$ , with  $seq_i^1 = (x_1^1, P_1^1) \dots (x_i^1, P_i^1)$  and  $f2 = \{seq_1^2, \dots, seq_{n_2}^2\}$  with  $seq_j^2 = (x_1^2, P_1^2) \dots (x_j^2, P_j^2)$

- $f1$  is equal to  $f2$  iff  $n_1 = n_2$  and  $\forall j = 1, \dots, n_1$   $\exists i = 1, \dots, n_2$  s.t.  $seq_j^1 = seq_i^2$
- $f1$  weakly subsumes  $f2$  iff  $n_1 = n_2$ ,  $\forall j = 1, \dots, n_1$   $\exists i = 1, \dots, n_2$  s.t.  $seq_j^2 \subseteq seq_i^1$ , and vice versa  $\forall i = 1, \dots, n_2$   $\exists j = 1, \dots, n_1$  s.t.  $seq_j^2 \subseteq seq_i^1$
- $f1$  strongly subsumes  $f2$  iff  $\forall j = 1, \dots, n_2$   $\exists i = 1, \dots, n_1$  s.t.  $seq_j^2 \subseteq seq_i^1$

Set theory

```

1 Input: event log L
2 let initialGraph = extract(L)
3 let I = mineInvariants(initialGraph)
4 let (V, E) = partition(initialGraph)
5 while (V, E) does not satisfy invariants I
6   // p: event → boolean, π: partition that will be split
7   let (p, π) = findSplit(V, E, I)
8   let π' = merge(π, p)
9   let π'' = merge(π', p)
10  V := V - {event | event ∈ π''}
11  E := {(π3, π4, r) ∈ V × V × R | ∃ event1 ∈ π3, ∃ event2 ∈ π4
12      : event1 r event2 ∈ initialGraph}
13 end while
14 if (hybrid)
15   (V, E) := kTail((V, E), 0, I)
16 Output: (V, E)

```

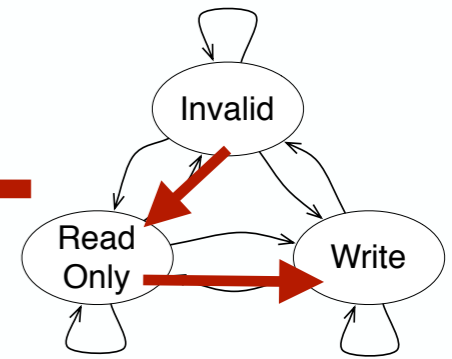
Pseudo-code

```

src : 2, dst : 0, timestamp : 0, type : prepare
src : 2, dst : 1, timestamp : 1, type : prepare
src : 0, dst : 2, timestamp : 2, type : commit
src : 1, dst : 2, timestamp : 3, type : commit
src : 2, dst : 0, timestamp : 4, type : tx_commit
src : 2, dst : 1, timestamp : 5, type : tx_commit
src : 0, dst : 2, timestamp : 6, type : ack
src : 1, dst : 2, timestamp : 7, type : ack
src : 2, dst : 0, timestamp : 8, type : prepare
src : 2, dst : 1, timestamp : 9, type : prepare
src : 0, dst : 2, timestamp : 10, type : commit
src : 1, dst : 2, timestamp : 11, type : commit
src : 2, dst : 0, timestamp : 12, type : tx_commit
src : 2, dst : 1, timestamp : 13, type : tx_commit
src : 0, dst : 2, timestamp : 14, type : ack
src : 1, dst : 2, timestamp : 15, type : ack
src : 2, dst : 0, timestamp : 16, type : prepare
src : 2, dst : 1, timestamp : 17, type : prepare
src : 0, dst : 2, timestamp : 18, type : commit
src : 1, dst : 2, timestamp : 19, type : commit
src : 2, dst : 0, timestamp : 20, type : tx_commit
src : 2, dst : 1, timestamp : 21, type : tx_commit
src : 0, dst : 2, timestamp : 22, type : ack
src : 1, dst : 2, timestamp : 23, type : ack
src : 2, dst : 0, timestamp : 0, type : prepare
src : 2, dst : 1, timestamp : 1, type : prepare
src : 0, dst : 2, timestamp : 2, type : commit
src : 1, dst : 2, timestamp : 3, type : commit
src : 2, dst : 0, timestamp : 4, type : tx_commit

```

Why is this execution admitted?



Observations

Model

Limitations:

- ✗ Transparency
- ✗ Extensibility

# FSM-inference representations

```

/**
 * Implements the KTails algorithm as defined in Biermann & Feldman '72.
 */
public class KTails {
    public static Logger logger;
    static {
        logger = Logger.getLogger("KTails");
    }

    // ...

    /**
     * Constructs
     * the given i
     */
    public static
    Partition(
    attemptMe)
    return pGraph;

    // ...

    /**
     * Finds all possible merges in pGraph. Requires making a new call to
     * attemptMerge after every merge in case previously un-merge-able pairs
     * become merge-able.
     */
    private static void attemptMerge(PartitionGraph pGraph, int k) {
        // Keeps track of the merges that we want to perform.
        Set<PartitionMultiMerge> merges = new LinkedHashSet<PartitionMultiMerge>();
    }
}

```

Feature A

Formally, given two sequences of length  $k$   $seq1 = (x_1, P_1) \dots (x_k, P_k)$  and  $seq2 = (x_1^2, P_1^2) \dots (x_k^2, P_k^2)$ , we say that

- $seq1 = seq2$  iff  $\forall i = 1, \dots, k$   $x_i^1 = x_i^2$  and  $P_i^1 \Leftrightarrow P_i^2$
- $seq1 \subseteq seq2$  iff  $\forall i = 1, \dots, k$   $x_i^1 = x_i^2$  and  $P_i^1 \Rightarrow P_i^2$

Given  $t$   $seq_i^1 =$   
 $x_1^1, P_1^1) \dots$   $seq_i^2 =$   
 $x_1^2, P_1^2) \dots$

- $f1$  is  $\exists j =$   
 $1, \dots$
- $f1$  weakly subsumes  $f2$  iff  $n_1 = n_2, \forall j = 1, \dots, n_1 \exists i =$   
 $1, \dots, n_1$  s.t.  $seq_j^2 \subseteq seq_i^1$ , and vice versa  $\forall i = 1, \dots, n_1$   
 $\exists j = 1, \dots, n_1$  s.t.  $seq_j^2 \subseteq seq_i^1$
- $f1$  strongly subsumes  $f2$  iff  $\forall j = 1, \dots, n_2 \exists i = 1, \dots, n_1$   
s.t.  $seq_j^2 \subseteq seq_i^1$

Base Algorithm

```

1 Input: event log L
2 let initialGraph = extract(L)
3 let I = mineInvariants(initialGraph)
4 let (V, E) = partition(initialGraph)
5 while (V, E) does not satisfy invariants I
6     // ...
7
8
9
10
11 E := {(π3, π4, r) ∈ V × V × K | ∃ event1 ∈ π3, ∃ event2 ∈ π4
12     : event1 revent2 ∈ initialGraph}
13 end while
14 if (hybrid)
15     (V, E) := kTail((V, E), 0, I)
16 Output: (V, E)

```

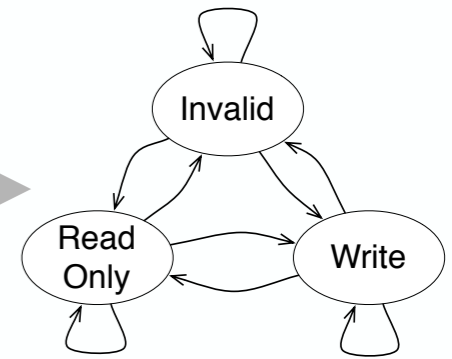
Feature B

```

src : 2, dst : 0, timestamp : 0, type : prepare
src : 2, dst : 1, timestamp : 1, type : prepare
src : 0, dst : 2, timestamp : 2, type : commit
src : 1, dst : 2, timestamp : 3, type : commit
src : 2, dst : 0, timestamp : 4, type : tx_commit
src : 2, dst : 1, timestamp : 5, type : tx_commit
src : 0, dst : 2, timestamp : 6, type : ack
src : 1, dst : 2, timestamp : 7, type : ack
src : 2, dst : 0, timestamp : 8, type : prepare
src : 2, dst : 1, timestamp : 9, type : prepare
src : 0, dst : 2, timestamp : 10, type : commit
src : 1, dst : 2, timestamp : 11, type : commit
src : 2, dst : 0, timestamp : 12, type : tx_commit
src : 2, dst : 1, timestamp : 13, type : tx_commit
src : 0, dst : 2, timestamp : 14, type : ack
src : 1, dst : 2, timestamp : 15, type : ack
src : 2, dst : 0, timestamp : 16, type : prepare
src : 2, dst : 1, timestamp : 17, type : prepare
src : 0, dst : 2, timestamp : 18, type : commit
src : 1, dst : 2, timestamp : 19, type : commit
src : 2, dst : 0, timestamp : 20, type : tx_commit
src : 2, dst : 1, timestamp : 21, type : tx_commit
src : 0, dst : 2, timestamp : 22, type : ack
src : 1, dst : 2, timestamp : 23, type : ack
src : 2, dst : 0, timestamp : 0, type : prepare
src : 2, dst : 1, timestamp : 1, type : prepare
src : 0, dst : 2, timestamp : 2, type : commit
src : 1, dst : 2, timestamp : 3, type : commit
src : 2, dst : 0, timestamp : 4, type : tx_commit

```

How to add features to an algorithm?



Observations

Model

Limitations:

- ✗ Transparency
- ✗ Extensibility



# FSM-inference representations

```
/**  
 * Implements the KTails algorithm as defined in Biermann & Feldman '72.  
 */  
public class KTails {  
    public static Logger logger;  
    static {  
        logger = Logger.getLogger("KTails");  
    }  
}
```

Formally, given two sequences of length  $k$   $seq1 = (x_1^1, P_1^1) \dots (x_k^1, P_k^1)$  and  $seq2 = (x_1^2, P_1^2) \dots (x_k^2, P_k^2)$ , we say that

- $seq1 = seq2$  iff  $\forall i = 1, \dots, k$   $x_i^1 = x_i^2$  and  $P_i^1 \Leftrightarrow P_i^2$
- $seq1 \subseteq seq2$  iff  $\forall i = 1, \dots, k$   $x_i^1 = x_i^2$  and  $P_i^1 \Rightarrow P_i^2$

```
1 Input: event log  $L$   
2 let initialGraph = extract( $L$ )  
3 let  $I$  = mineInvariants(initialGraph)  
4 let  $(V, E)$  = partition(initialGraph)  
5 while  $(V, E)$  does not satisfy invariants  $I$ 
```

## 1. Low-level specifications

```
/*  
 * Finds all possible merges in pGraph. Requires making a new call to  
 * attemptMerge after every merge in case previously un-merge-able pairs  
 * become merge-able.  
 */  
private static void attemptMerge(PartitionGraph pGraph, int k) {  
    // Keeps track of the merges that we want to perform.  
    Set<PartitionMultiMerge> merges = new LinkedHashSet<PartitionMultiMerge>();  
}
```

$1, \dots, n_1$  s.t.  $seq_j^2 \subseteq seq_i^1$ , and vice versa  $\forall i = 1, \dots, n_1$   
 $\exists j = 1, \dots, n_1$  s.t.  $seq_j^2 \subseteq seq_i^1$

- $f1$  strongly subsumes  $f2$  iff  $\forall j = 1, \dots, n_2$   $\exists i = 1, \dots, n_1$   
s.t.  $seq_j^2 \subseteq seq_i^1$

```
12 ... event1 / event2 ∈ initialGraph;  
13 end while  
14 if (hybrid)  
15  $(V, E) := kTail((V, E), 0, I)$   
16 Output:  $(V, E)$ 
```

## 2. Monolithic architecture

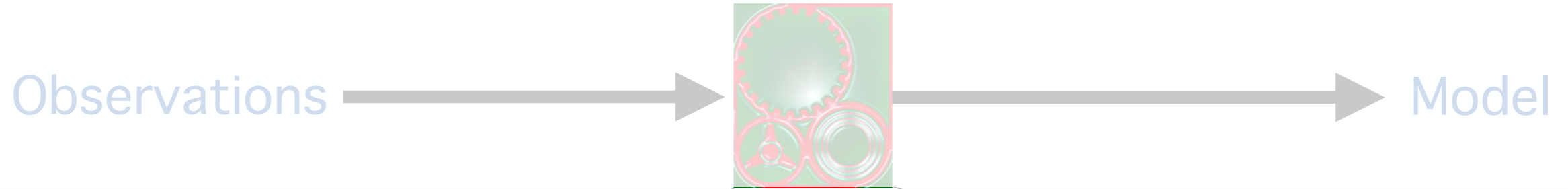
Limitations:

✗ Transparency

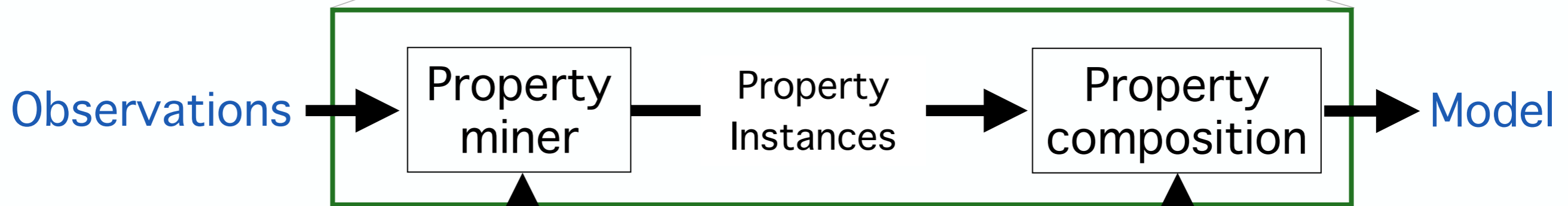
✗ Extensibility

# InvariMint: modular and declarative

Prior work:



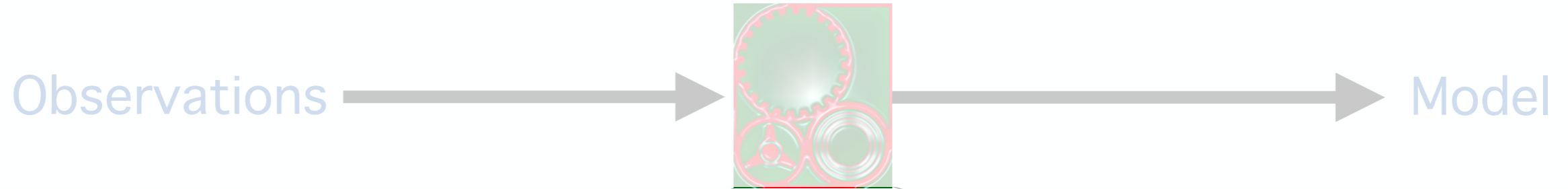
InvariMint:



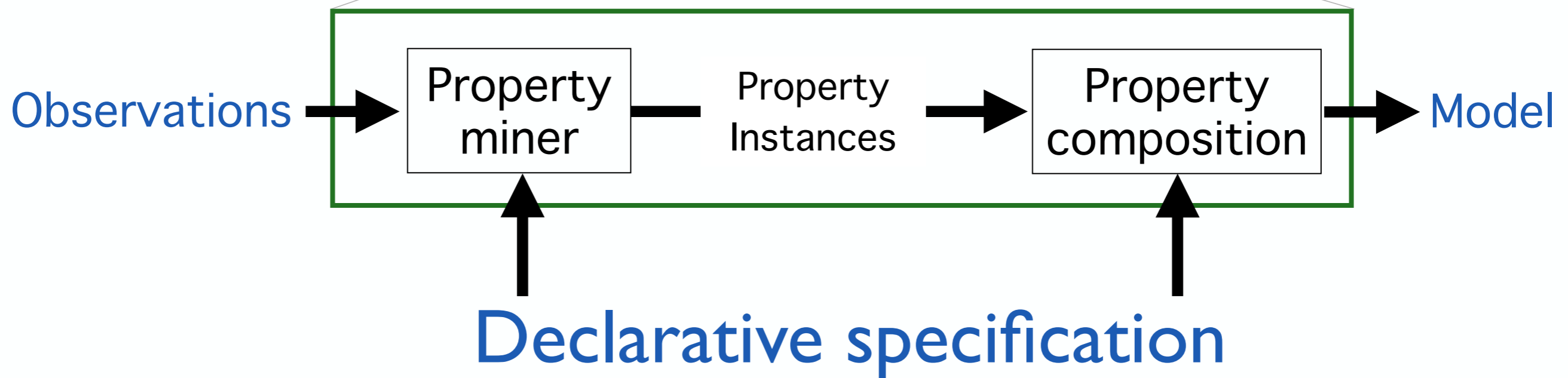
Declarative specification  
of FSM-inference algorithm

# InvariMint: modular and declarative

Prior work:



InvariMint:

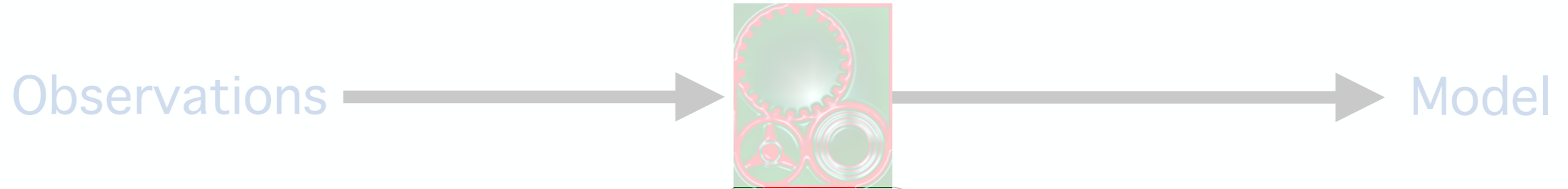


✓ Transparent

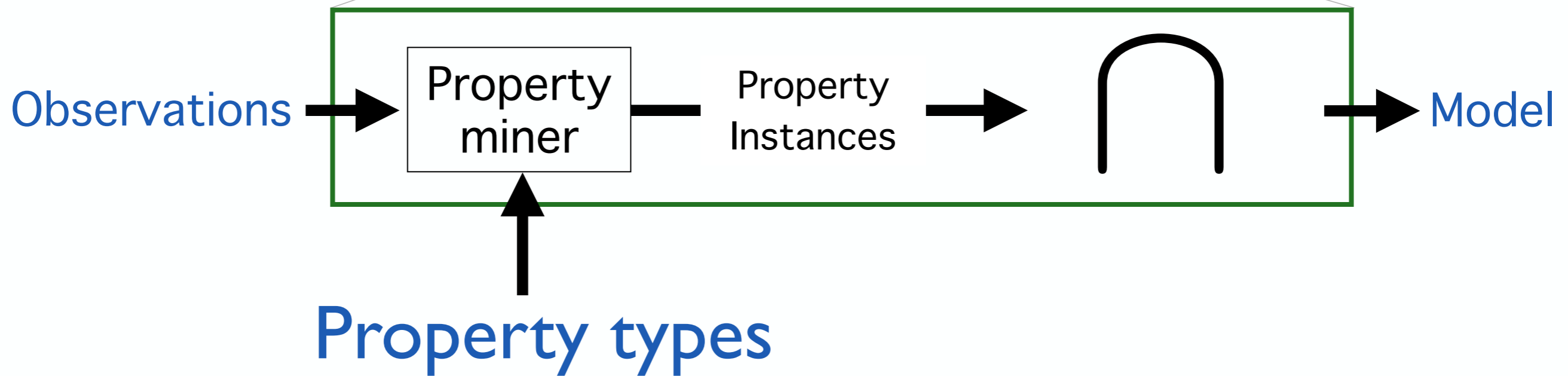
✓ Extensible

# InvariMint: modular and declarative

Prior work:



InvariMint:



✓ Transparent

✓ Extensible

# Outline

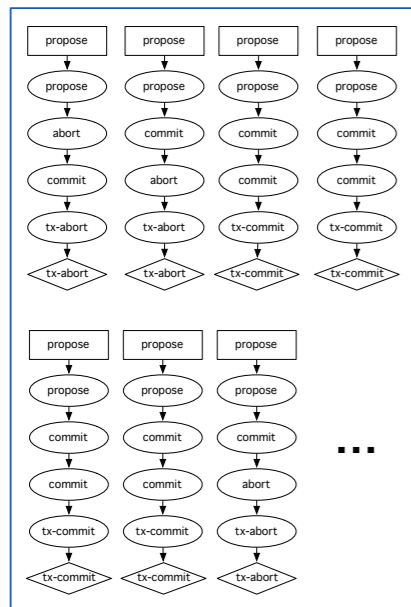
- FSM-inference algorithms not transparent/extensible
  - ▶ **InvariMint** -- modular and declarative FSM-inference
- Expressing algorithm using **InvariMint**
  - ▶ kTails -- a state-merging algorithm
  - ▶ Synoptic -- a state-splitting algorithm
- General specification formalism
- **InvariMint** evaluation



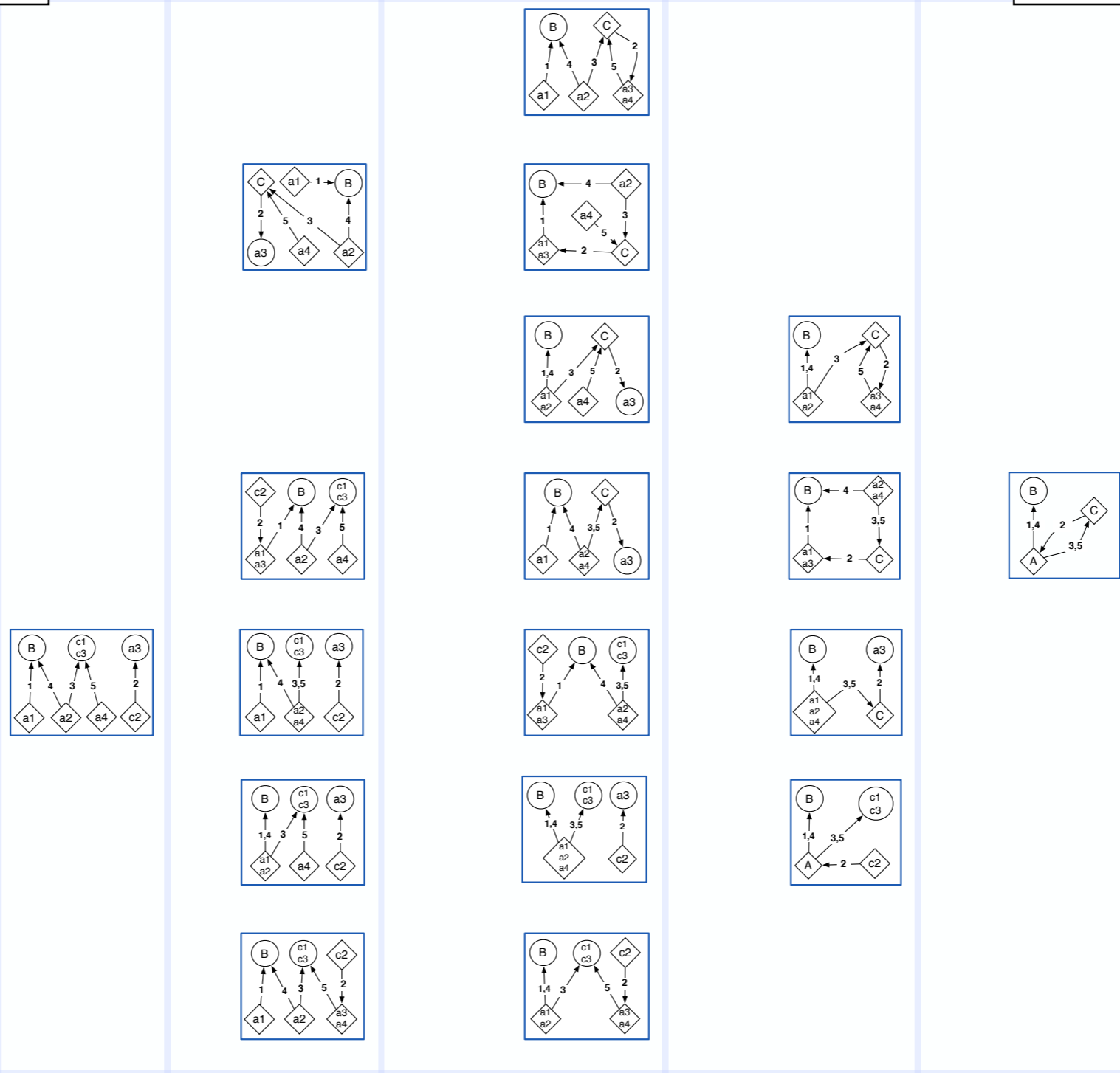
# kTails and Synoptic overview

Larger models:  
fewer behaviors

Smaller models:  
more behaviors



Observations

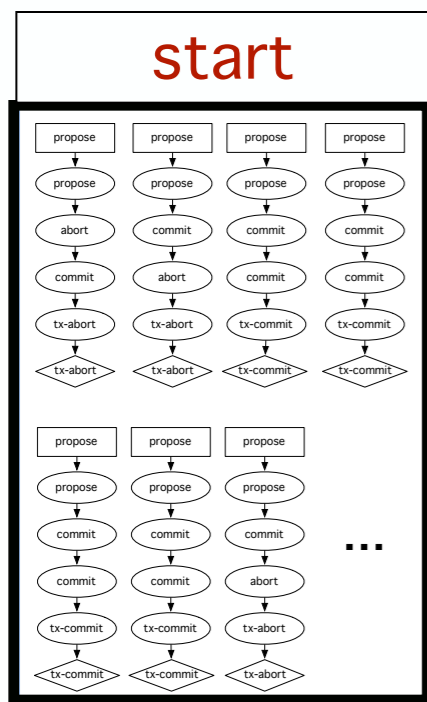


Smallest model

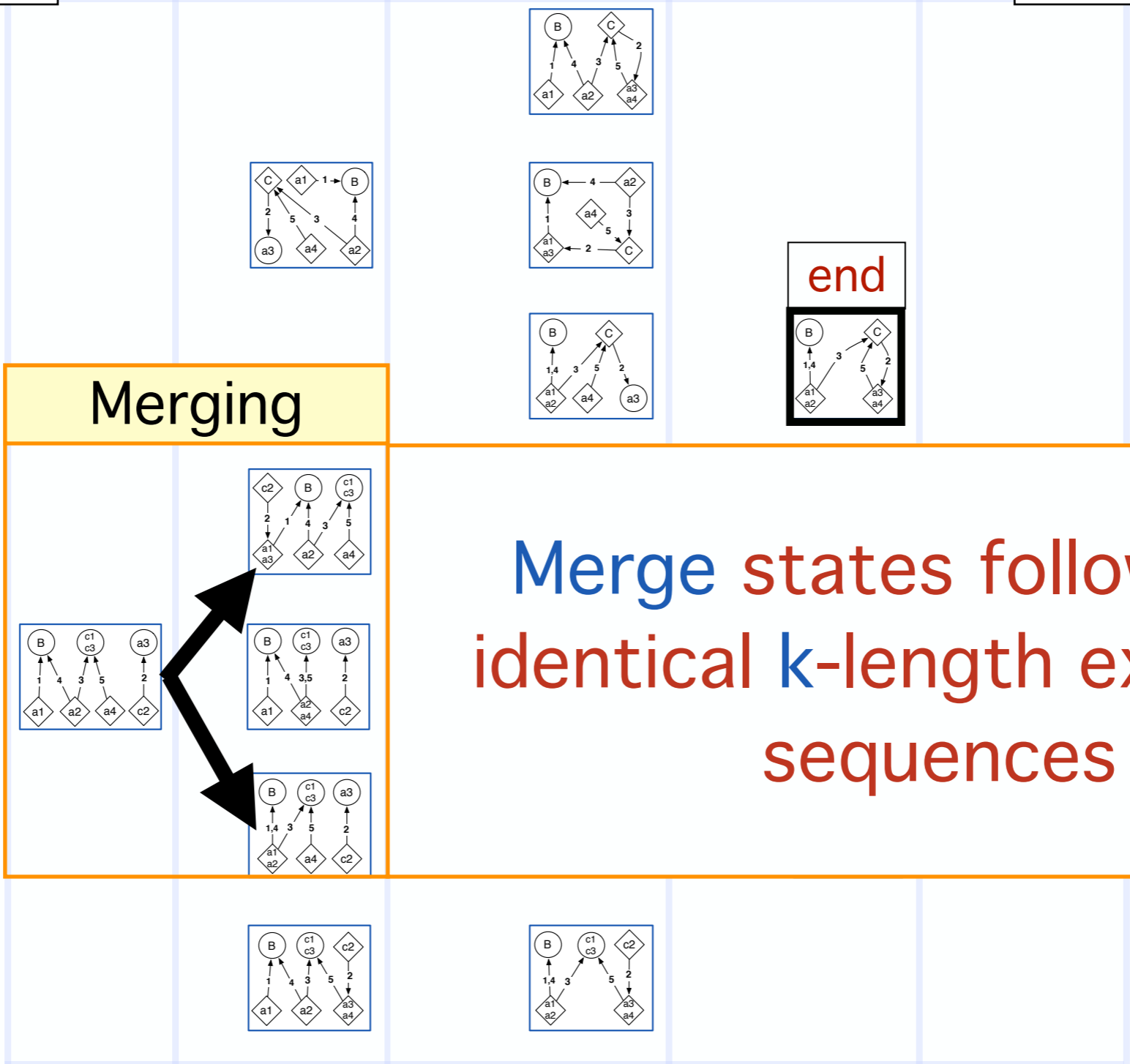
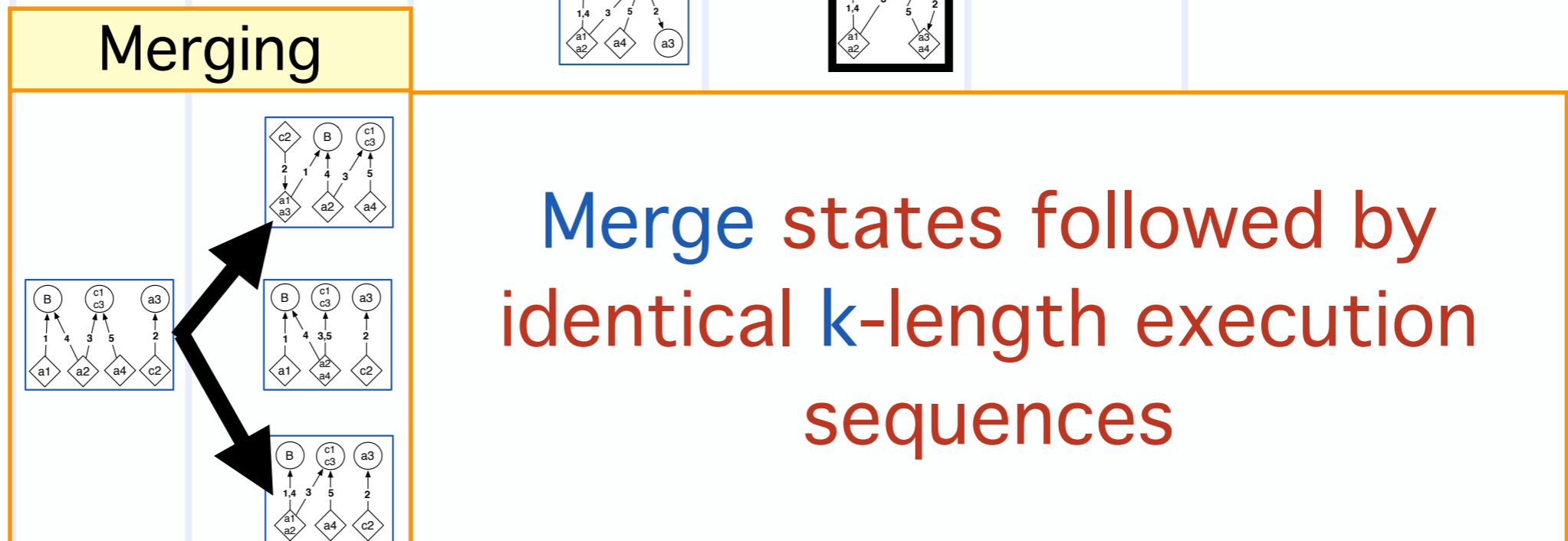
# kTails(k) and Synoptic overview

Larger models:  
fewer behaviors

Smaller models:  
more behaviors



Observations



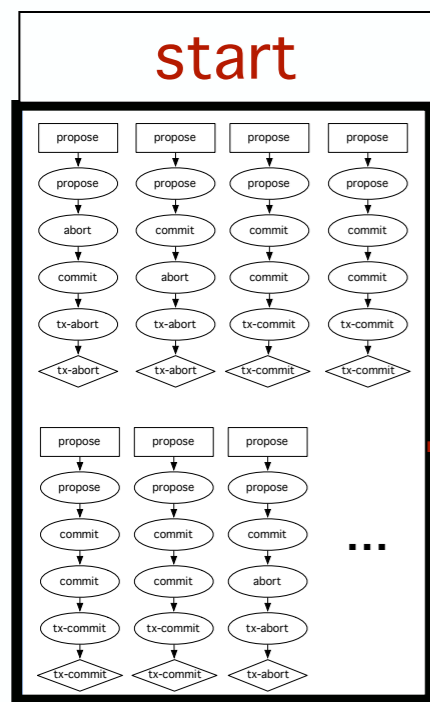


# kTails(k) and Synoptic overview

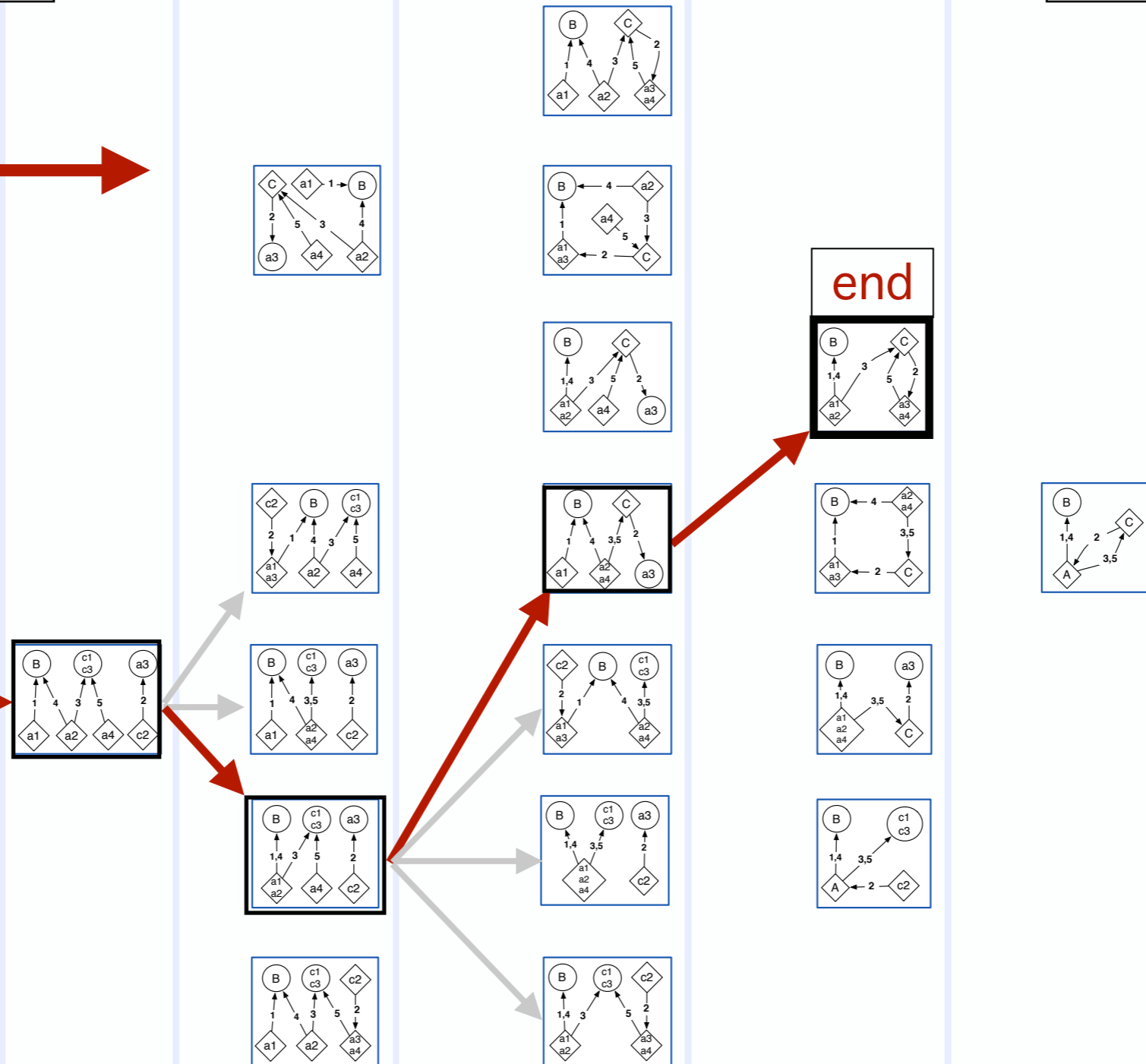
Larger models:  
fewer behaviors

Smaller models:  
more behaviors

Merging



Observations

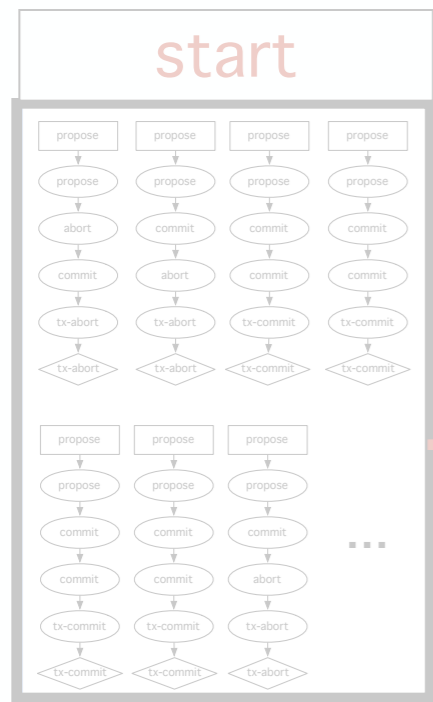


# kTails(k) and Synoptic overview

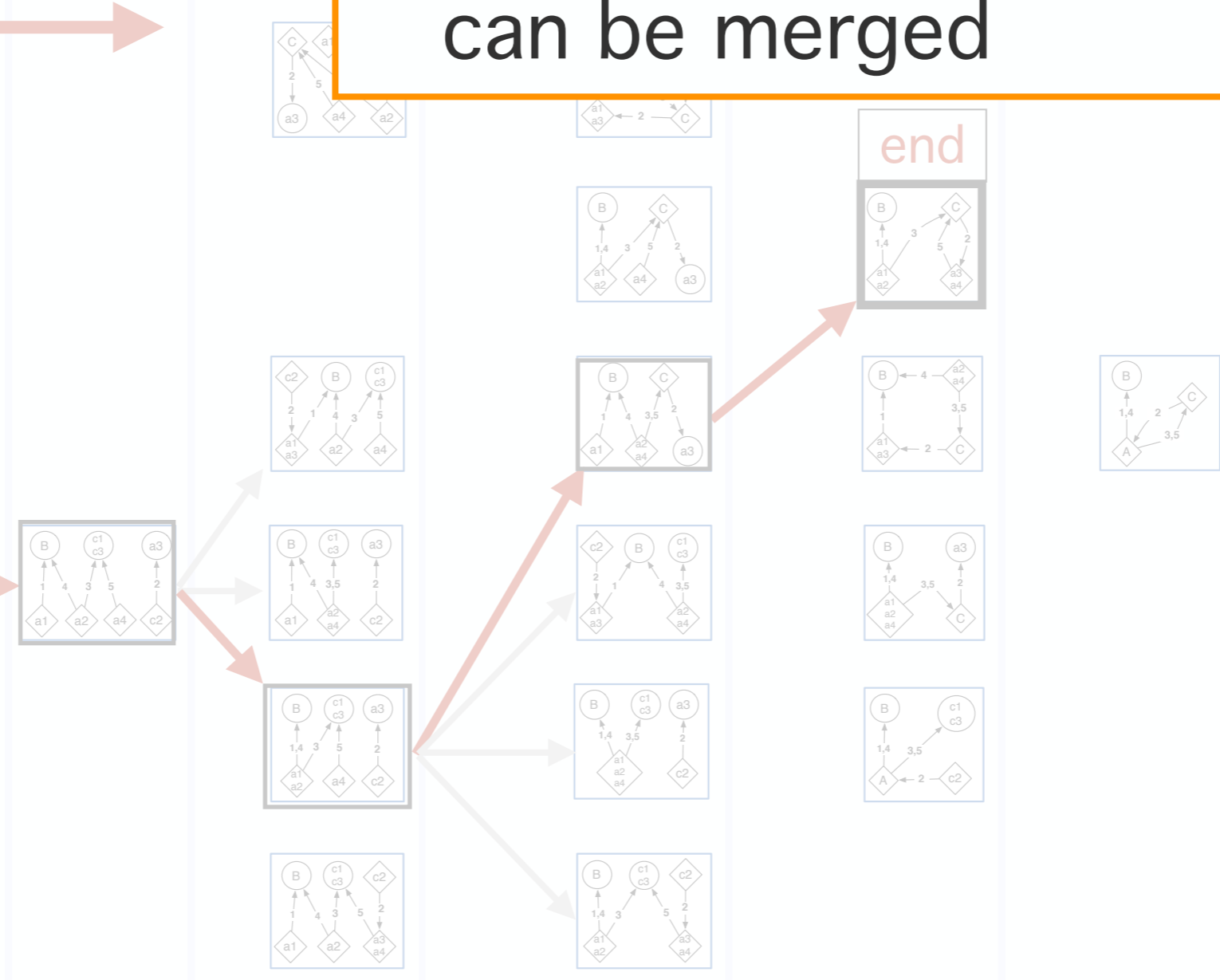
Larger models:  
fewer behaviors

Merging

1. Choose any merge
2. Terminate when no more states can be merged



Observations



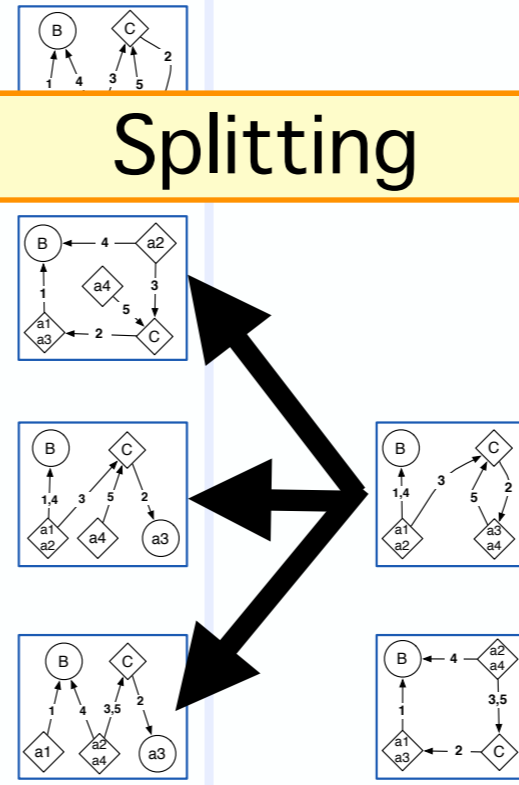
# kTails(k) and Synoptic overview

Larger models:  
fewer behaviors

Smaller models:  
more behaviors

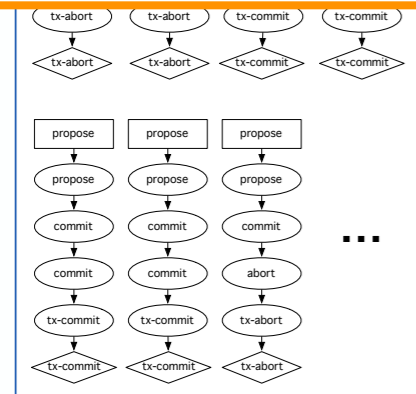
Split states to eliminate executions that violate a set of mined observation invariants

## Splitting

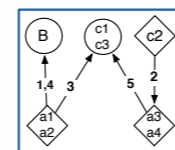
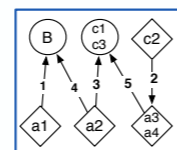
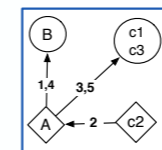
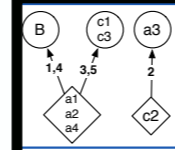
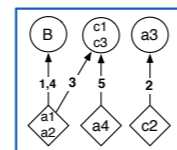
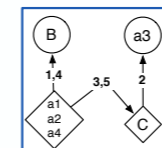
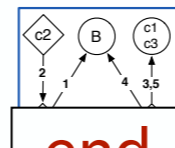
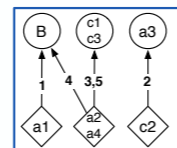
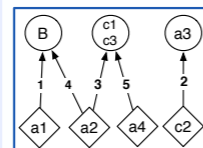


start

Smallest model



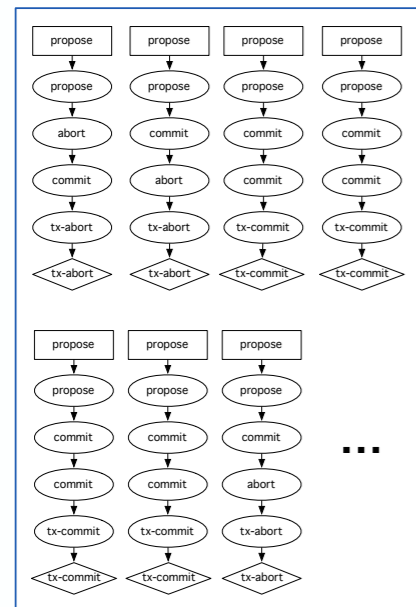
Observations



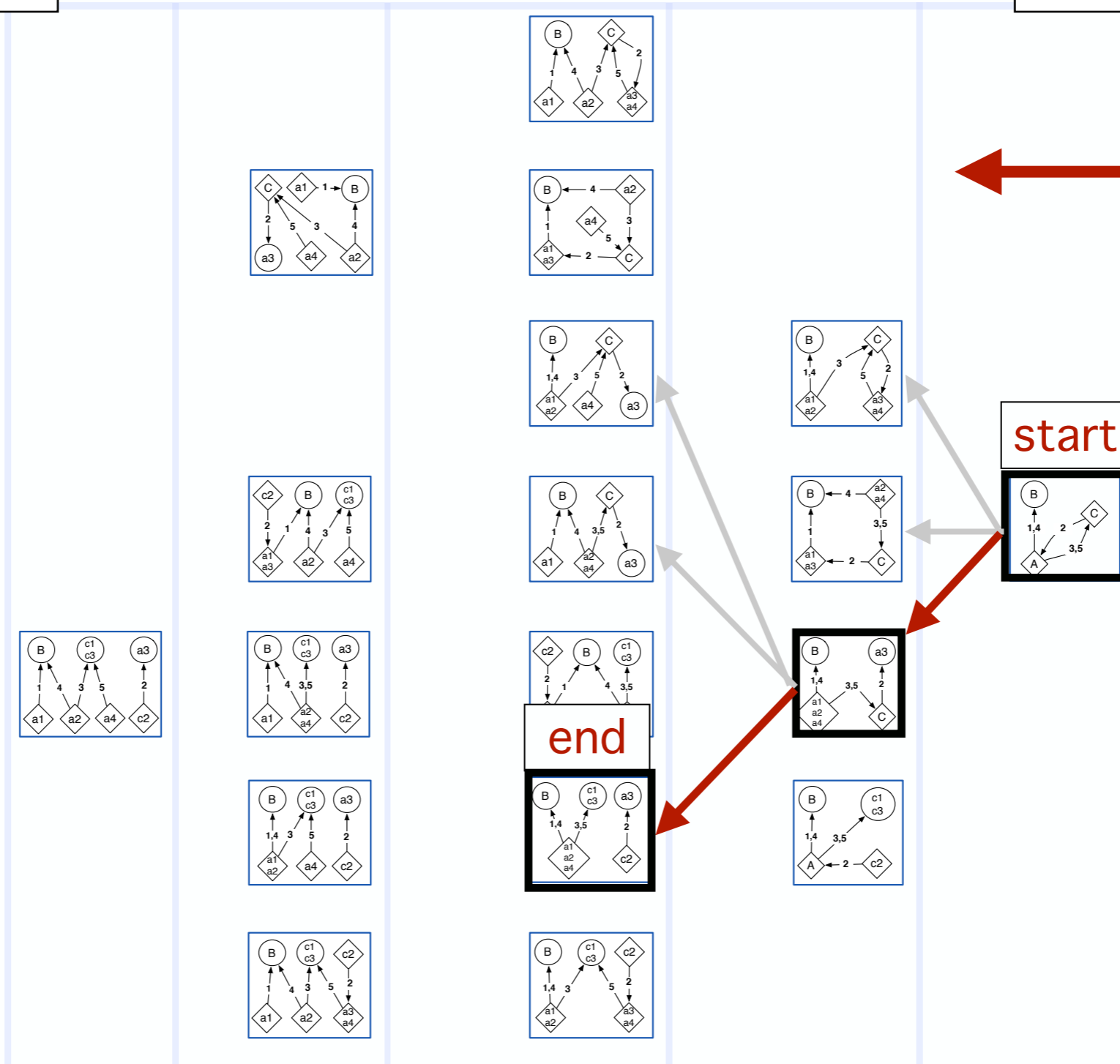
# kTails(k) and Synoptic overview

Larger models:  
fewer behaviors

Smaller models:  
more behaviors



Observations



Splitting

start

Smallest  
model

end

# kTails(k) and **Synoptic** overview

**Synoptic is non-deterministic:**

Final model depends on splitting choices

1. Choose split to eliminate counter-examples to unsatisfied mined invariants:

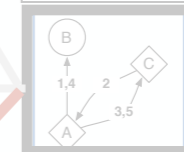
- **x** AlwaysFollowedBy **y**
- **x** NeverFollowedBy **y**
- **x** AlwaysPrecedes **y**

2. Terminate when model satisfies all mined invariants

Smaller models:  
more behaviors

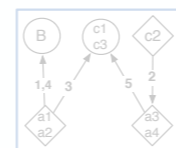
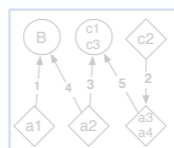
Splitting

start



Smallest  
model

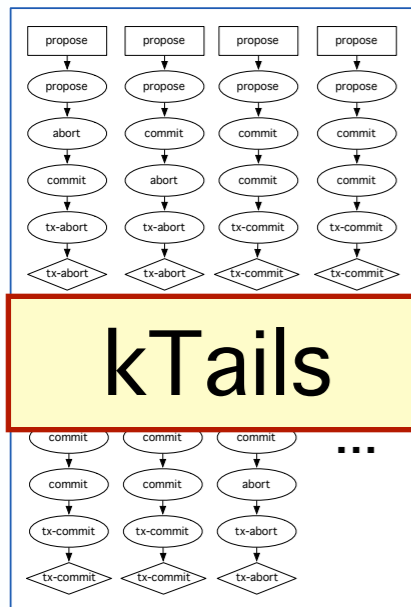
Observations



# Motivating questions

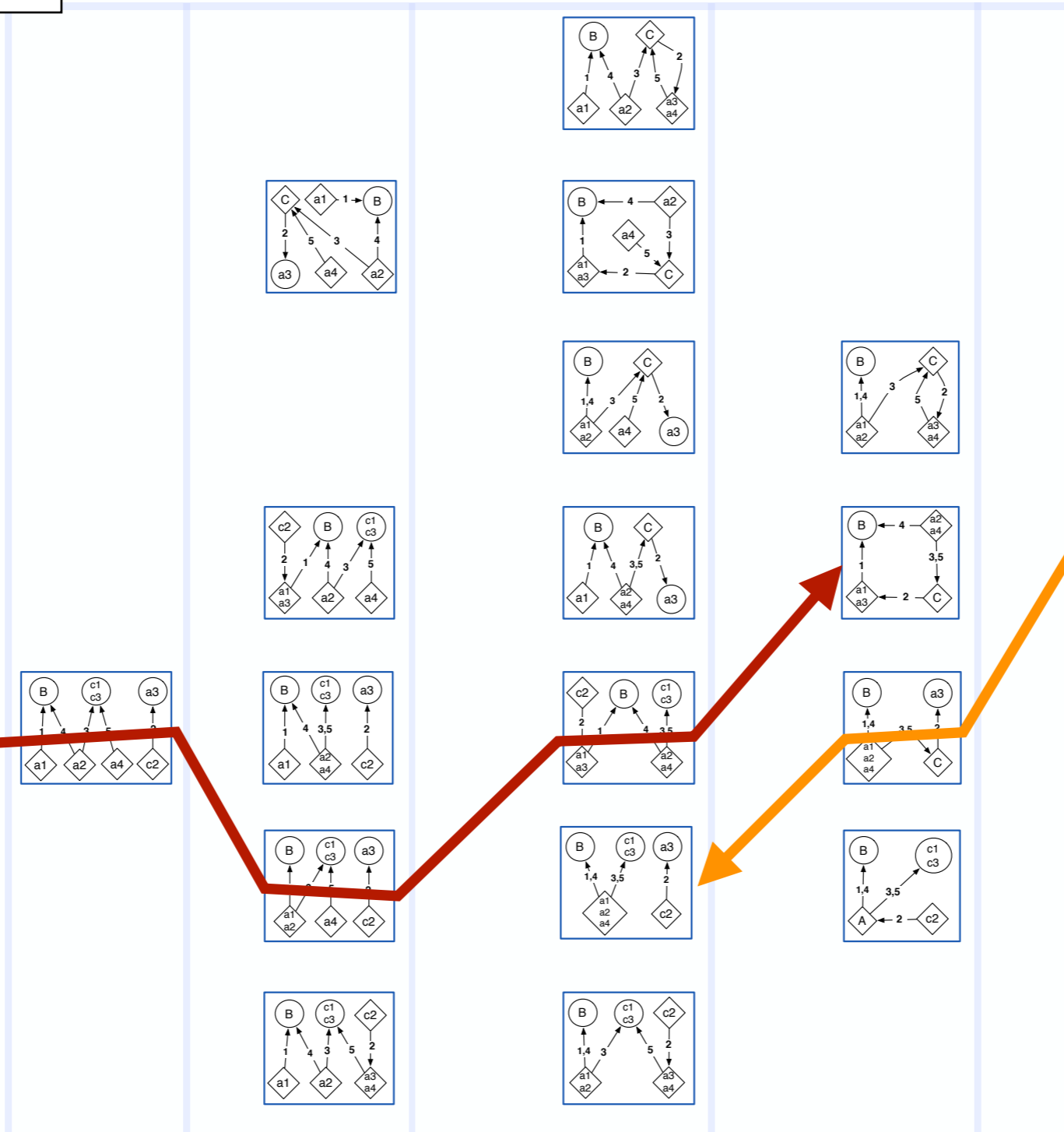
Larger models:  
fewer behaviors

Smaller models:  
more behaviors



Observations

kTails



Synoptic

# Motivating questions

Larger models:  
fewer behaviors

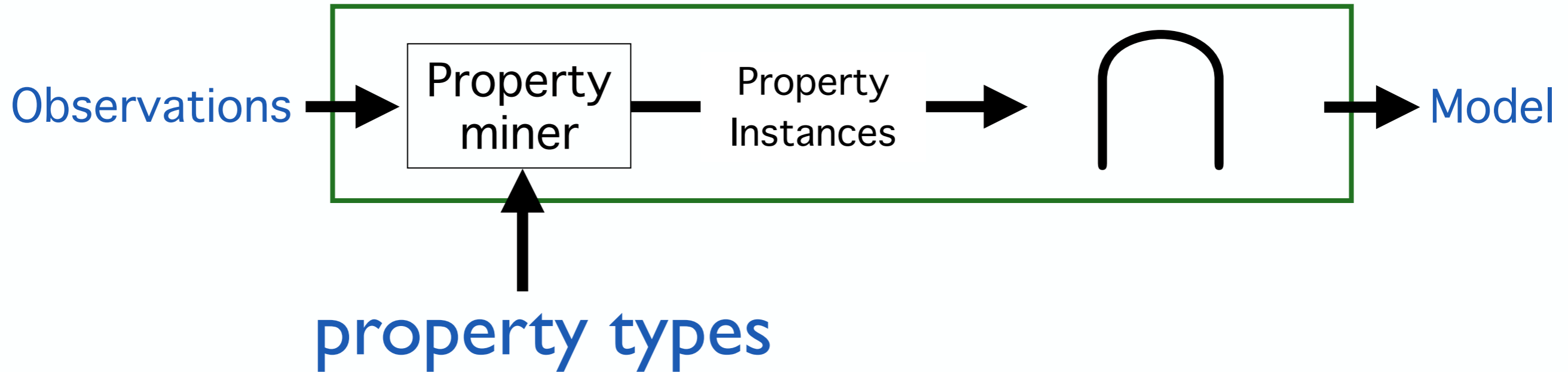
Smaller models:  
more behaviors

How can we easily:

- ... get kTails to ignore certain  $k$ -length sequences?
- ... add the  $x$  AlwaysFollowedBy  $y$  invariant to kTails?
- ... make Synoptic deterministic?
- ... add a new kind of invariant to Synoptic?
- ... learn which properties kTails/Synoptic preserve?

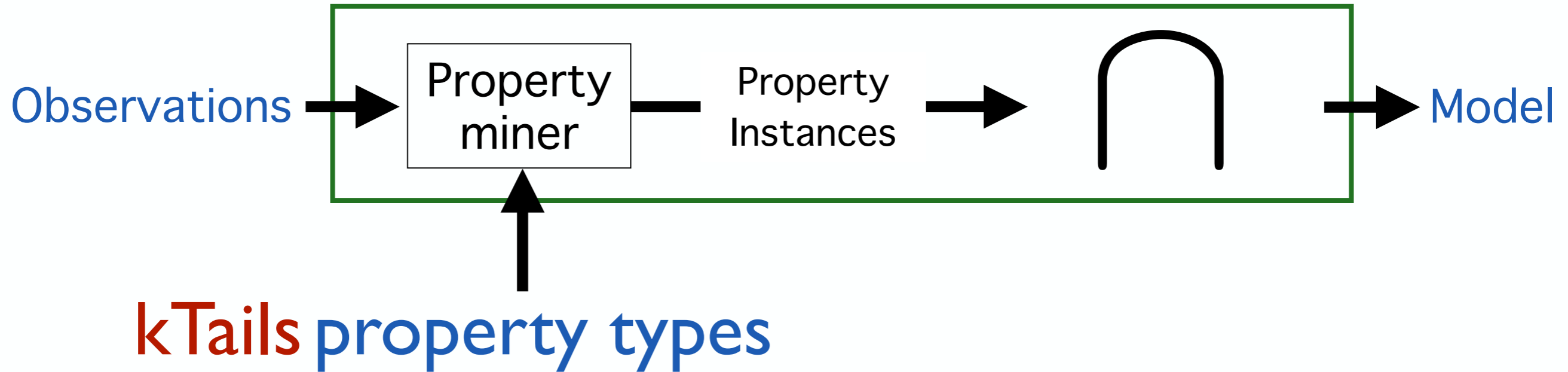
We can answer all of these questions by representing these algorithms with **InvariMint**

# InvariMint: modular and declarative

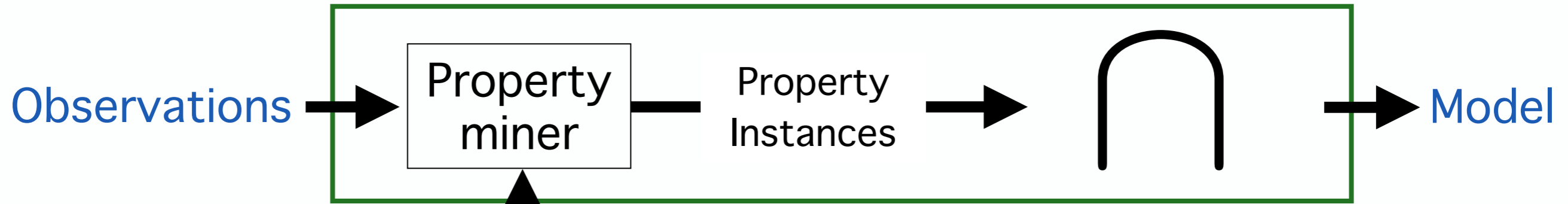




# Expressing **kTails**( $k=1$ )



# Expressing $k$ Tails ( $k=1$ )



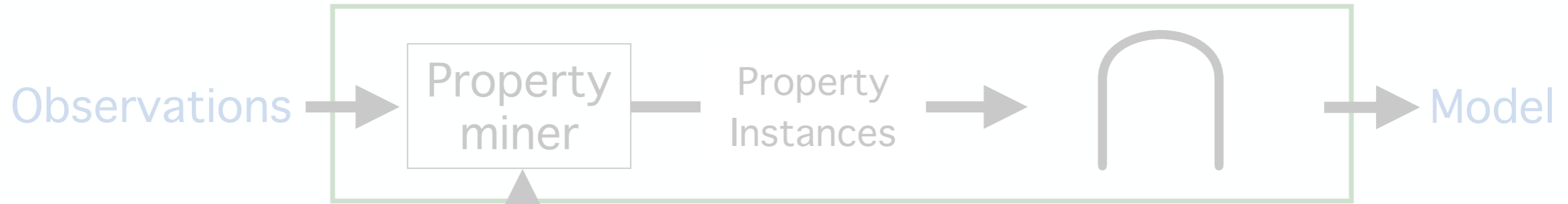
A template to express merging of observation sequences that are identical in the first **1** event(s)

=

A property type:

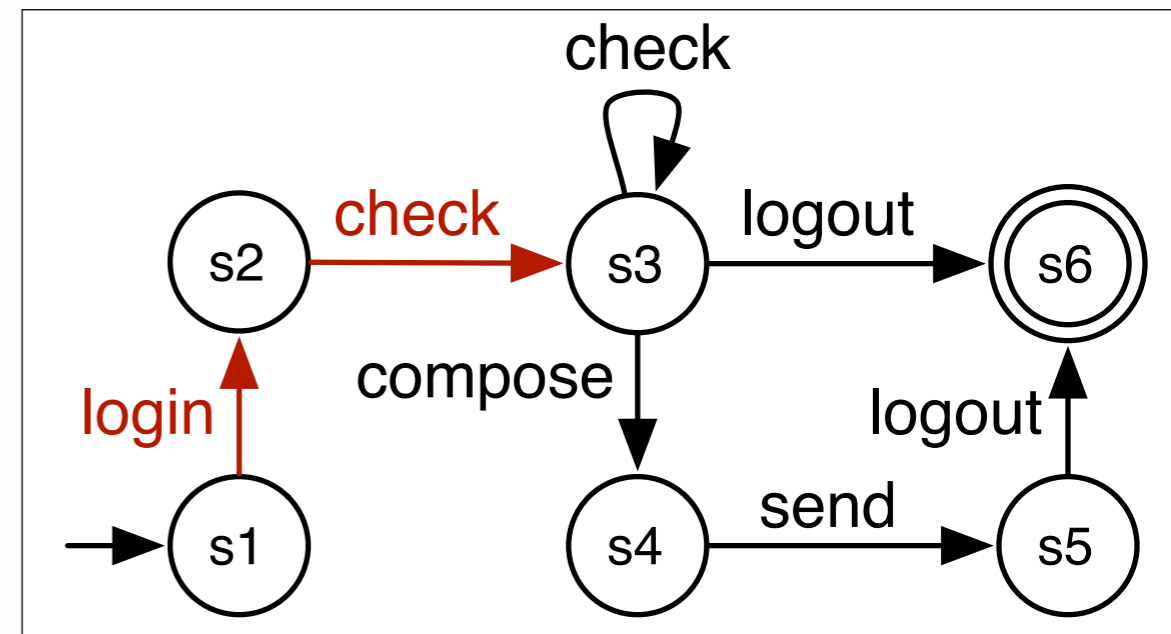
“**x** can be immediately followed by one of **Y**”

# Expressing $k$ Tails ( $k=1$ )



“ $x$  can be immediately followed by one of  $Y$ ”

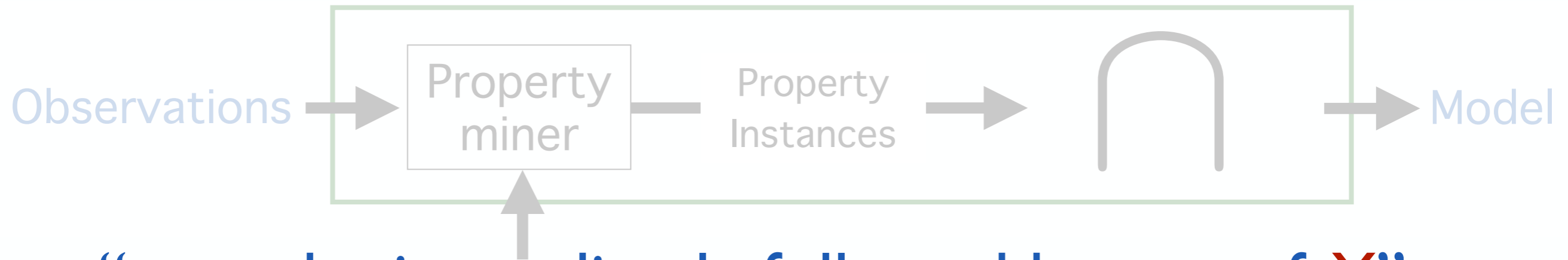
trace 1:	trace 2:
login check check logout	login check compose send logout



Observations

$k$ Tails output

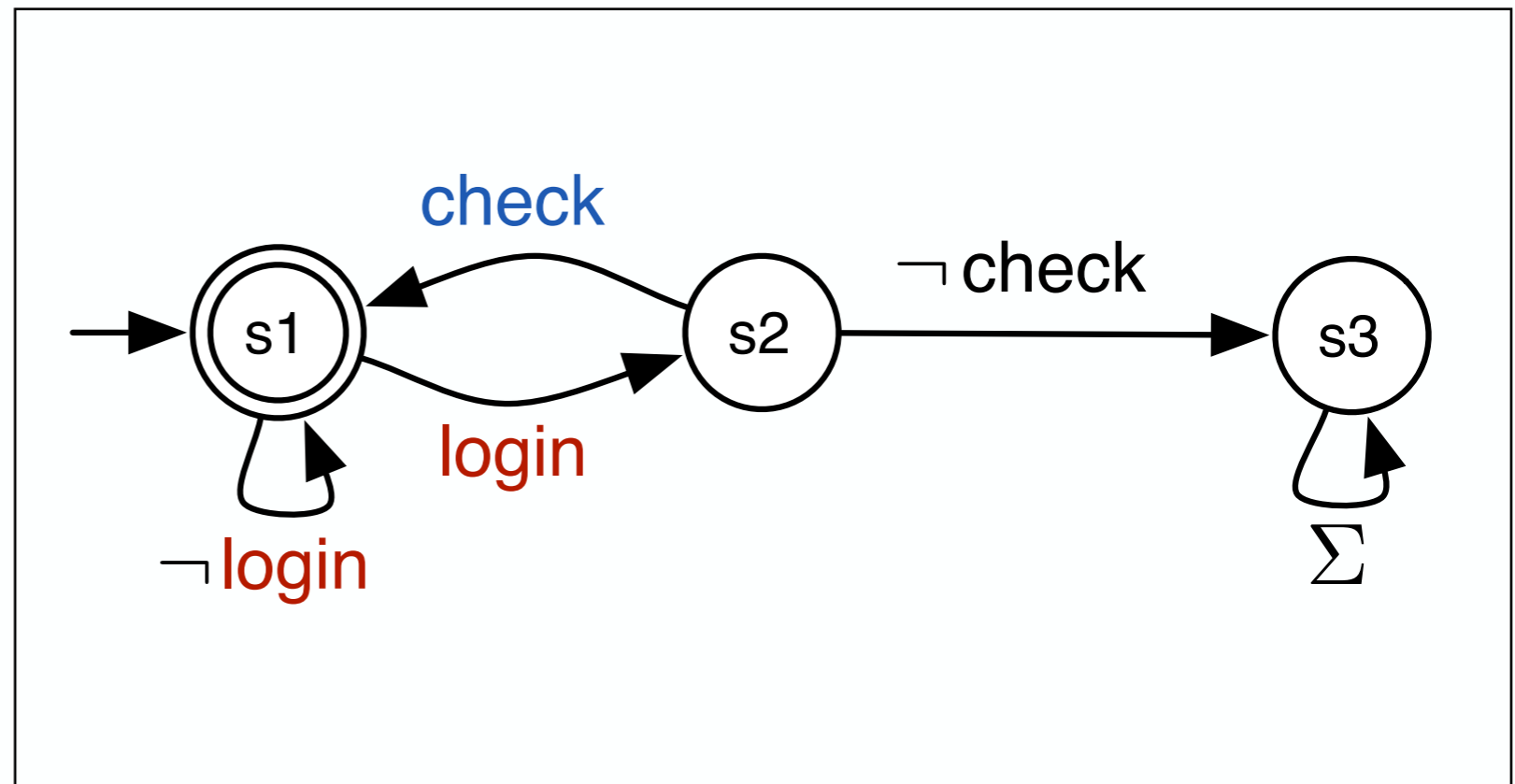
# Expressing $k$ Tails ( $k=1$ )



“ $x$  can be immediately followed by one of  $Y$ ”

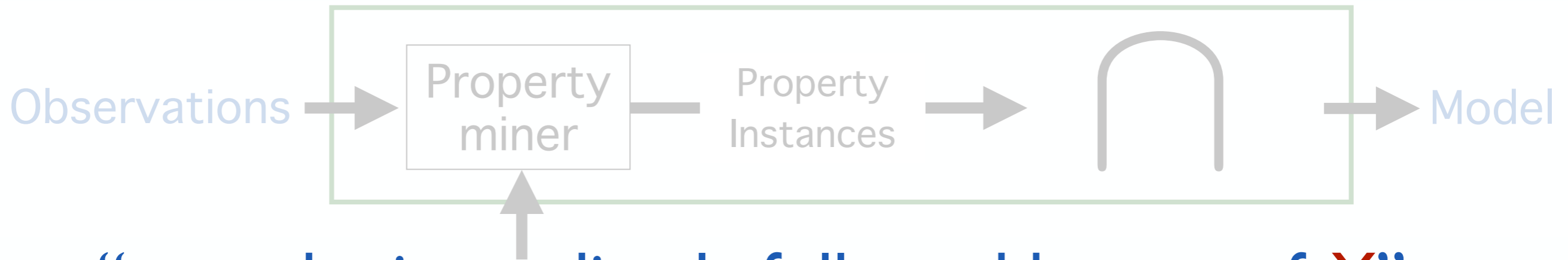
trace 1:	trace 2:
login	login
check	check
check	compose
logout	send
	logout

Observations



$x$ =login property instance

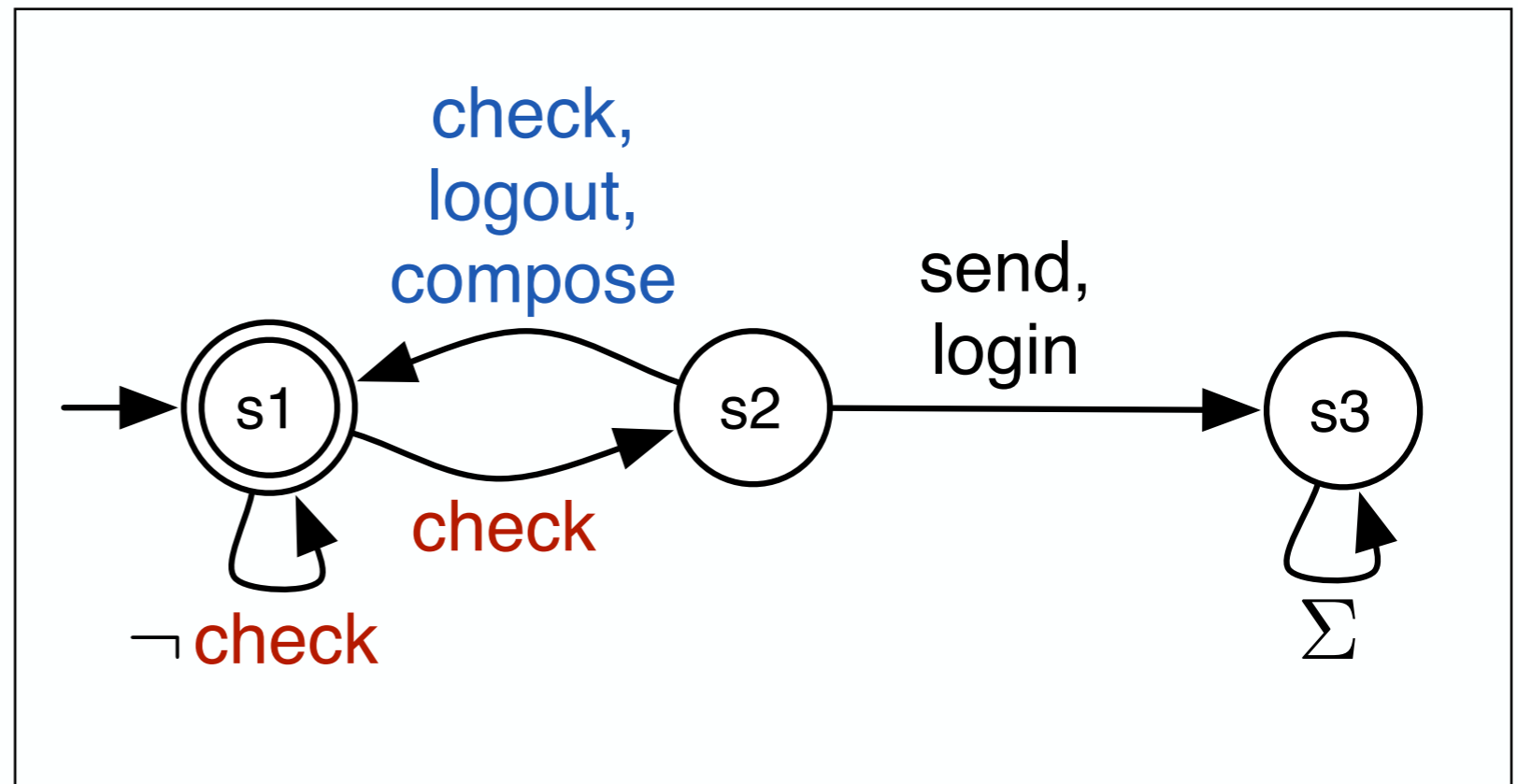
# Expressing $k$ Tails ( $k=1$ )



“ $x$  can be immediately followed by one of  $Y$ ”

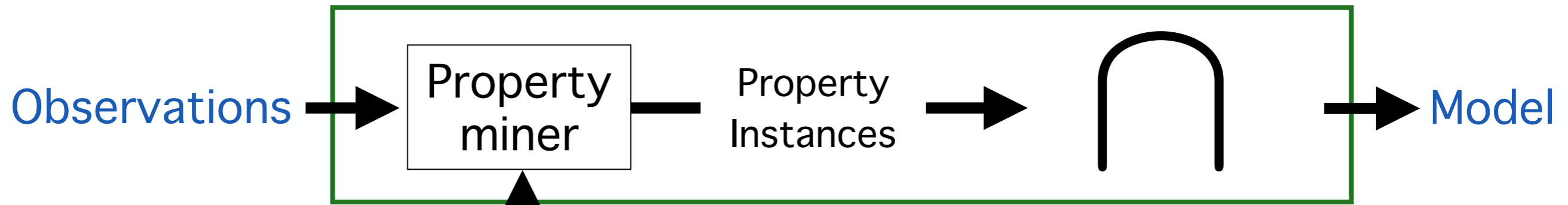
trace 1:	trace 2:
login check check logout	login check compose send logout

Observations



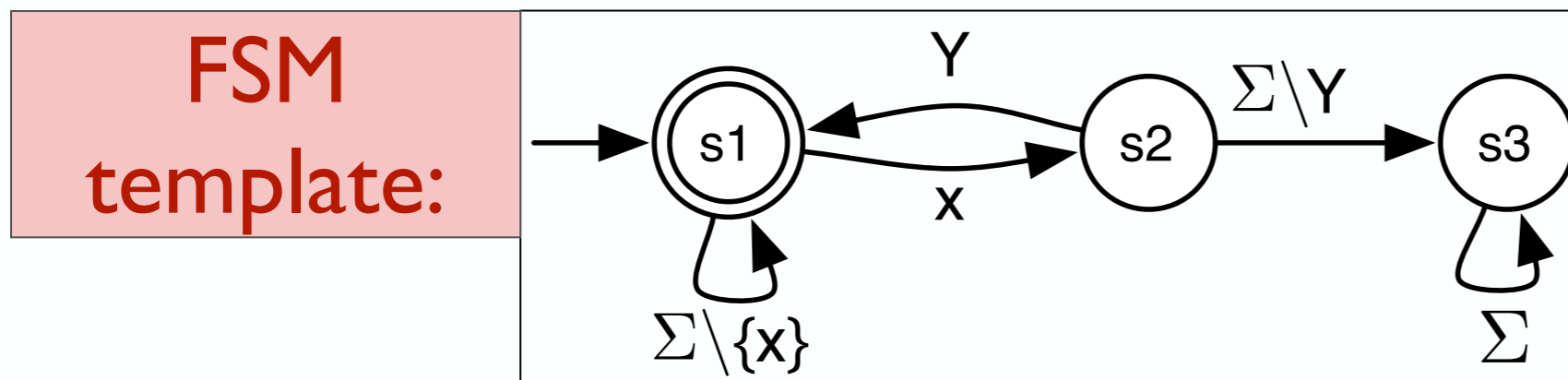
$x = \text{check}$  property instance

# Expressing $k$ Tails ( $k=1$ )



“ $x$  can be immediately followed by one of  $Y$ ”

=

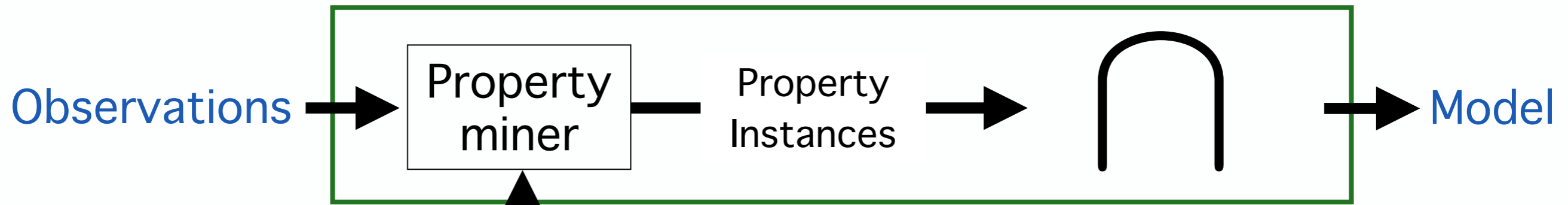


+

**When to instantiate:**

$$Eval(\text{Log } L, \langle x=a, Y=B \rangle) = \begin{cases} \text{true} : \forall t \in L, \exists b \in B, \diamond(a \rightarrow \bigcirc b) \text{ in } t \wedge \\ \quad \forall b \in B, \exists t \in L, \diamond(a \rightarrow \bigcirc b) \text{ in } t \\ \text{false} : \text{otherwise} \end{cases}$$

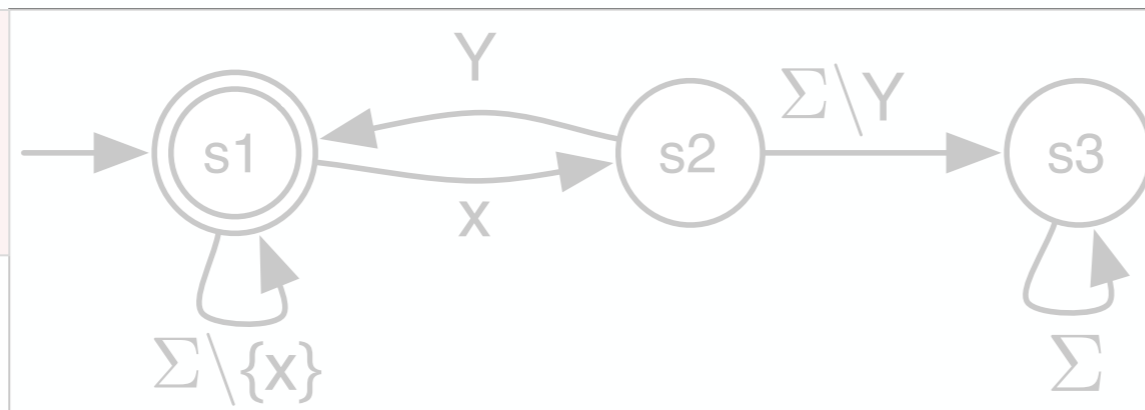
# Expressing **kTails**( $k=1$ )



“**x** can be immediately followed by one of **Y**”

=

FSM  
template:



Complete  
template

+

When to  
instantiate:

$$Eval(\text{Log } L, \langle x=a, Y=B \rangle) = \begin{cases} \text{true} : \forall t \in L, \exists b \in B, \diamond(a \rightarrow \bigcirc b) \text{ in } t \wedge \\ \quad \forall b \in B, \exists t \in L, \diamond(a \rightarrow \bigcirc b) \text{ in } t \\ \text{false} : \text{otherwise} \end{cases}$$

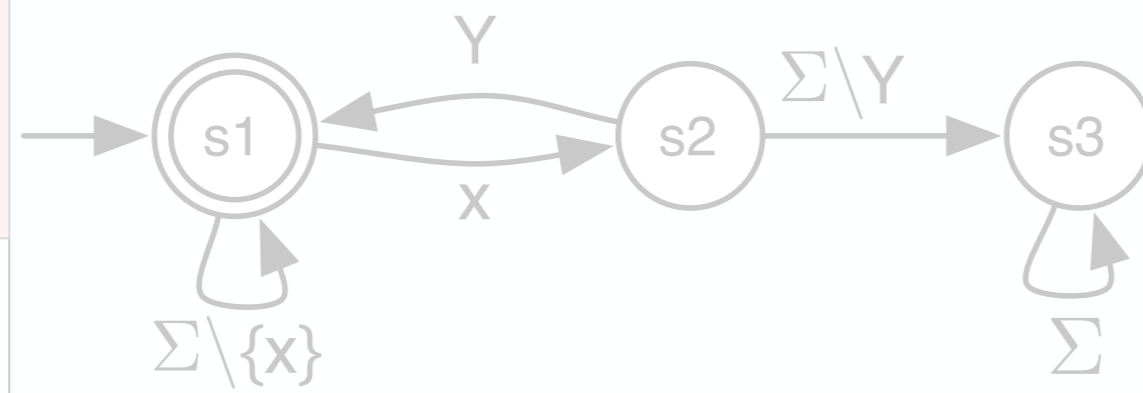
# Expressing $k$ Tails ( $k=1$ )

- This formulation is exact (identical models)
- Generalizes to arbitrary  $k$  value

“ $x$  can be immediately followed by one of  $Y$ ”

=

FSM  
template:



Complete  
template

+

When to  
instantiate:

$$Eval(\text{Log } L, \langle x=a, Y=B \rangle) = \begin{cases} \text{true} : \forall t \in L, \exists b \in B, \diamond(a \rightarrow \bigcirc b) \text{ in } t \wedge \\ \quad \forall b \in B, \exists t \in L, \diamond(a \rightarrow \bigcirc b) \text{ in } t \\ \text{false} : \text{otherwise} \end{cases}$$



```

package synoptic.algorithms;

import java.util.ArrayList;
import java.util.Collections;
import java.util.LinkedHashMap;
import java.util.LinkedHashSet;
import java.util.List;
import java.util.Map;
import java.util.Set;
import java.util.logging.Logger;

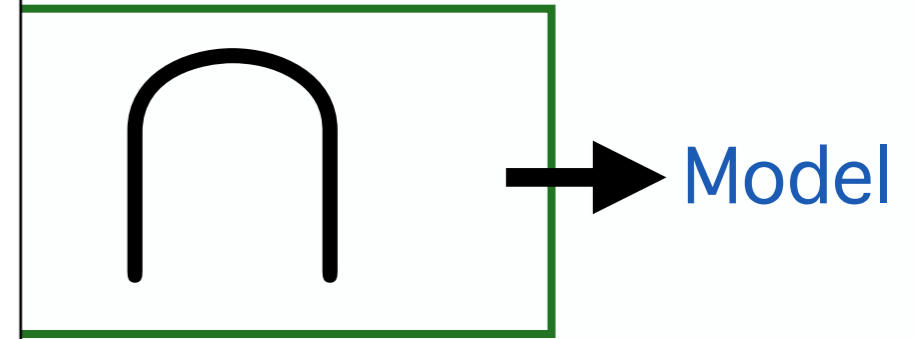
import synoptic.algorithms.graphops.PartitionMultiMerge;
import synoptic.model.ChainsTraceGraph;
import synoptic.model.Partition;
import synoptic.model.PartitionGraph;
import synoptic.model.event.EventType;
import synoptic.model.interfaces.INode;
import synoptic.model.interfaces.ITransition;
import synoptic.util.InternalSynopticException;
import synoptic.util.NotImplementedException;

/**
 * Implements the KTails algorithm as defined in Biermann & Feldman '72.
 */
public class KTails {
    public static Logger logger;
    static {
        logger = Logger.getLogger("KTails");
    }

    /**
     * Constructs and returns a PartitionGraph generated by applying kTails with
     * the given k value to the given trace graph
     */
    public static PartitionGraph performKTails(ChainsTraceGraph g, int k) {
        PartitionGraph pGraph = new PartitionGraph(g, false, null);
        attemptMerge(pGraph, k);
        return pGraph;
    }

    /**
     * Finds all possible merges in pGraph. Requires making a new call to
     * attemptMerge after every merge in case previously unmergeable pairs

```



by one of Y”

Complete  
template

$\in L, \exists b \in B, \diamond(a \rightarrow \bigcirc b)$  in  $t \wedge$   
 $\in B, \exists t \in L, \diamond(a \rightarrow \bigcirc b)$  in  $t$   
 otherwise

# kTails Procedural description

```

    allVisitedMatches.put(n1, n2);
    for (ITransition<NodeType> t1 : n1Trans) {
        NodeType c1 = t1.getTarget();
        // Skip c1 if it was visited by this method earlier.
        if (visitedN1Children.contains(c1)) {
            continue;
        }

        boolean kEqual = false;

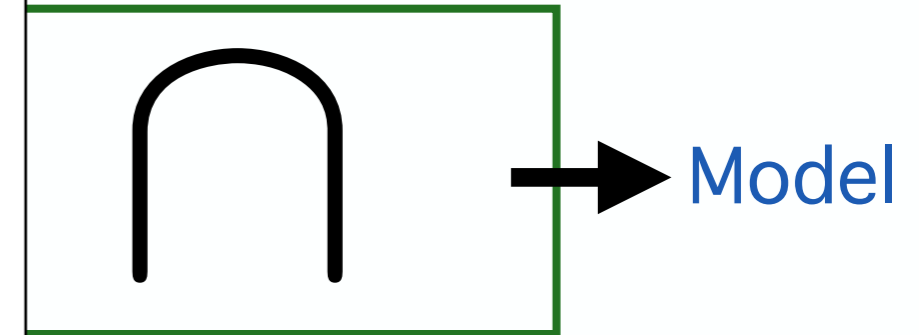
        // Make sure to get transitions of the same relation.
        for (ITransition<NodeType> t2 : n2
            .getTransitionsWithExactRelations(t1.getRelation())) {
            NodeType c2 = t2.getTarget();

            // Skip c2 if it was visited by this method earlier.
            if (visitedN2Children.contains(c2)) {
                continue;
            }

            // Skip c2 if its already been mapped to a c1 previously in the
            // outer loop.
            if (childKEquivMatches.contains(c2)) {
                continue;
            }

            if (kEqualsWithoutSubsumption(c1, c2, k - 1, allVisitedMatches)) {
                kEqual = true;
                childKEquivMatches.add(c2);
                break;
            }
        }
        // Could not find any kEqual c2 to match with c1.
        if (!kEqual) {
            // Remove the record of visiting n1 and n2.
            allVisitedMatches.remove(n1);
            return false;
        }
    }
    return true;
}
}
}

```



by one of Y”

Complete  
template

$\in L, \exists b \in B, \diamond(a \rightarrow \bigcirc b)$  in  $t \wedge$   
 $\in B, \exists t \in L, \diamond(a \rightarrow \bigcirc b)$  in  $t$   
 otherwise

# Expressing **kTails(k)**

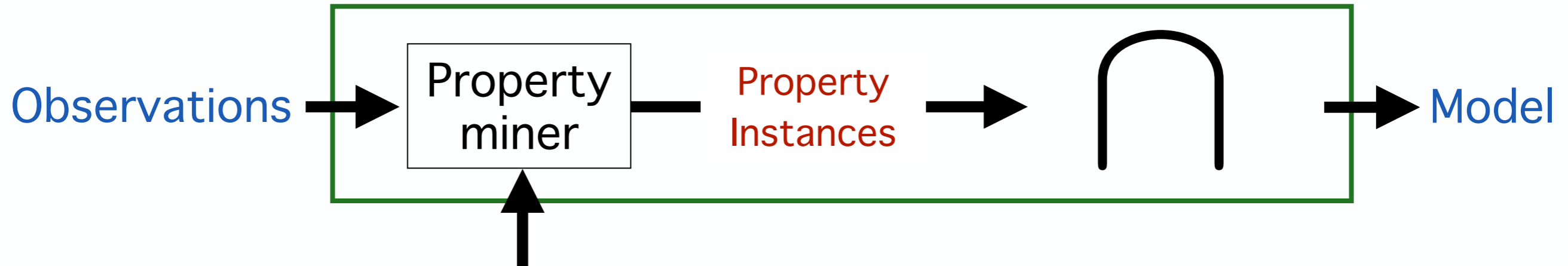
- This formulation is exact (identical models)
- Generalizes to arbitrary  $k$  value

“ $x$  can be immediately followed by one of  $Y$ ”

How can we:

- ... get **kTails** to ignore certain  $k$ -length sequences?
- ... add the  $x$  **AlwaysFollowedBy**  $y$  invariant to **kTails**?
- ... make **Synoptic** deterministic?
- ... add a new kind of invariant to **Synoptic**?
- ... learn which properties **kTails/Synoptic** preserve?

# Expressing kTails(k)



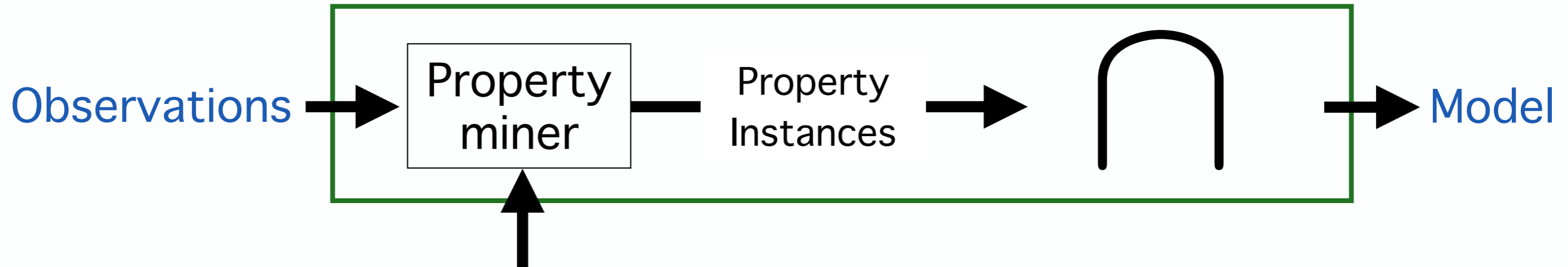
“x can be immediately followed by one of Y”

How can we:

- ... get kTails to ignore certain k-length sequences?
- ... add the x AlwaysFollowedBy y invariant to kTails?
- ... make Synoptic deterministic?
- ... add a new kind of invariant to Synoptic?
- ... learn which properties kTails/Synoptic preserve?

Ev

# Expressing kTails(k)



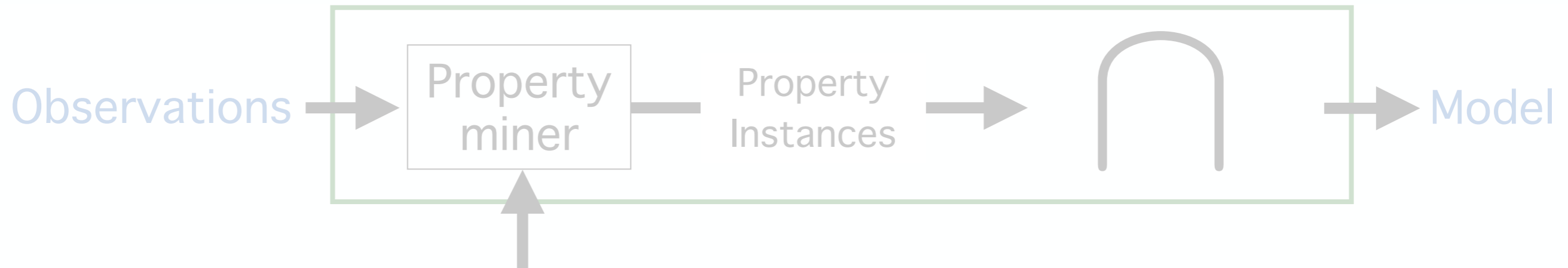
“x can be immediately followed by one of Y”

How can we:

- ... get kTails to ignore certain k-length sequences? ✓
- ... add the x AlwaysFollowedBy y invariant to kTails?
- ... make Synoptic deterministic?
- ... add a new kind of invariant to Synoptic?
- ... learn which properties kTails/Synoptic preserve?

Ev

# Expressing kTails(k)



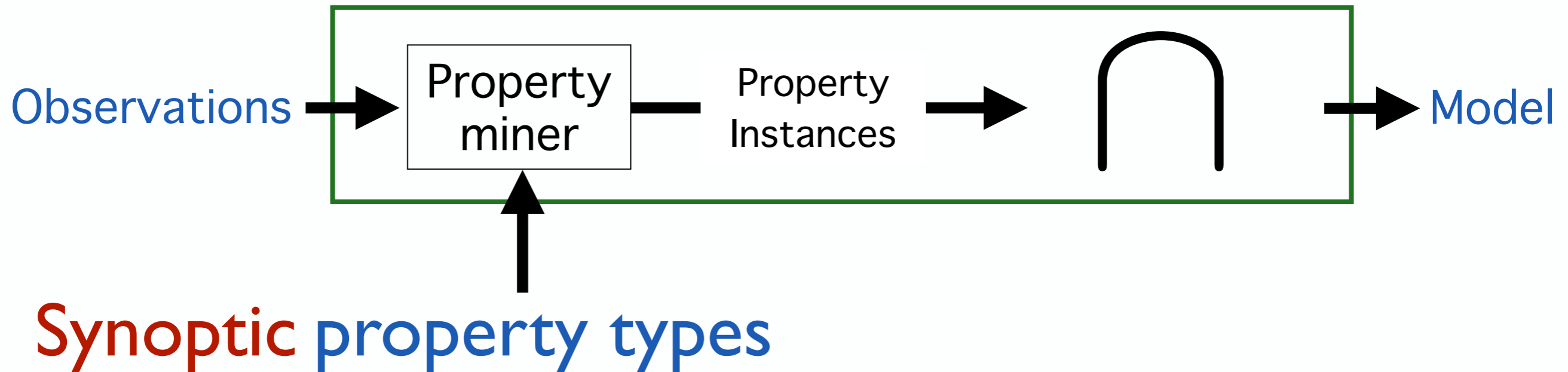
“x can be immediately followed by one of Y”

How can we:

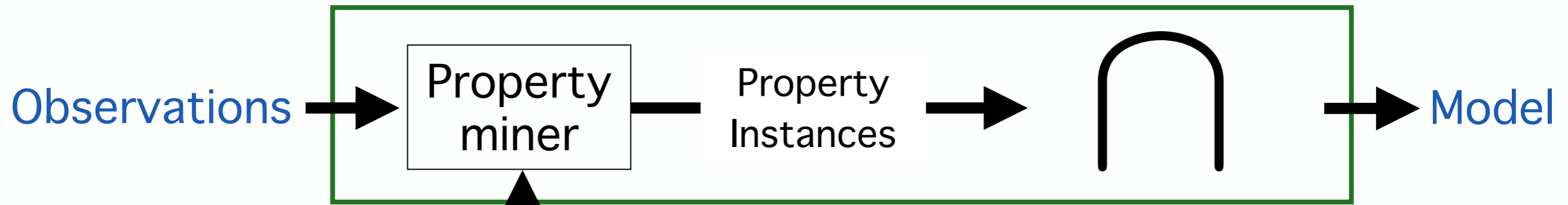
- ... get kTails to ignore certain k-length sequences? ✓
- ... add the x AlwaysFollowedBy y invariant to kTails?
- ... make Synoptic deterministic?
- ... add a new kind of invariant to Synoptic?
- ... learn which properties kTails/Synoptic preserve?

Ev

# Expressing **Synoptic** with **InvariMint**



# Expressing **Synoptic** with InvariMint

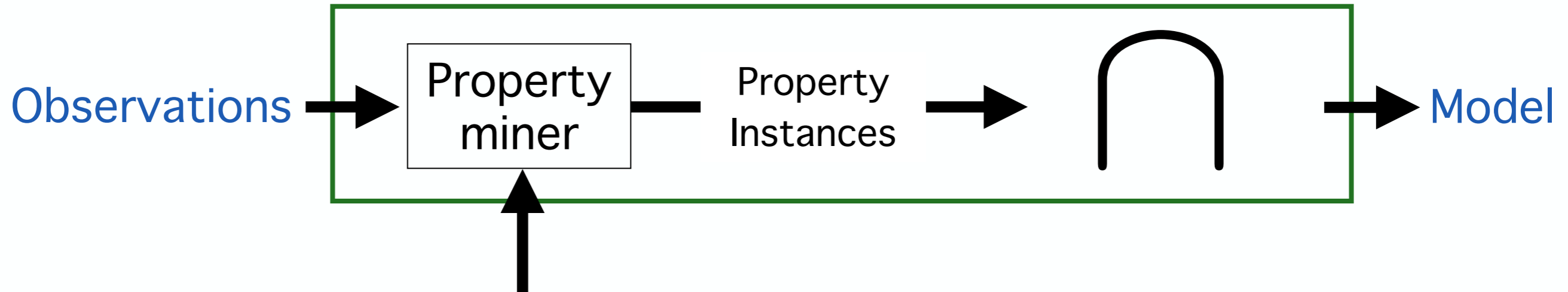


- x AlwaysFollowedBy *y*
- x AlwaysPrecedes *y*
- x NeverFollowedBy *y*

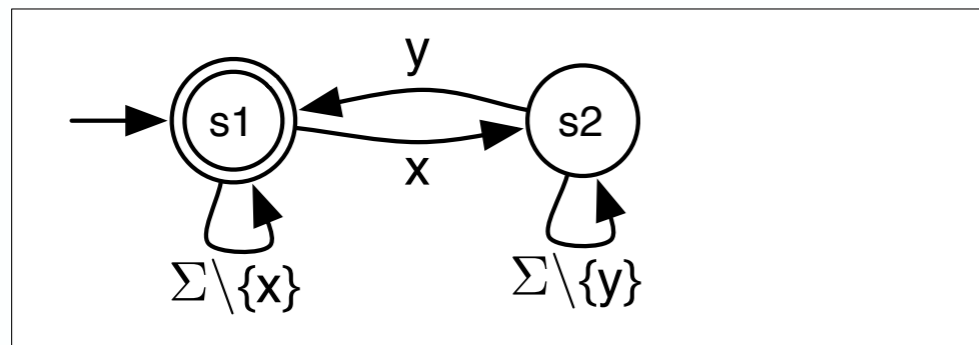
**Synoptic invariants**



# Expressing **Synoptic** with InvariMint



x AlwaysFollowedBy y

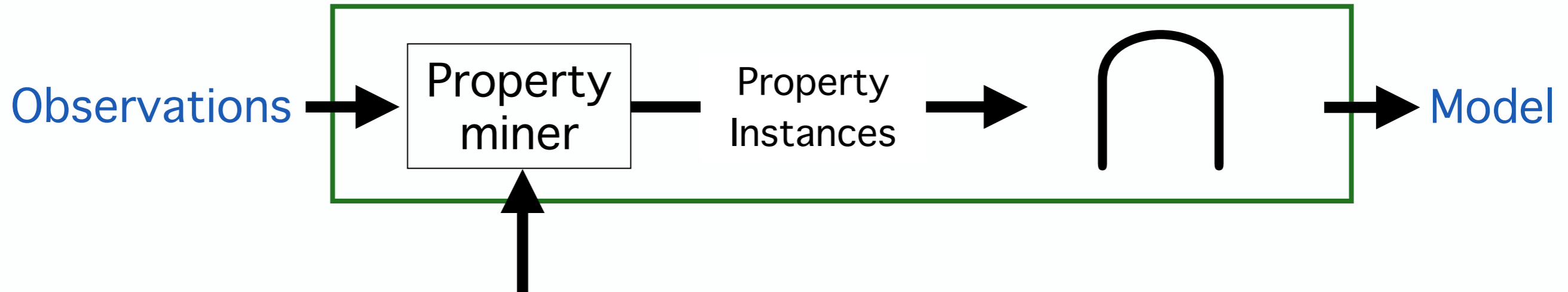


x AlwaysPrecedes y

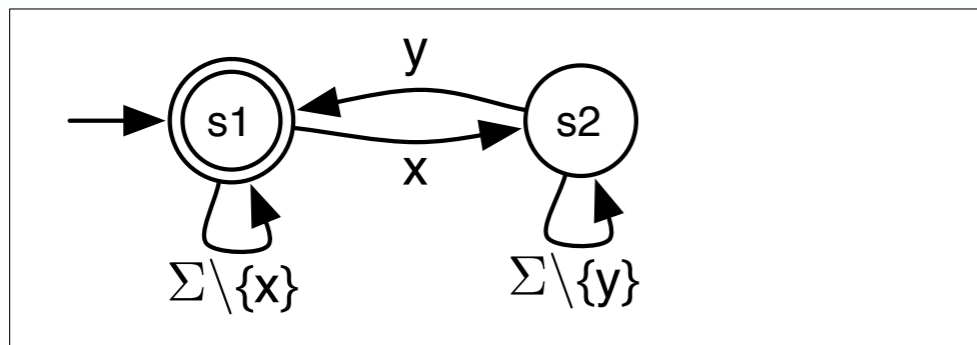
x NeverFollowedBy y

## Synoptic invariants

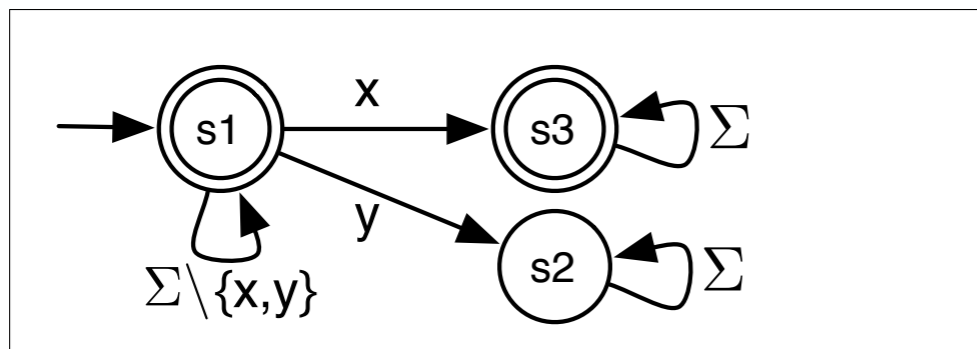
# Expressing **Synoptic** with InvariMint



**x** AlwaysFollowedBy **y**



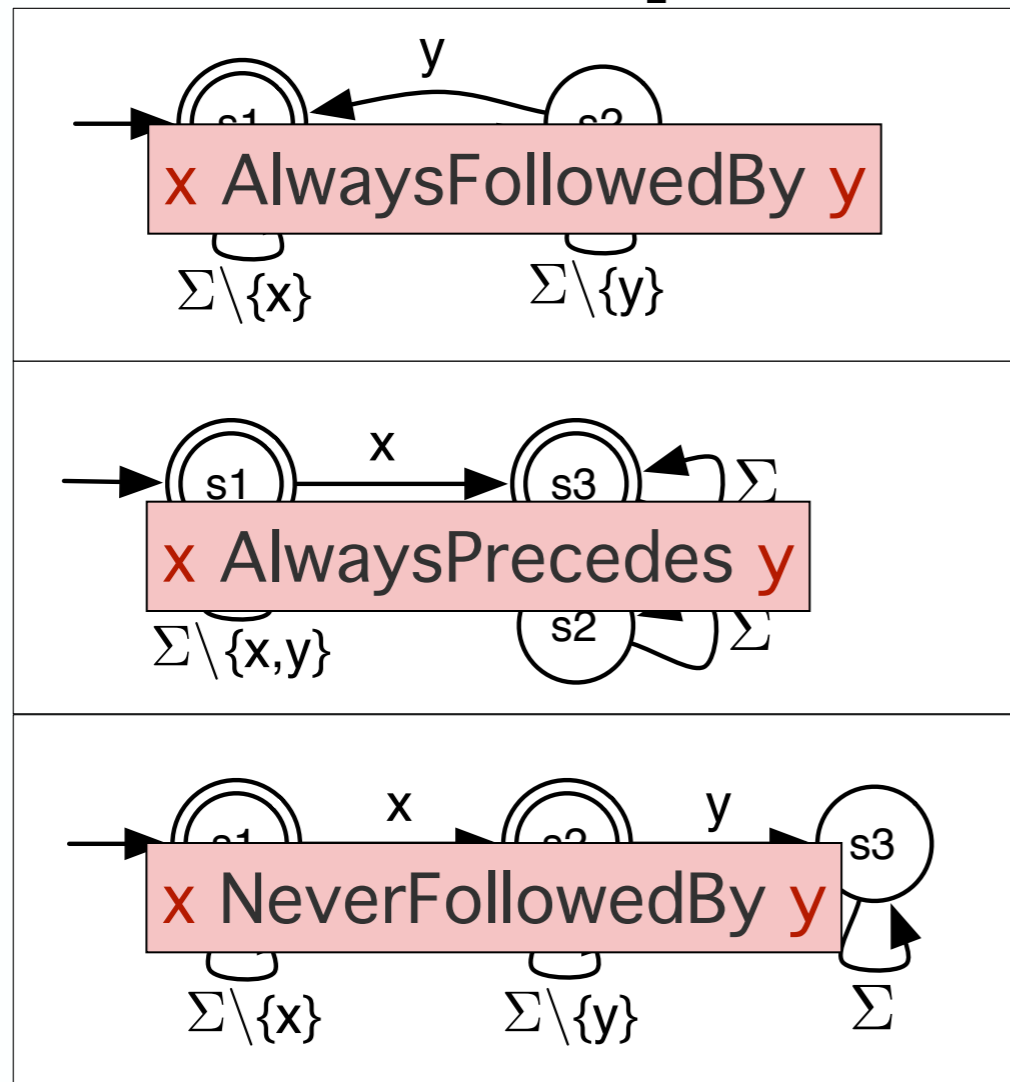
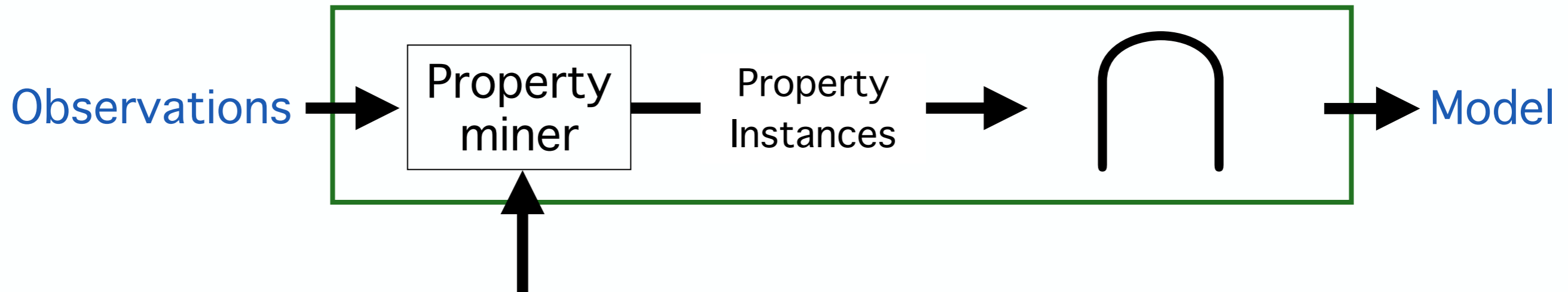
**x** AlwaysPrecedes **y**



**x** NeverFollowedBy **y**

**Synoptic invariants**

# Expressing **Synoptic** with InvariMint



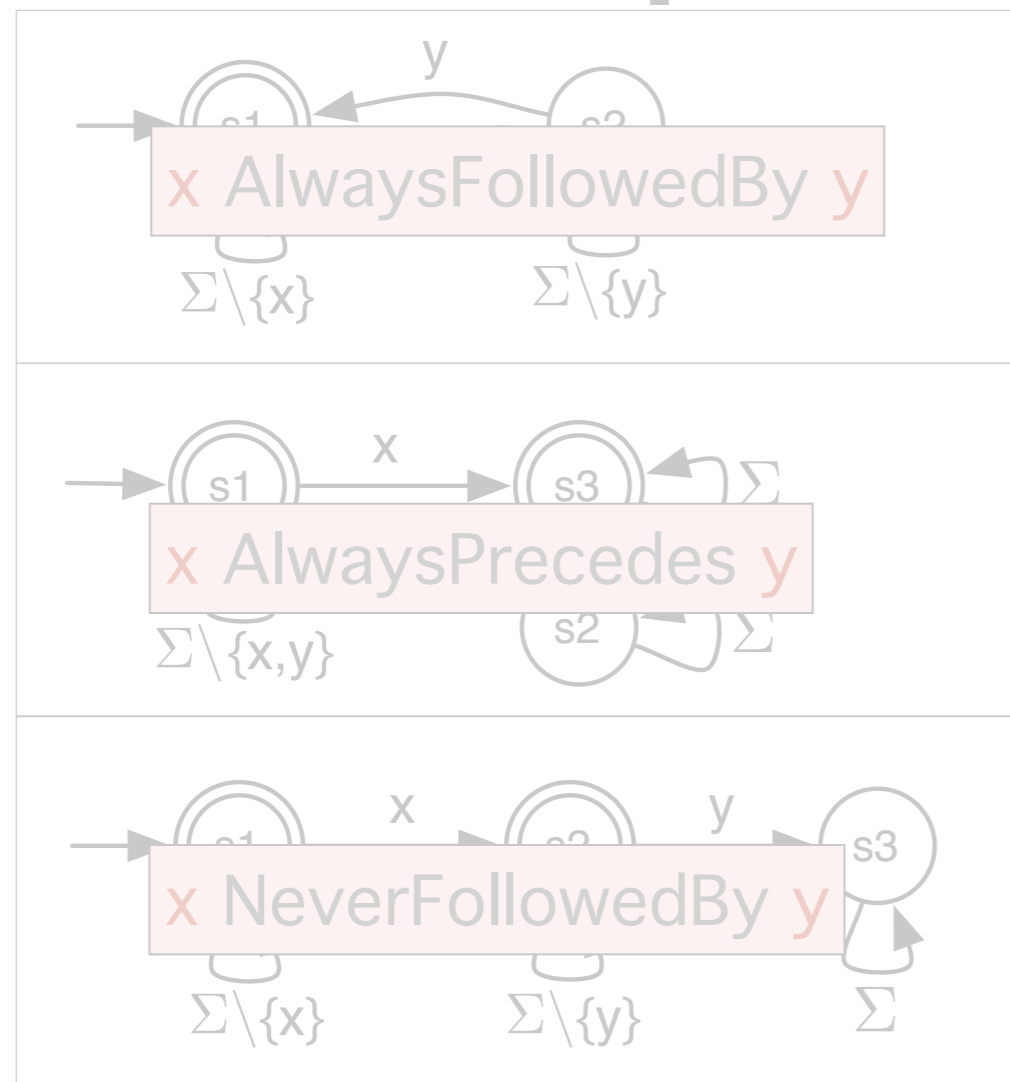
**Synoptic invariants**

# Expressing **Synoptic** with **InvariMint**

Observer

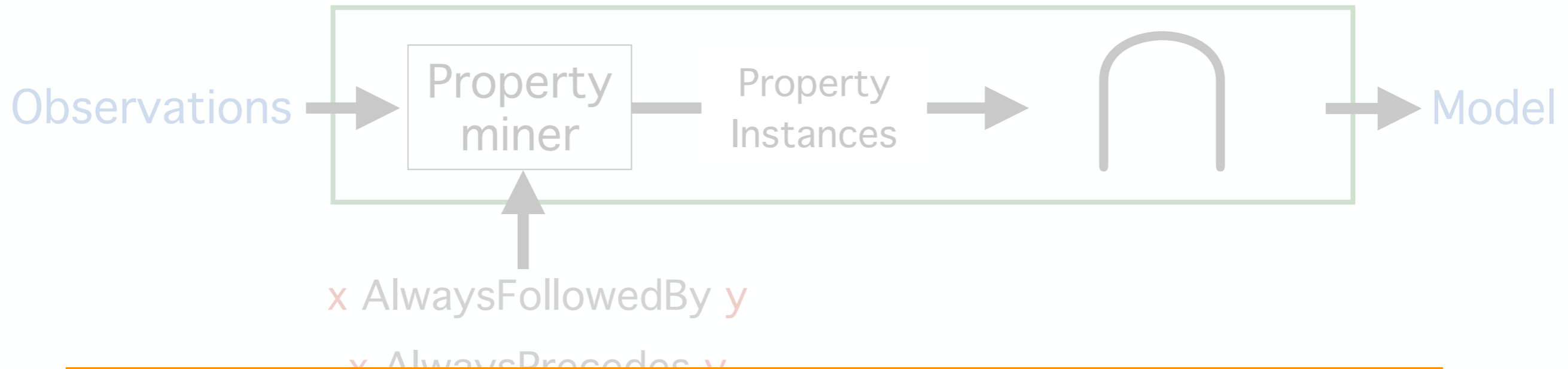
This formulation is approximate:

InvariMint model **superset** of Synoptic models



Synoptic invariants

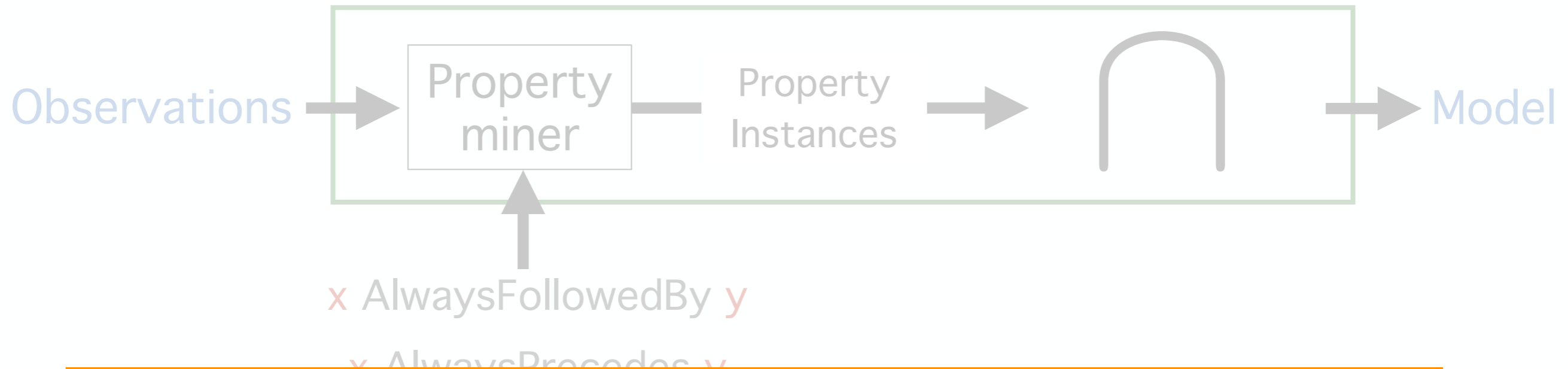
# Expressing **Synoptic** with **InvariMint**



How can we:

- ... get kTails to ignore certain k-length sequences? ✓
- ... add the x AlwaysFollowedBy y invariant to kTails?
- ... make Synoptic deterministic?
- ... add a new kind of invariant to Synoptic?
- ... learn which properties kTails/Synoptic preserve?

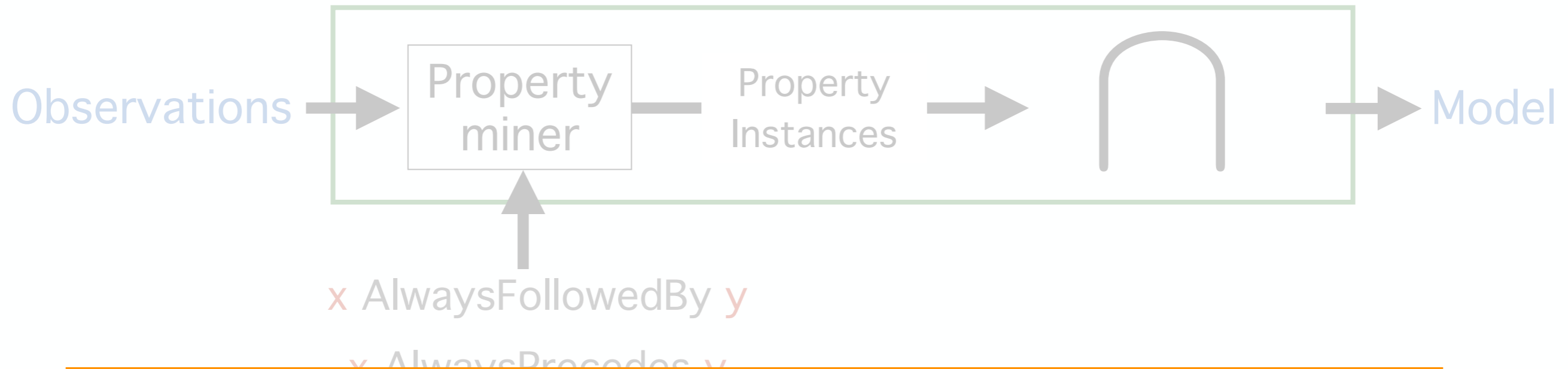
# Expressing **Synoptic** with **InvariMint**



How can we:

- ... get kTails to ignore certain k-length sequences? ✓
- ... add the x AlwaysFollowedBy y invariant to kTails? ✓
- ... make Synoptic deterministic?
- ... add a new kind of invariant to Synoptic?
- ... learn which properties kTails/Synoptic preserve?

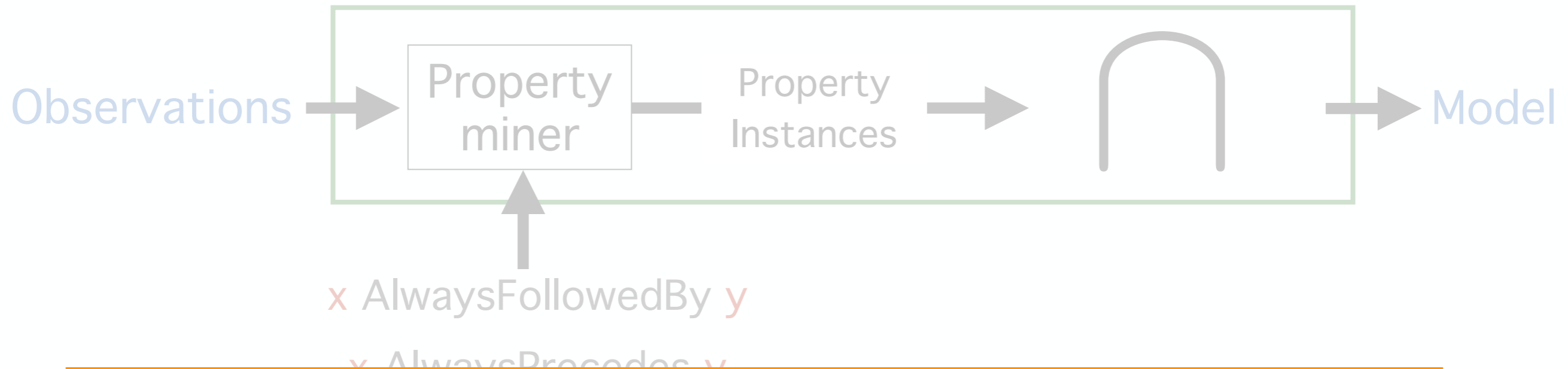
# Expressing **Synoptic** with **InvariMint**



How can we:

- ... get kTails to ignore certain k-length sequences? ✓
- ... add the x AlwaysFollowedBy y invariant to kTails? ✓
- ... **make Synoptic deterministic?**
- ... add a new kind of invariant to Synoptic?
- ... learn which properties kTails/Synoptic preserve?

# Expressing **Synoptic** with **InvariMint**

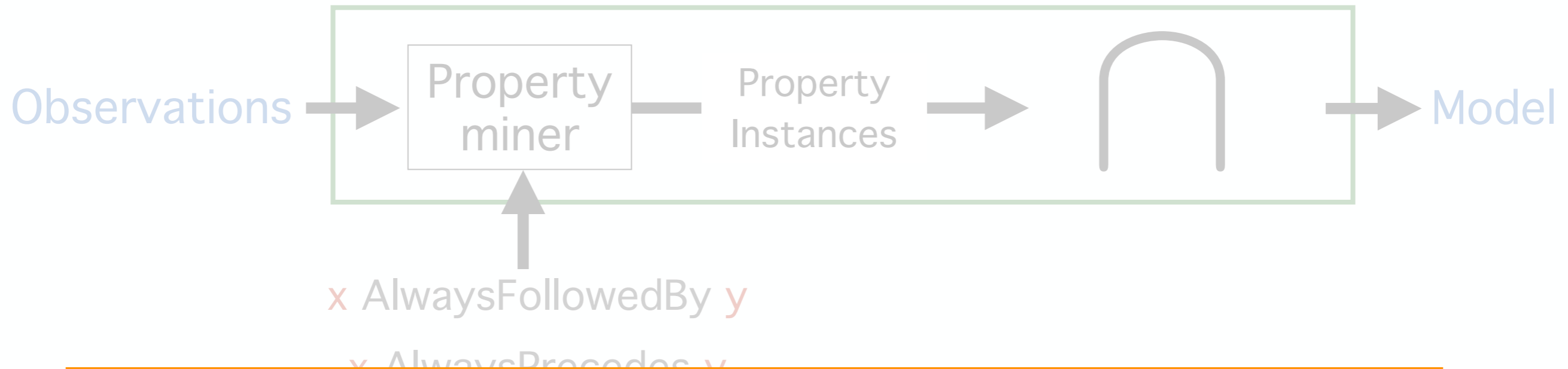


How can we:

- ... get kTails to ignore certain k-length sequences? ✓
- ... add the x AlwaysFollowedBy y invariant to kTails? ✓
- ... **make Synoptic deterministic?** ✓
- ... add a new kind of invariant to Synoptic?
- ... learn which properties kTails/Synoptic preserve?



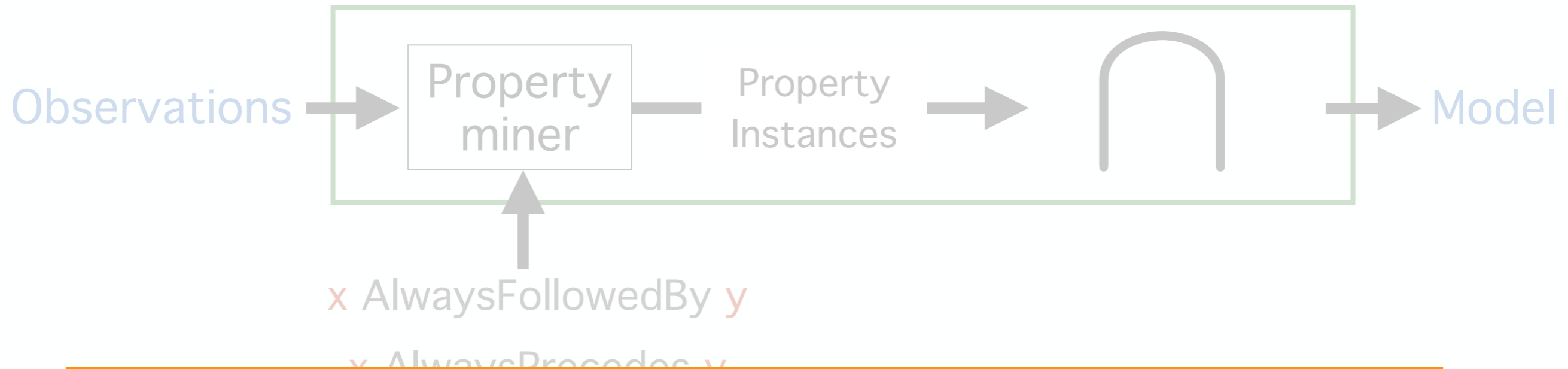
# Expressing **Synoptic** with **InvariMint**



How can we:

- ... get kTails to ignore certain k-length sequences? ✓
- ... add the x AlwaysFollowedBy y invariant to kTails? ✓
- ... make Synoptic deterministic? ✓
- ... **add a new kind of invariant to Synoptic?**
- ... learn which properties kTails/Synoptic preserve?

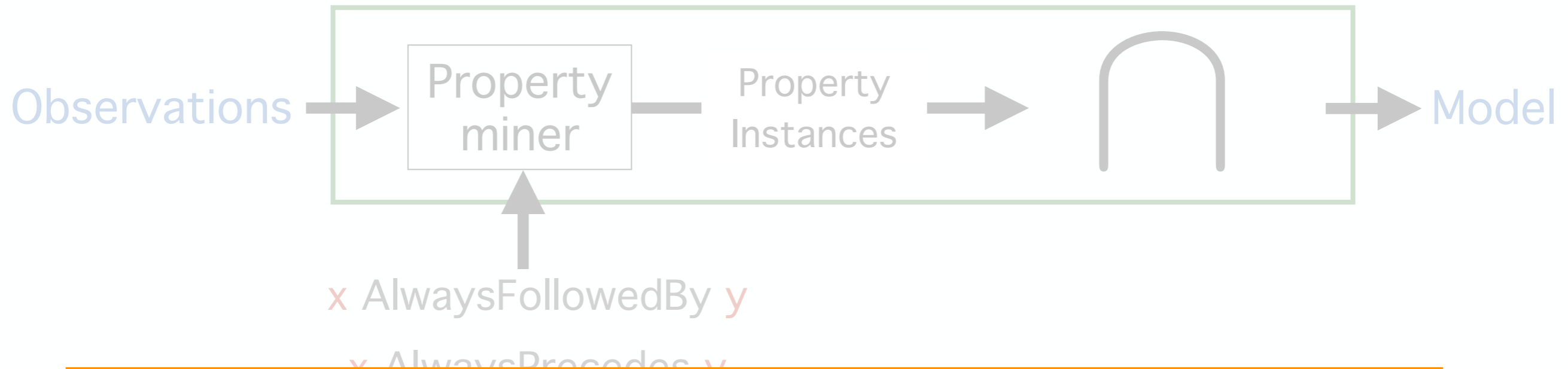
# Expressing **Synoptic** with **InvariMint**



How can we:

- ... get kTails to ignore certain k-length sequences? ✓
- ... add the x AlwaysFollowedBy y invariant to kTails? ✓
- ... make Synoptic deterministic? ✓
- ... **add a new kind of invariant to Synoptic? ✓**
- ... learn which properties kTails/Synoptic preserve?

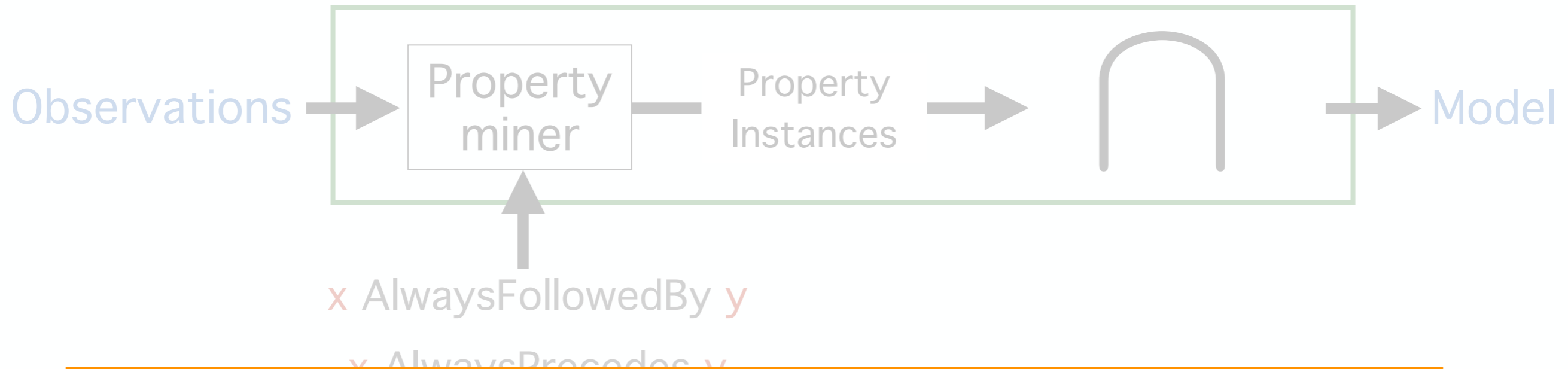
# Expressing **Synoptic** with **InvariMint**



How can we:

- ... get kTails to ignore certain k-length sequences? ✓
- ... add the x AlwaysFollowedBy y invariant to kTails? ✓
- ... make Synoptic deterministic? ✓
- ... add a new kind of invariant to Synoptic? ✓
- ... learn which properties kTails/Synoptic preserve?

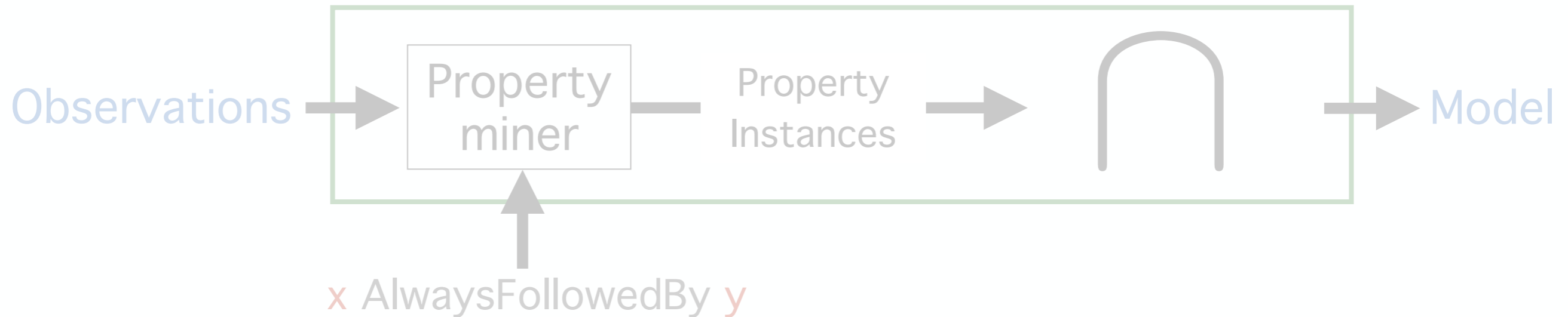
# Expressing **Synoptic** with **InvariMint**



How can we:

- ... get kTails to ignore certain k-length sequences? ✓
- ... add the x AlwaysFollowedBy y invariant to kTails? ✓
- ... make Synoptic deterministic? ✓
- ... add a new kind of invariant to Synoptic? ✓
- ... learn which properties kTails/Synoptic preserve? ✓

# Expressing **Synoptic** with **InvariMint**



How can we:

... get kTails from a Synoptic instance? 

... add the Synoptic instance to kTails?

... make Synoptic transparent?

... add a Synoptic instance to kTails?

... learn which properties kTails/Synoptic preserve?

## InvariMint advantages:

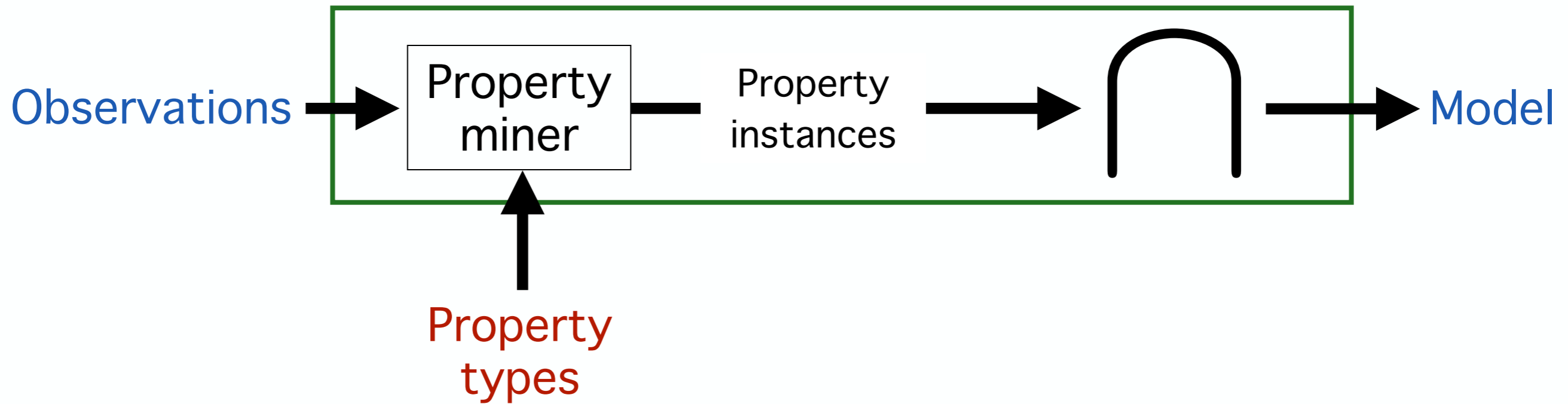
✓ Transparency

✓ Extensibility

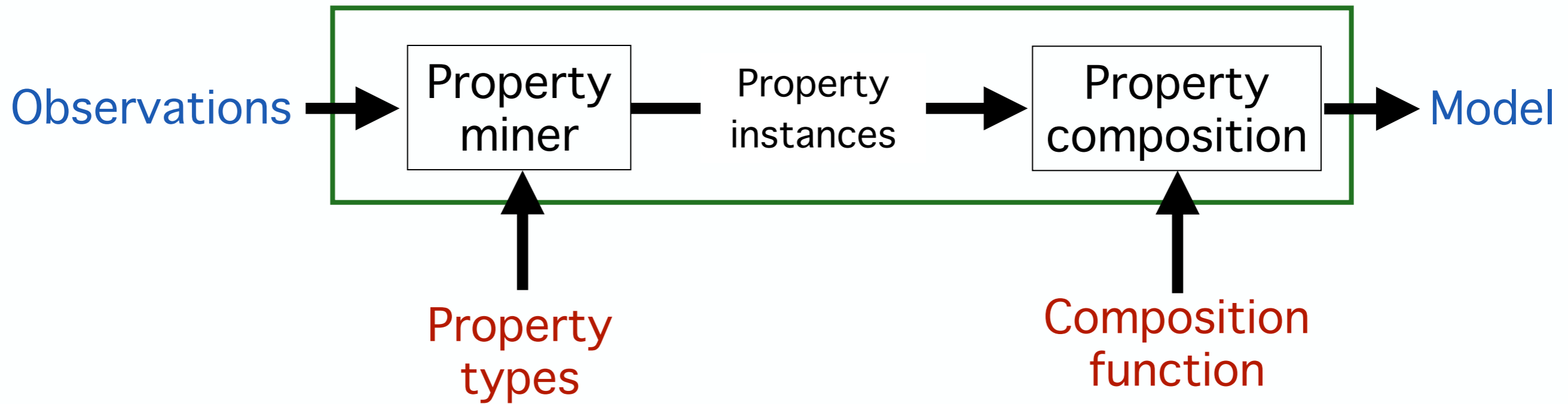
# Outline

- FSM-inference algorithms not transparent/extensible
  - ▶ **InvariMint** -- modular and declarative FSM-inference
- Expressing algorithm using **InvariMint**
  - ▶ kTails -- a state-merging algorithm
  - ▶ Synoptic -- a state-splitting algorithm
- General specification formalism
- **InvariMint** evaluation

# InvariMint specifications

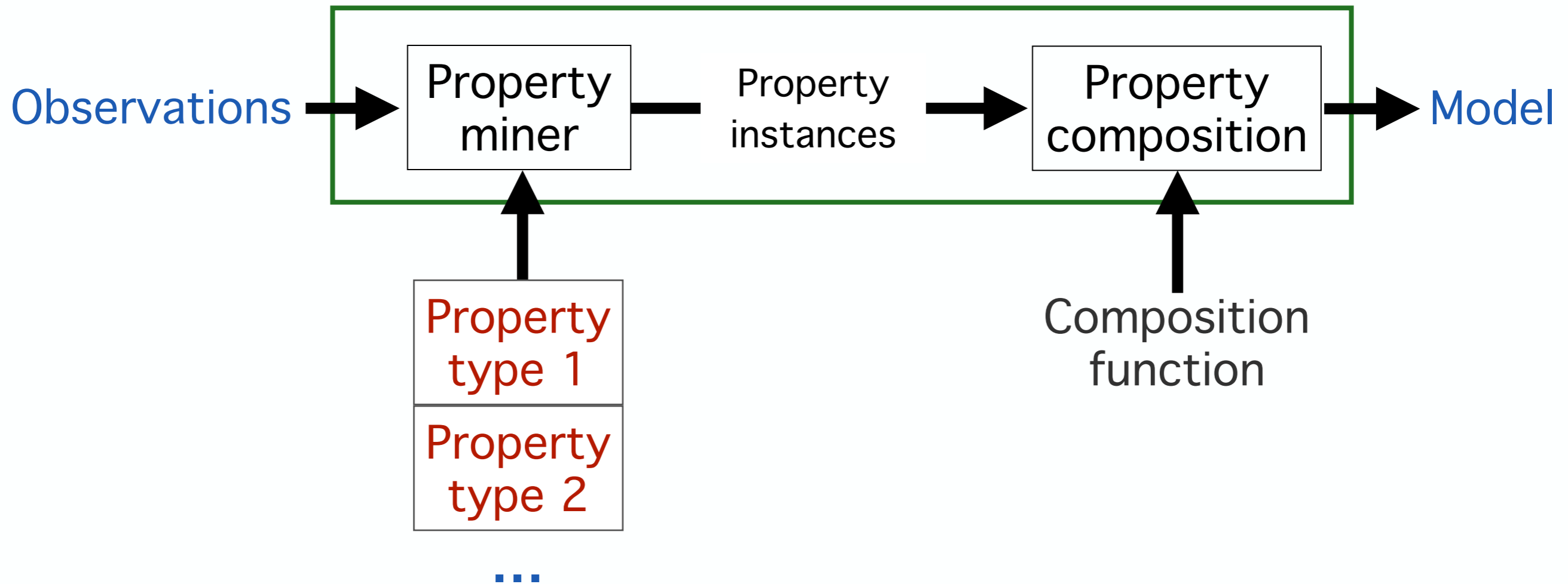


# InvariMint specifications

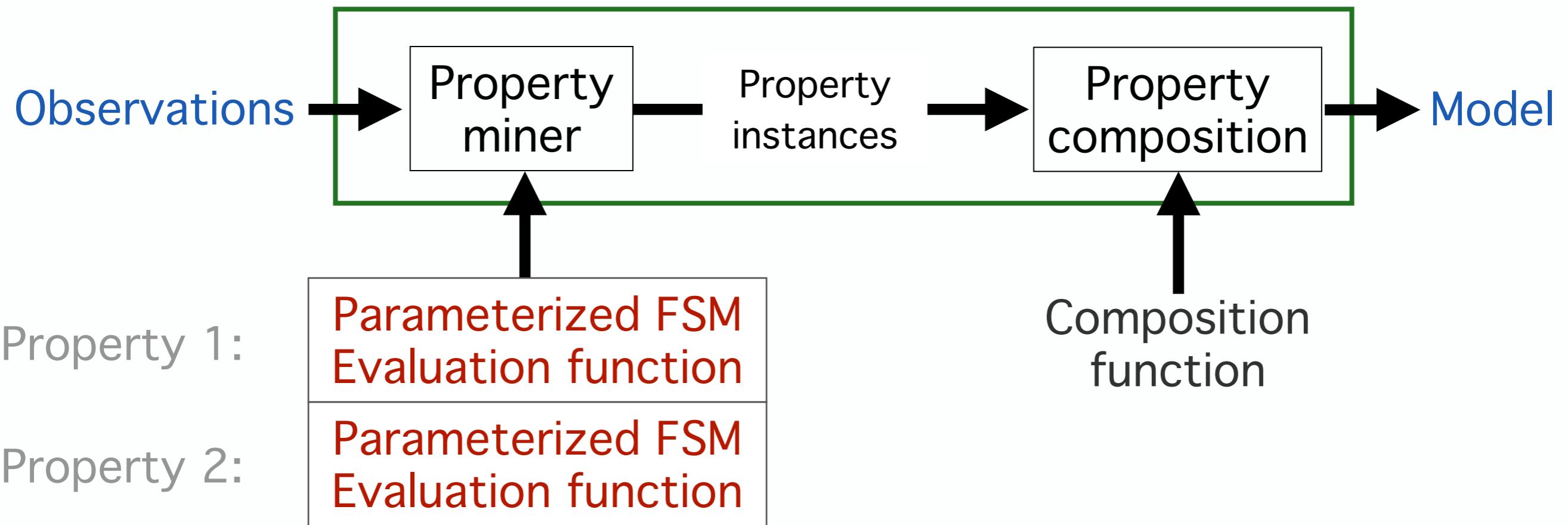




# InvariMint specifications

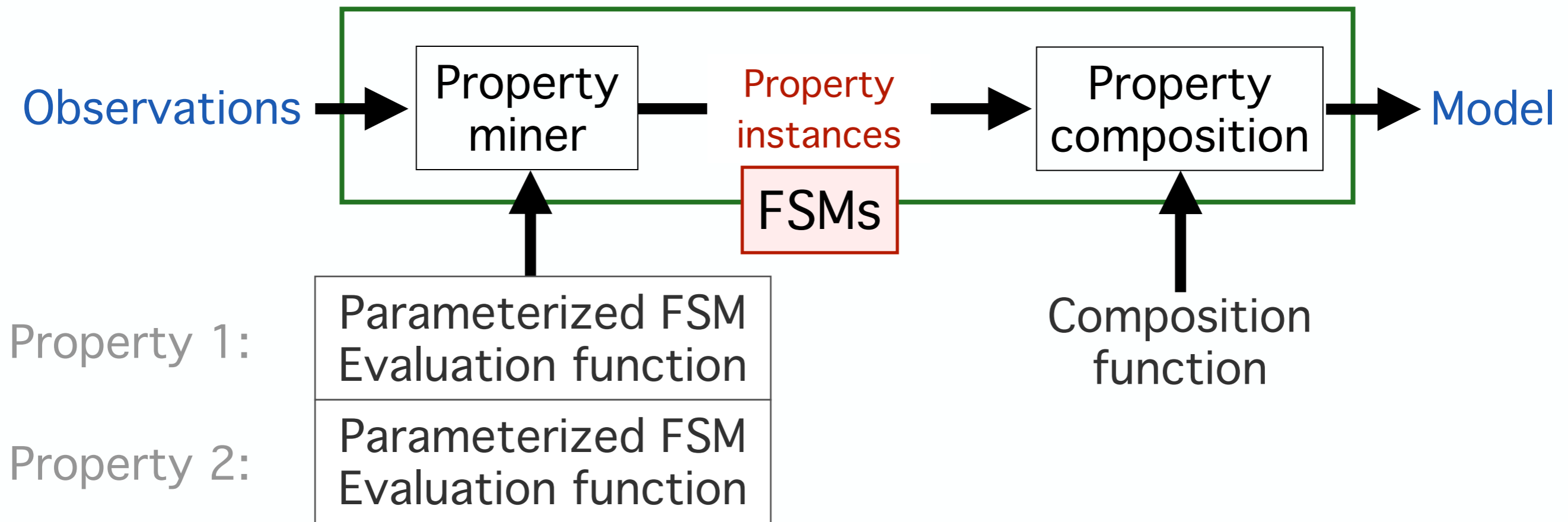


# InvariMint specifications



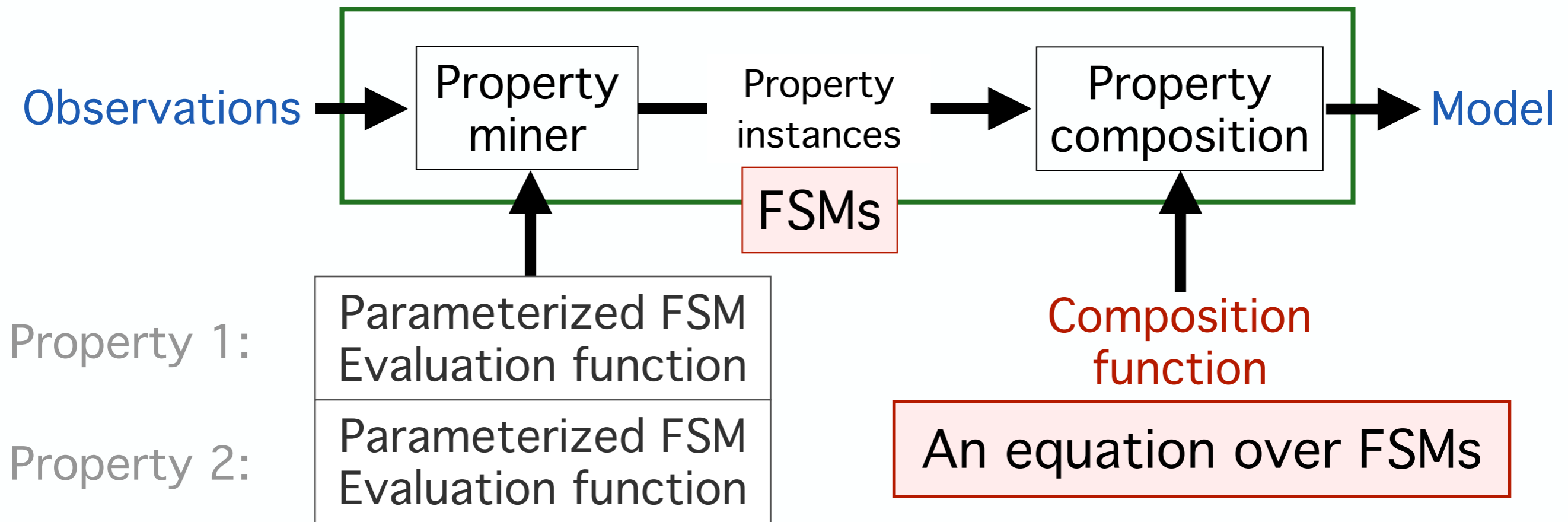
- Parameterized FSM (PFSM) -- property template
  - An FSM with variable-labeled transitions
- Evaluation function
  - Controls when a PFSM is instantiated as an FSM

# InvariMint specifications



- Parameterized FSM (PFSM) -- property template
  - An FSM with variable-labeled transitions
- Evaluation function
  - Controls when a PFSM is instantiated as an FSM

# InvariMint specifications



- Parameterized FSM (PFSM) -- property template
  - An FSM with variable-labeled transitions
- Evaluation function
  - Controls when a PFSM is instantiated as an FSM

# Outline

- FSM-inference algorithms not transparent/extensible
  - ▶ **InvariMint** -- modular and declarative FSM-inference
- Expressing algorithm using **InvariMint**
  - ▶ kTails -- a state-merging algorithm
  - ▶ Synoptic -- a state-splitting algorithm
- General specification formalism
- **InvariMint** evaluation

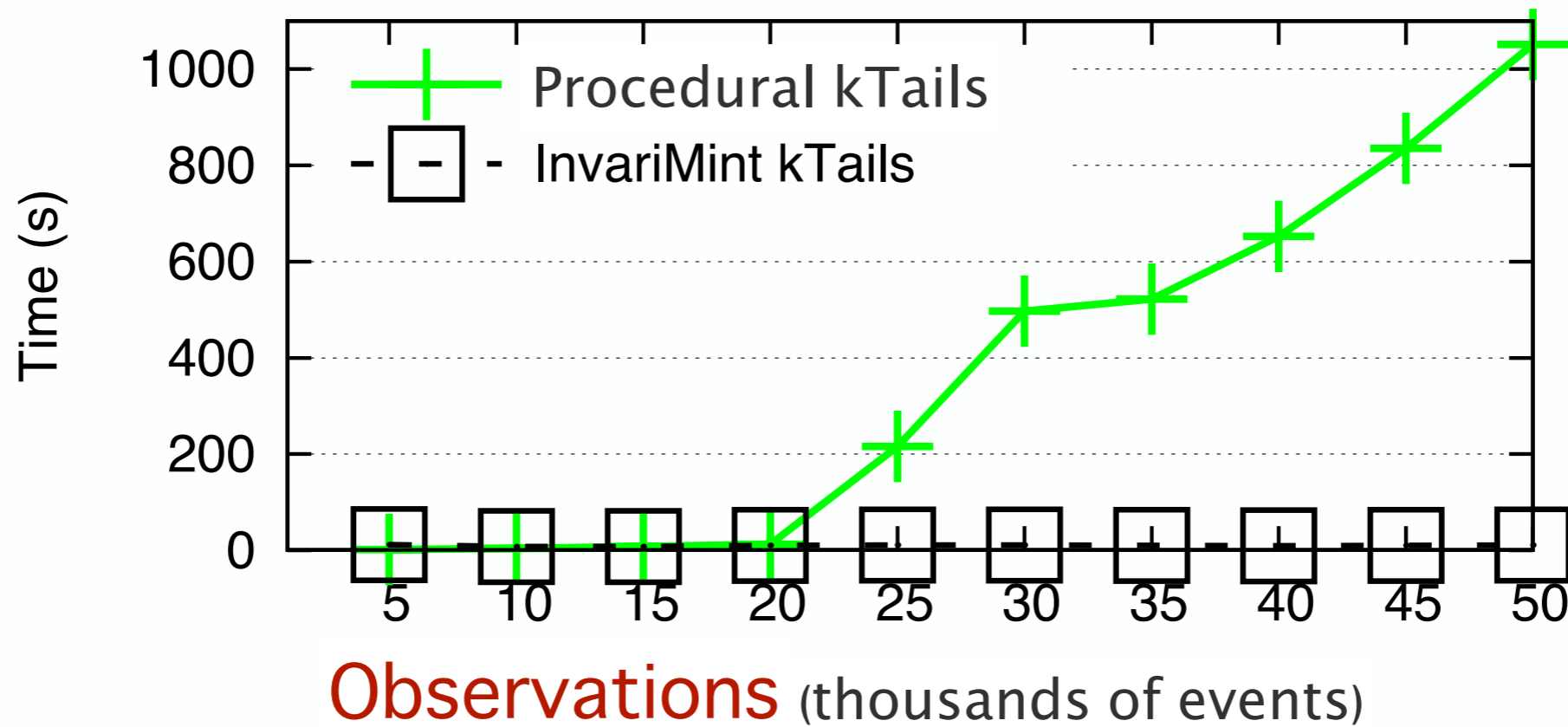
# InvariMint expressiveness

- FSM-inference algorithms from prior work
  - ▶ Expressed kTails and Synoptic
  - ▶ Useful for understanding both algorithms
  - ▶ Noticed overlap in their InvariMint specifications!
- What can and can't we express with InvariMint?
  - ▶ Limited to FSM algorithms that consider temporality
  - ▶ Constrains inference to composition of properties
  - ▶ Benefits from a formal background

# InvariMint performance

- FSM inference useful for large systems
  - ▶ Compact representation of mass executions
  - ▶ A node at Google generates a million messages/day
- Does transparency and extensibility impact performance?
  - ▶ InvariMint versions of kTails and Synoptic over **100x faster** than procedural variants

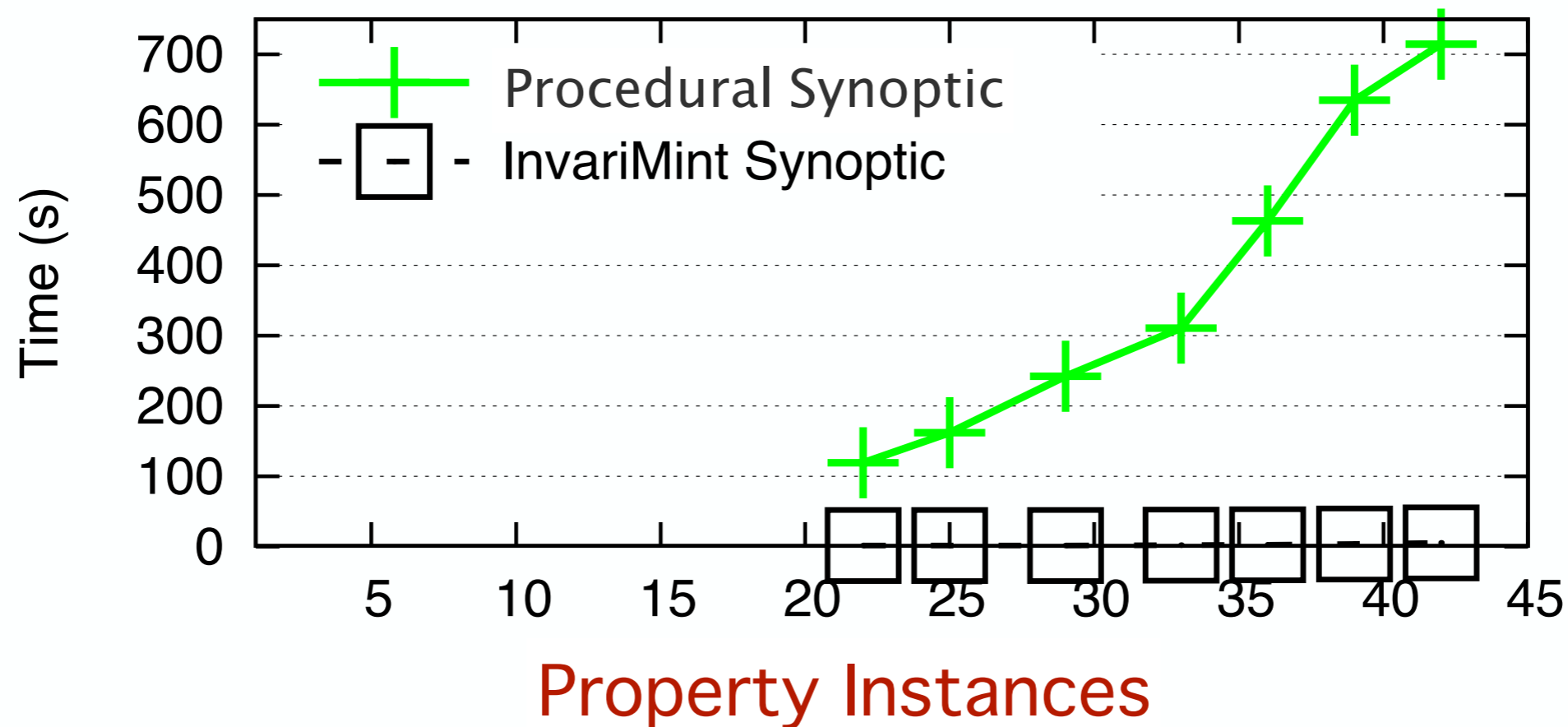
# Performance: kTails



- InvariMint kTails much faster than procedural kTails
  - ▶ Efficient FSM composition (e.g., intersection)
  - ▶ The full log is never maintained in memory



# Performance: Synoptic

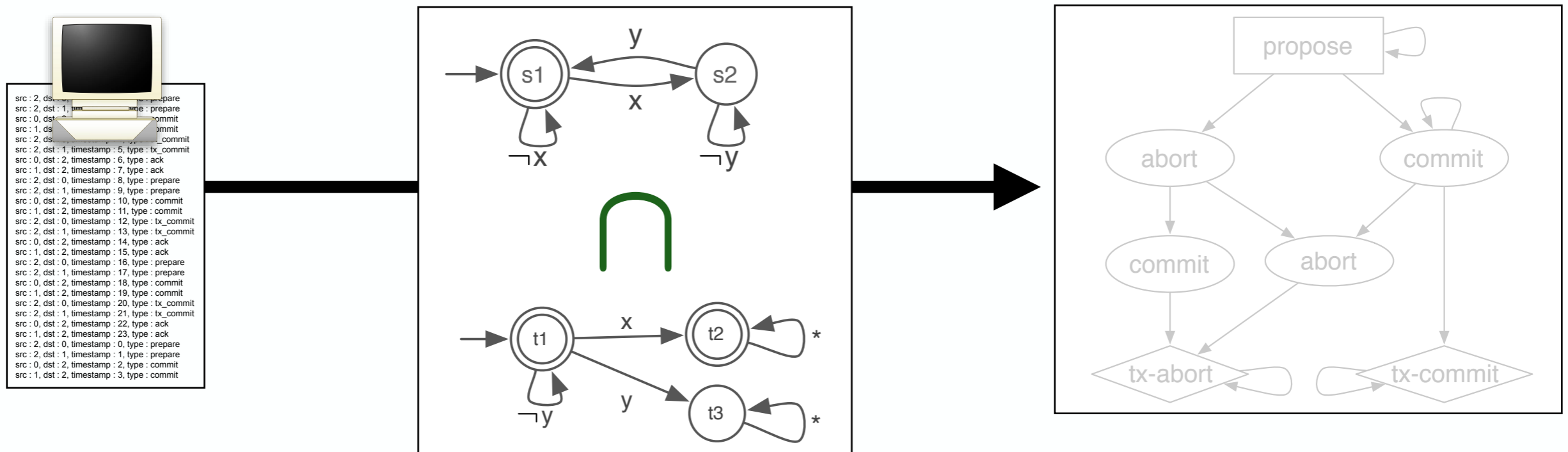


- Procedural Synoptic much slower than InvariMint Synoptic
  - ▶ Synoptic's modeling check invariants
  - ▶ Synoptic's refinement deals with concrete event instances

# Contributions

FSM-inference algorithms are **not transparent/extensible**

## InvariMint



Open source

[synoptic.googlecode.com](https://synoptic.googlecode.com)

# Contributions

FSM-inference algorithms are not transparent/extensible

## InvariMint

Declarative specification of FSM-inference algorithms

- Provides **insight** into how an algorithms works
- A **common language** for inference algorithms
- Simplifies **extension** of existing algorithms
- **Faster** than procedural equivalents

Open source

[synoptic.googlecode.com](https://synoptic.googlecode.com)